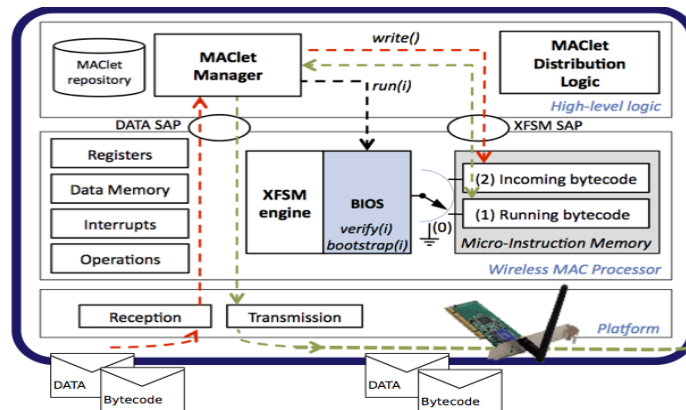


MAClets

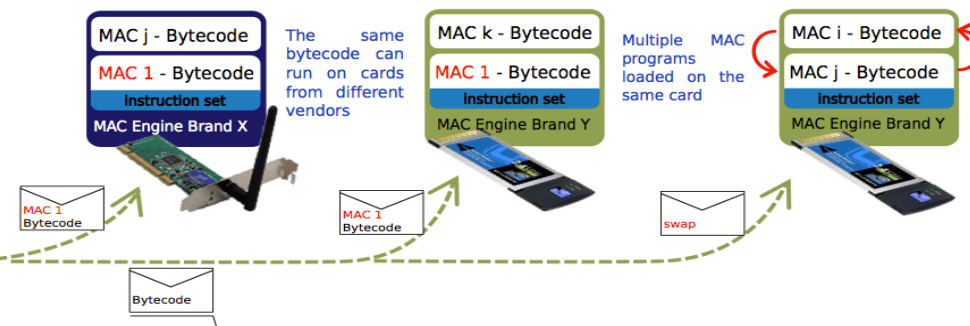
Active MAC protocols over hard-coded devices

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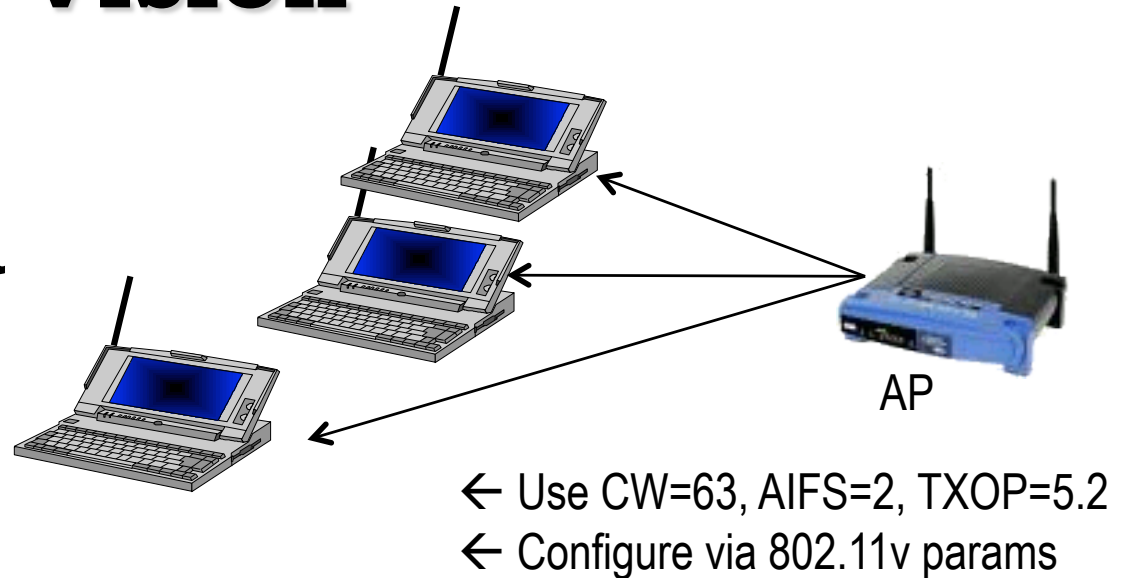


The WMP architecture permits the **design MAC once, run everywhere** paradigm and the decoupling between the platform and the MAC protocol logic.

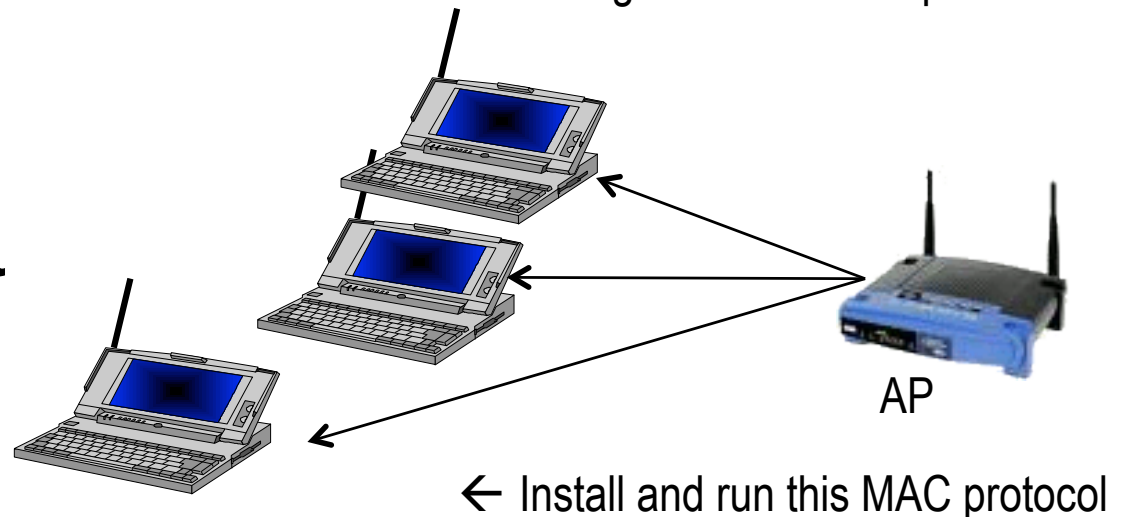


Vision

→ From STA management via parameter settings...



→ To STA management via full MAC stack reprogramming!



Whole MAC protocol stack as a sort of JAVA applet?

...more opportunities...

→ **Flexibly Adapt Access Protocol to scenario/context**

- ⇒ Dynamic spectrum access
- ⇒ Niche scenario optimization
 - home, industrial, ...
- ⇒ Context/application-specific protocol design
- ⇒ Faster paper-to-field deployment
- ⇒ Improved support for PHY enhancements

→ **Virtualization**

- ⇒ Each operator can design its own resource management
 - frame forging, scheduling, timing, channel switching, PHY selection, ...
- ⇒ Different MAC coexisting on same AP/net

Real world blockers

→ Lower MAC protocol ops are real time!

- ⇒ O(us): TX, RX, slot times, set IFS, set timers, etc
- ⇒ Driver to NIC interface: too slow → MUST run on NIC

→ Vendors will **HARDLY** give us open source, fully programmable, NICs

- ⇒ SDR is 20 years old but...
 - ...still no real world commodity SDR NICs
- ⇒ NIC design extensively leveraging HW
 - non programmable, unless FPGA NICs...
 - Your commodity card is NOT an FPGA!
- ⇒ Why a vendor should renounce to its internal Intellectual Property??

→ **But even if stack gets opened...which programmability model?**

- ⇒ Current practice (in most cases):
 - *patch/hack existing SW/FW/HW code base*
 - Huge skills/experience, low level languages, slow development, inter-module dependencies

Our contribution

→ Exploiting a new abstraction model for run-time MAC protocol reconfigurations!

⇒ based on the Wireless MAC Processor (WMP)

→ INFOCOM 2012

→ Enabling active MAC protocols and remote MAC injection

⇒ Ultra-fast (below ms) reconfiguration

⇒ MAC multi-threading

⇒ Virtualization

Learn from computing systems?

→ 1: Instruction sets

perform elementary tasks on the platform

→ A-priori given by the platform

→ Can be VERY rich in special purpose computing platforms

» Crypto accelerators, GPUs, DSPs, etc

→ 2: Programming languages

sequence of such instructions + conditions

⇒ Convey desired platform's operation or algorithm

→ 3: Central Processing Unit (CPU)

execute program over the platform

⇒ Unaware of what the program specifically does

⇒ Fetch/invoke instructions, update registers, etc

Clear decoupling between:

- | | |
|---------------------|---|
| - platform's vendor | → implements (closed source!) instruction set & CPU |
| - programmer | → produces SW code in given language |

1: Which elementary MAC tasks?

(“our” instruction set!)

→ ACTIONS

⇒ **frame management, radio control, time scheduling**

→ TX frame, set PHY params, RX frame,
set timer, freeze counter, build header,
forge frame, switch channel, etc

→ EVENTS

⇒ **available HW/SW signals/interrupts**

→ Busy channel signal, RX indication,
inqueued frame, end timer, etc

→ CONDITIONS

⇒ **boolean/arithmetic tests on available registers/info**

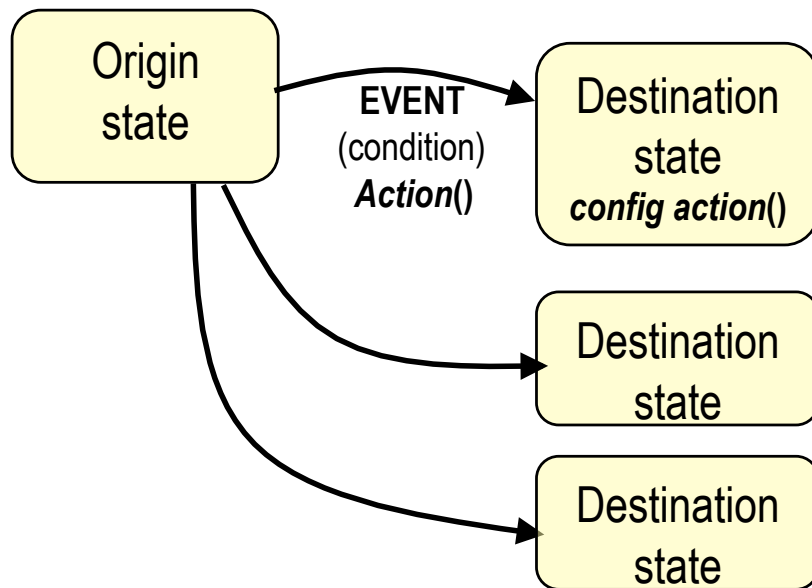
→ Frame address == X, queue length > 0,
ACK received, power level < P, etc

2: How to compose MAC tasks?

(“our” programming language!)

→ Convenient “language”: XFSM **eXtended Finite State Machines**

⇒ Compact way for composing available acts/ev/cond
to form a custom MAC protocol logic



XFSM formal notation		meaning
S	symbolic states	MAC protocol states
I	input symbols	Events
O	output symbols	MAC actions
D	n-dimensional linear space $D_1 \times \dots \times D_n$	all possible settings of n configuration registers
F	set of enabling functions $f_i : D \rightarrow \{0, 1\}$	Conditions to be verified on the configuration registers
U	set of update functions $u_i : D \rightarrow D$	Configuration commands, update registers' content
T	transition relation $T : S \times F \times I \rightarrow S \times U \times O$	Target state, actions and configuration commands associated to each transition

3: How to run a MAC program?

(MAC engine – XFSM onboard executor - our CPU!)

→ **MAC engine: specialized XFSM executor** *(unaware of MAC logic)*

⇒ Fetch state

⇒ Receive events

⇒ Verify conditions

⇒ Perform actions and state transition

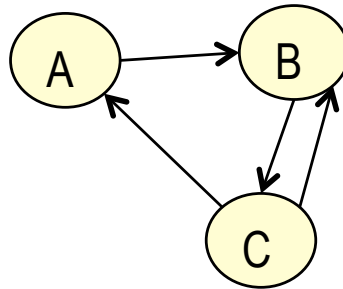
→ **Once-for-all “vendor”-implemented in NIC** *(no need for open source)*

⇒ “close” to radio resources = straightforward real-time handling

MAC Bytecode

→ MAC description:

⇒ XFSM



→ XFSM → tables

	A	B	C
A		T(A,B)	
B			T(B,C)
C	T(C,A)	T(C,B)	

→ Transitions

⇒ «byte»-code event, condition, action

→ **Portable over different vendors' devices, as long as API is the same!!**

⇒ Pack & optimize in WMP «machine-language» bytecode

A	T(A,B)	
B	T(B,C)	
C	T(C,A)	T(C,B)

MAC protocol specification:
XFSM design
(e.g. Eclipse GMF)

Machine-readable code

Custom language compiler

Code injection
in radio HW platform

MAC Bytecode

MAC Engine



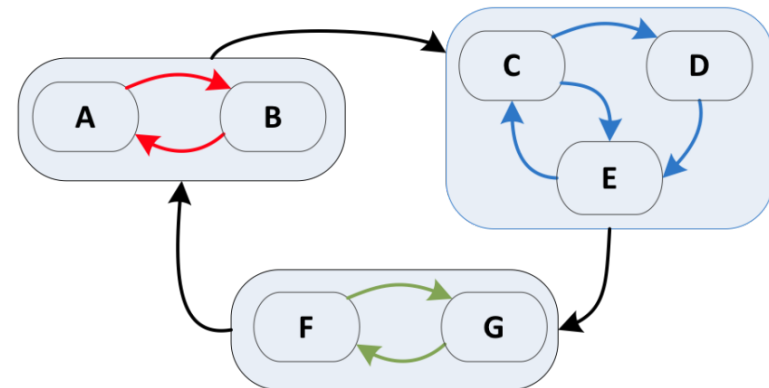
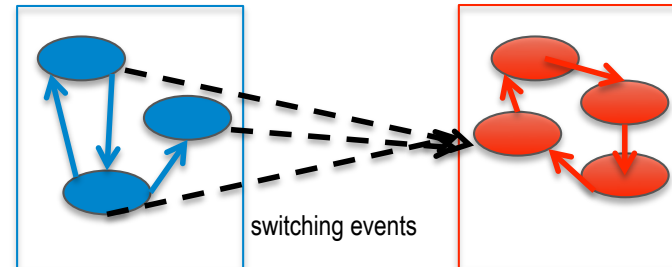
Multi-Thread Support

→ **The MAC Engine does not need to know to which MAC program a new fetched state belongs!**

⇒ Code switching can be easily supported by moving to a state in a different transition table

→ **It is enough to:**

- ⇒ Define Meta State Machines for programming code switching
- ⇒ Verify MAC switching events from each state of the program under execution
- ⇒ Re-load system configuration registers at MAC transitions



From MAC Programs to MAClets

→ Upload MAC program on NIC from remote

⇒ While another MAC is running

⇒ Embed code in ordinary packets

→ WMP Control Primitives

⇒ load(XFSM)

⇒ run(XFSM)

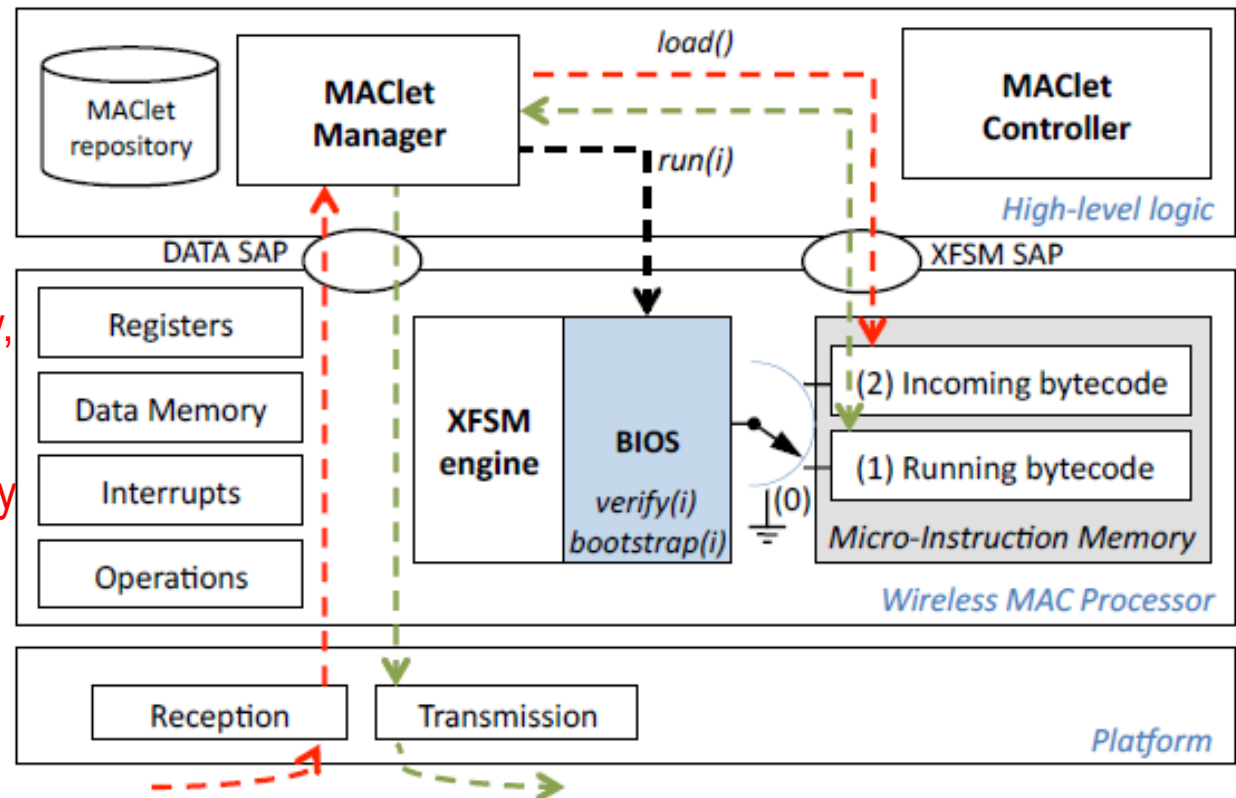
⇒ verify(XFSM)

⇒ switch(XFSM1, XFSM2, ev, cond)

→ Further primitives

⇒ Distribution protocol (run by the MAClet Manager)

⇒ Synchro support for distributed start of same MAC operation)



“Bios” state machine: DEFAULT protocol (e.g. wifi) which all terminals understand

MAClets

→ **An entire MAC program can be coded in a single frame!**

⇒ our abstractions and machine codes allow to code DCF in about 500 bytes

→ **Other fields:**

- ⇒ type (distribution protocol and action messages)
- ⇒ destination IDs
- ⇒ initial state
- ⇒ command (load, run, switch..)
- ⇒ activation event

Memory address	Memory Initial State Descriptor										Description
0x0BC0:	0100	FFFF	0800	0014	A5FF	6ADA	0014	A5FF			
0x0BD0:	6ADA	6C00	80A4	FF00	FF00	3600	80EE	FF00			
0x0BE0:	FF00	0000	80BB	FF00	FF00	0600	2C01	0600			
0x0BF0:	0000	0000	0000	0000	0000	0000	0000	0000			

	00	01	02	03	Coded state machine						
0x0C00:	0100	0100	0100	0401	0108	0508	1C01	010B			
0x0C10:	010B	3001	010D	0200	FFFF	5101	010E	030D			
0x0C20:	0000	0100	010F	C100	0102	0602	E100	0106			
0x0C30:	0106	0401	0108	0508	1C01	010B	030B	FFFF			
0x0C40:	CD00	0104	080C	0000	010B	0D00	FFFF	0E01			
0x0C50:	0109	0909	1C01	010B	0D0B	FFFF	C700	0103			
0x0C60:	0C03	E100	0106	0106	FFFF	A601	0110	1600			
0x0C70:	0000	0100	0100	FFFF	0E01	0109	0109	1C01			
0x0C80:	010B	010B	FFFF	5F01	010F	0A00	0000	0100			
0x0C90:	0D00	FFFF	C100	0102	0A02	C700	0103	0B03			
0x0CA0:	E100	0106	0D06	FFFF	D300	0105	0D05	E100			
0x0CB0:	0106	0D06	FFFF	D300	0105	0705	E100	0106			
0x0CC0:	0106	FFFF	6D01	0111	1800	0000	0100	0100			
0x0CD0:	0000	0100	0D10	7401	0112	1512	0000	0100			
0x0CE0:	1111	9601	0113	0513	0000	0100	0500	0000			
0x0CF0:	0100	0304	E100	0106	1206	0401	0108	0508			
0x0D00:	1C01	010B	120B	FFFF	A901	0115	0100	E401			
0x0D10:	0117	1200	0000	0100	0100	0000	0100	1214			
0x0D20:	B901	0118	0310	0000	0100	0300	0401	010B			
0x0D30:	1708	1501	010A	010A	1C01	010B	010B	C501			
0x0D40:	0119	0800	0000	0100	0500	3001	010D	0200			
0x0D50:	0401	0108	0508	1C01	010B	180B	CB01	011A			
0x0D60:	0200	0000	0100	0100	0000	0000	0000	0000			
0x0D70:	0000	0000	0000	0000	0000	0000	0000	0000			
0x0D80:	0000	0000	0000	0000	0000	0000	0000	0000			

	00	01								
0x0D90:	00F0	03FE	08F2	13FB	20FB	27FB	28FB	35FB		
0x0DA0:	3CFE	43FE	4AFE	54FE	5BFE	62F2	68F0	6BF2		
0x0DB0:	71F2	77F0	7AFE	84F4	8DF0	90F2	96F4	9FF2		
0x0DC0:	A5F4	A8F2	0000	0000	0000	0000	0000	0000		
0x0DD0:	0000	0000	0000	0000	0000	0000	0000	0000		

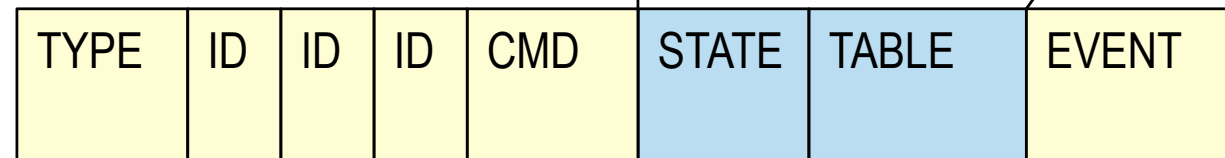
Outgoing transitions for state 01
 0401 0108 0508 = trans. 1
 1C01 010B 010B = trans. 2
 3001 010D 0200 = trans. 3
 FFFF = delimiter

Transition 1
 0401 = event pointer
 01 = event parameter
 08 = event index
 05 = target state
 08 = action

State 01
 03 = transitions offset (9 bits)
 E = FFFF delimiter



MACLET



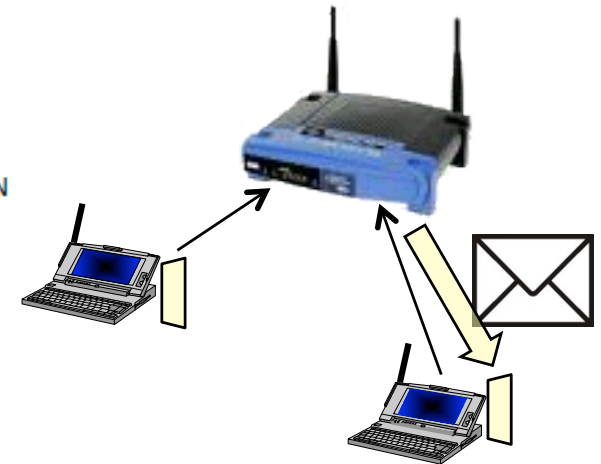
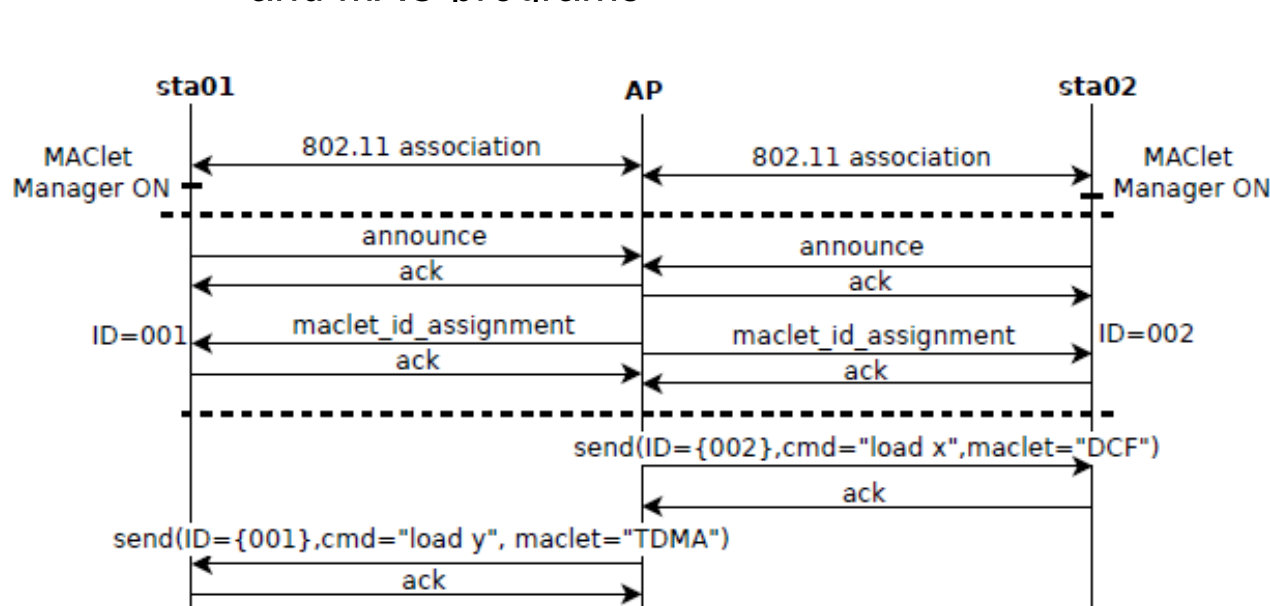
MAClet Distribution Protocol

→ Defined for allow the AP to remotely access the WMP control interface of the associated nodes

⇒ Binding MAClet Managers of each node to the AP MAClet Controller

→ Notification of activation/de-activation, ID assignment

⇒ Transporting Action Messages coding WMP commands (load/run/switch) and MAC programs



Synchronization Primitives

→ When to switch to a new MAC protocol?

⇒ Mechanisms available, but final solutions left to the MAClet programmers

→ Triggering events and signals

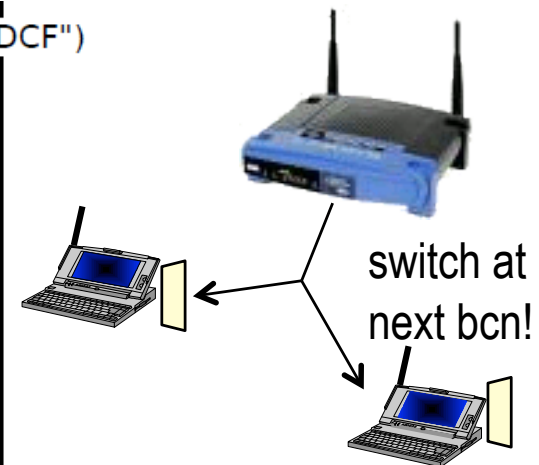
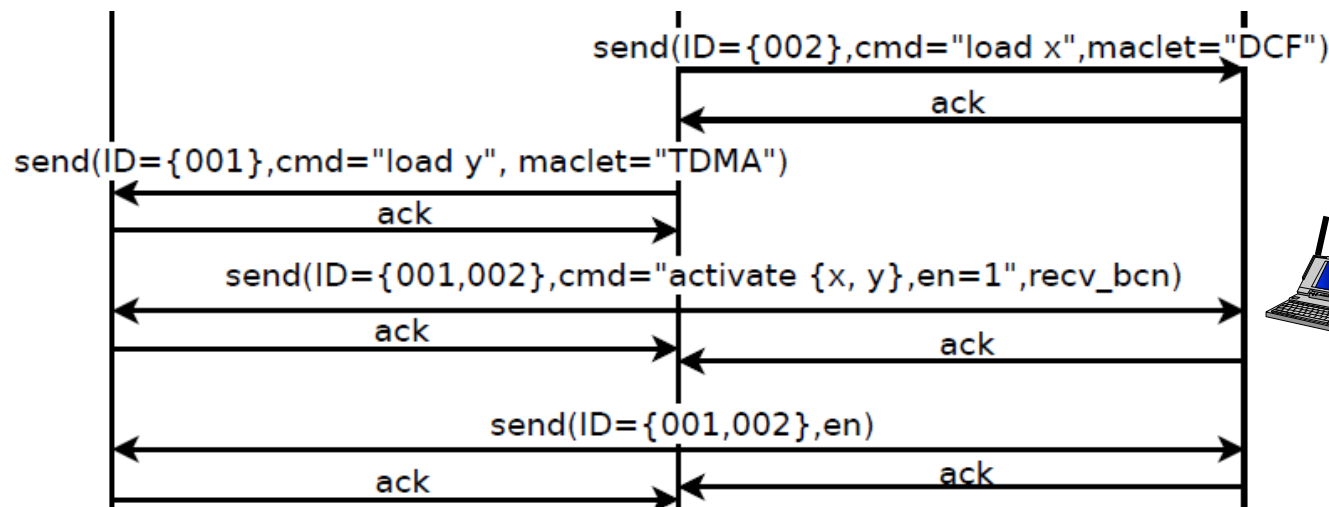
⇒ No trigger: asynchronous activation

⇒ Control frames sent by the AP

⇒ Expiration of relative or absolute timer

→ Absolute timers built on top of the time-synchronization function included in DCF

⇒ 1-way or 3-way handshakes

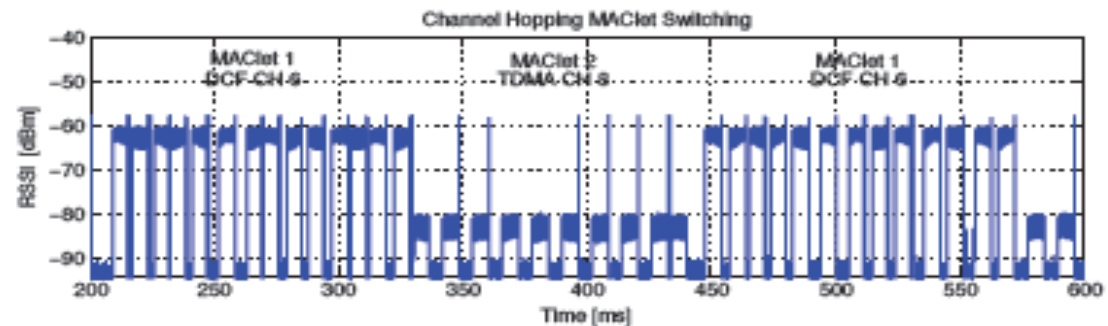
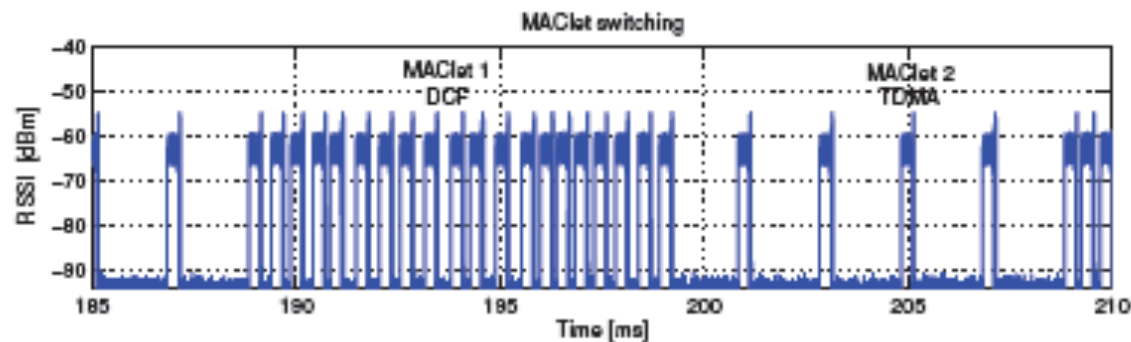


Switching Operation

→ From a configuration to another..

→ From a program to another!

⇒ (with latency of about 1 microsecond)



Implementation at a glance

(on commodity hardware!)

Reference platform: broadcom Airforce54g 4311/4318

→ WMP:

- ⇒ replace both Broadcom and openFWWF firmware with
 - Implementation of actions, events, conditions
 - MAC engine: XFSM executor
- ⇒ Develop “machine language” for MAC engine
 - Custom made “bytecode” specified and implemented

→ WMP Control Architecture:

- ⇒ At firmare level:
 - WMP Control Interface
- ⇒ At the application level:
 - MAClet Manager: receive/transmit MAClets and other messages of the MAClet Distribution Protocol
 - MAClet Controller: Intelligent part of the system, dealing with network-level decisions
 - Current implementation based on classical client-server model!

Application Examples

AP Virtualization with MAClets

→ Two operators on same AP/infrastructure

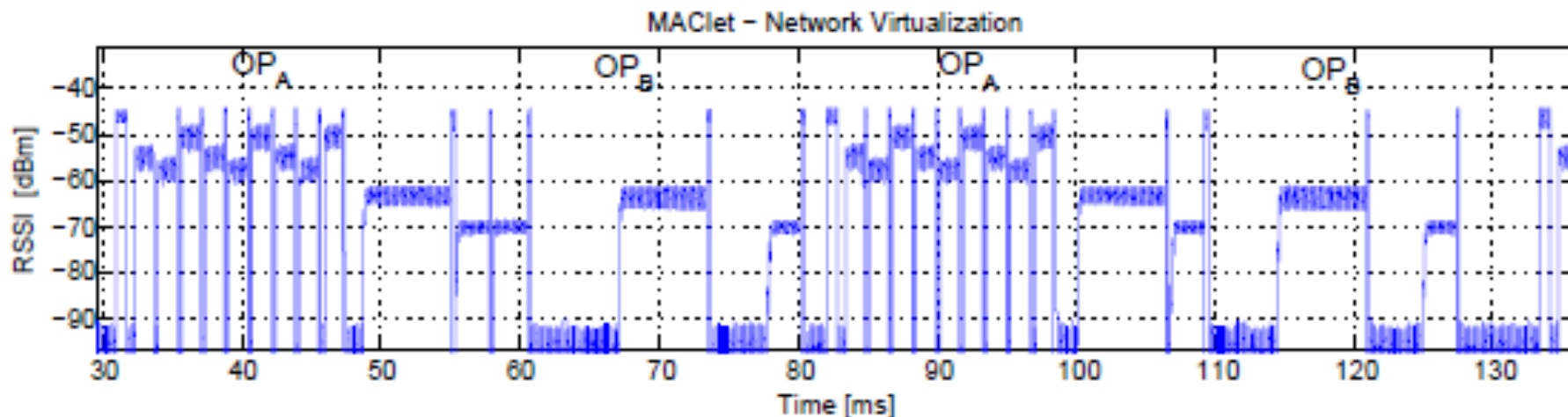
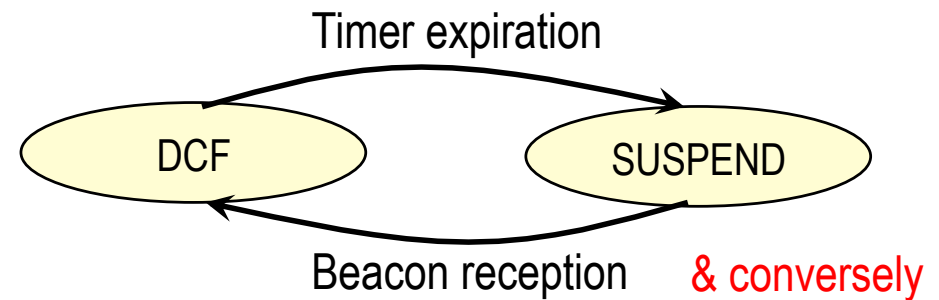
- ⇒ A: wants TDM, fixed rate
- ⇒ B: wants best effort DCF

→ Trivial with MAClets!

- ⇒ Customers of A/B download respective TDM/DCF MAClets!

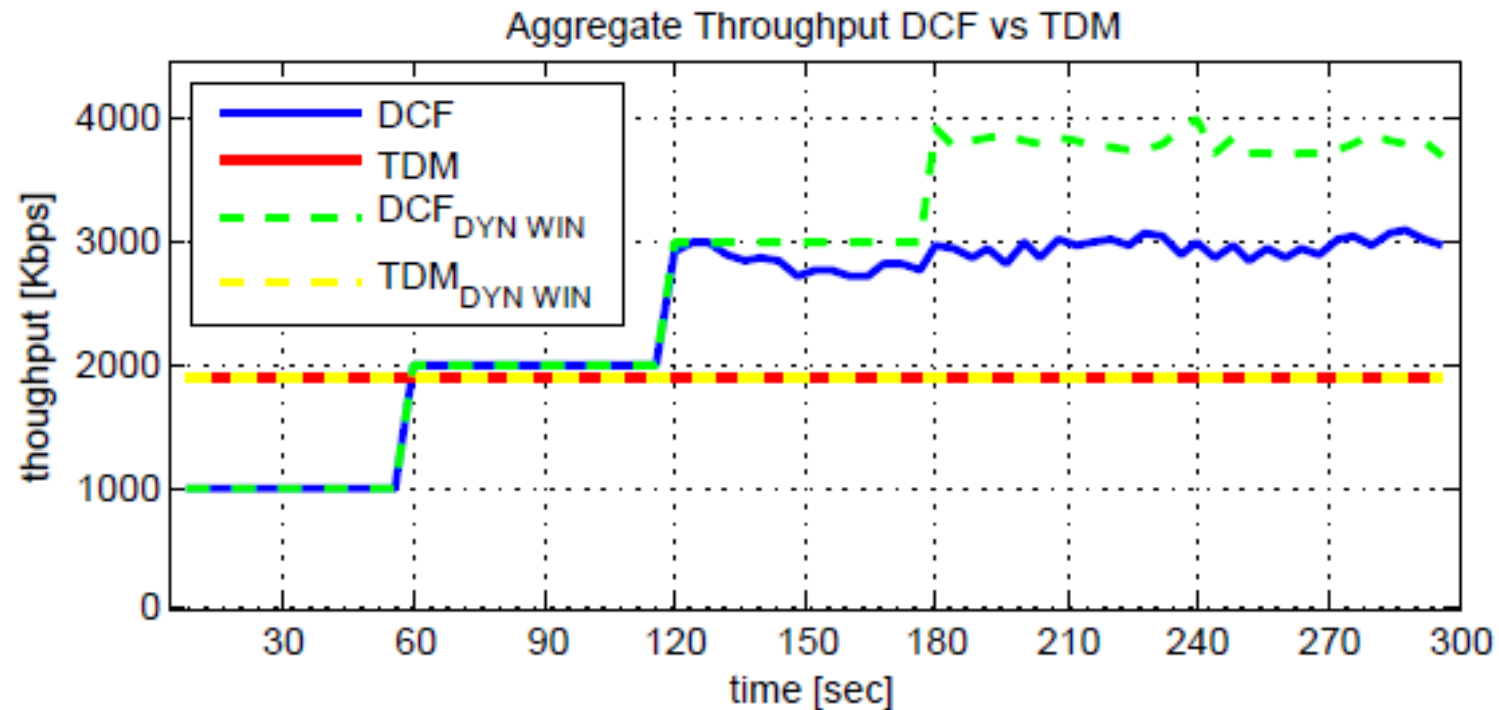
→ Isolation via MAClet design

- ⇒ Time slicing DESIGNED INTO the MAClets! (static or dynamic)



Throughput Performance

3 FIXED stations @ 0.63 Mbps vs. 5 BEST stations @ 1Mbps



Home Networks with MAClets

→ Heterogeneous applications at home

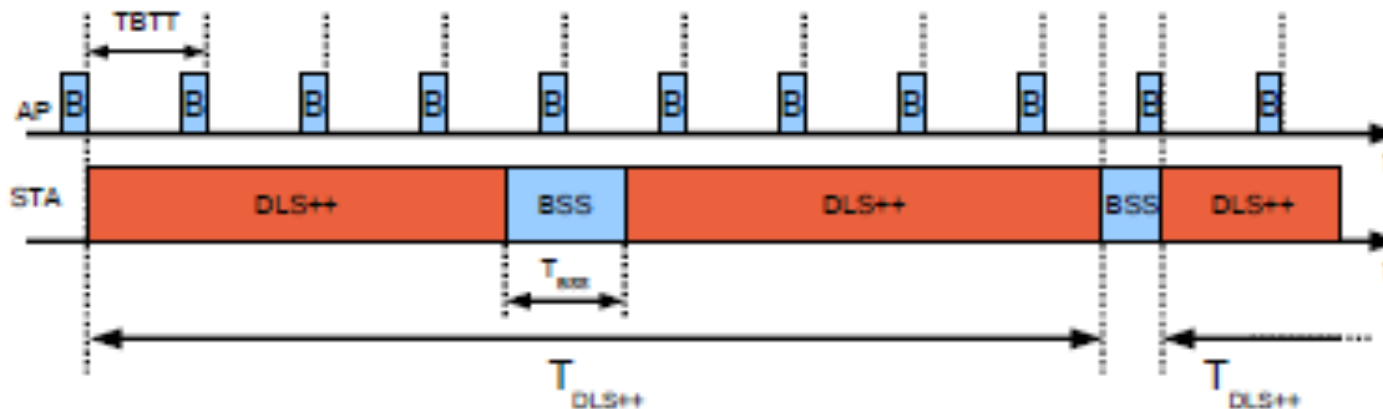
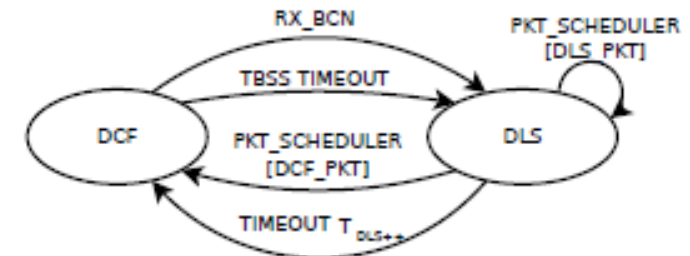
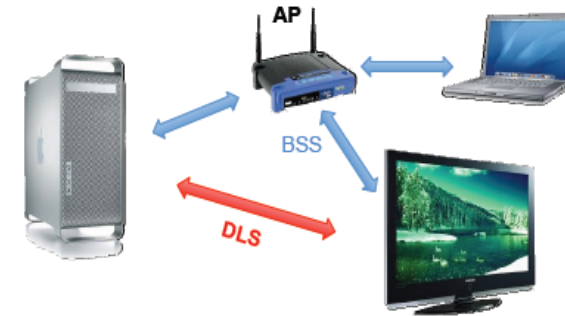
⇒ E.g. Video streaming and web browsing

→ Trivial with MAClets!

⇒ The Smart TV is not expected to implement any specific standard amendment

⇒ DLS protocol can be loaded when necessary

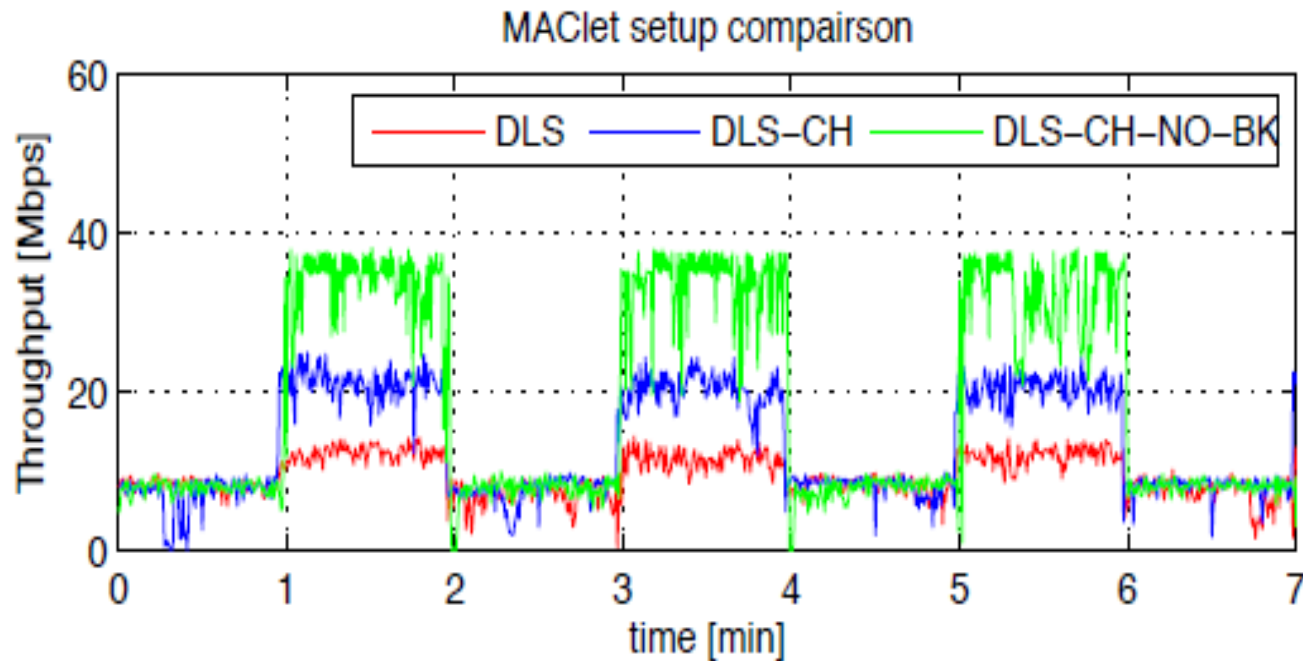
⇒ The network owner can push further optimizations:
→ additional channel for direct link channel, without losing association
→ Additional channel for direct link with greedy backoff



Throughput Performance

→ Experiment with a periodic switching from DLS++ to DCF

⇒ For testing multithreading and synchronization mechanisms



Conclusions

→ **New vision:**

- ⇒ MAC no more an all-size-fits-all protocol
- ⇒ Can be made context-dependent
- ⇒ Complex scenarios (e.g. virtualization) become trivial!

→ **Very simple and viable model**

- ⇒ Byte-coded XFSM injection
- ⇒ Does NOT require open source NICs!

→ **Next steps**

- ⇒ We focused on the «act» phase; what about the decision and cognitive plane using such new weapons?
- ⇒ can we think to networks which «self-program» themselves?
 - Not too far, as it just suffices to generate and inject a state machine...

Public-domain Platform

→ Supported by the FLAVIA EU FP7 project

⇒ <http://www.ict-flavia.eu/>

→ general coordinator: giuseppe.bianchi@uniroma2.it

→ Technical coordinator: ilenia.tinnirello@tti.unipa.it

→ Public domain release in alpha version

⇒ <https://github.com/ict-flavia/Wireless-MAC-Processor.git>

⇒ Developer team:

→ ilenia.tinnirello@tti.unipa.it

→ domenico.garlisi@dieet.unipa.it

→ fabrizio.giuliano@dieet.unipa.it

→ francesco.gringoli@ing.unibs.it

→ Released distribution:

⇒ Binary image for WMP

⇒ Source code for MAClet Manager

⇒ You DO NOT need it open source!

Remember the “hard-coded” device philosophy...

→ Conveniently mounted and run on Linksis or Alix

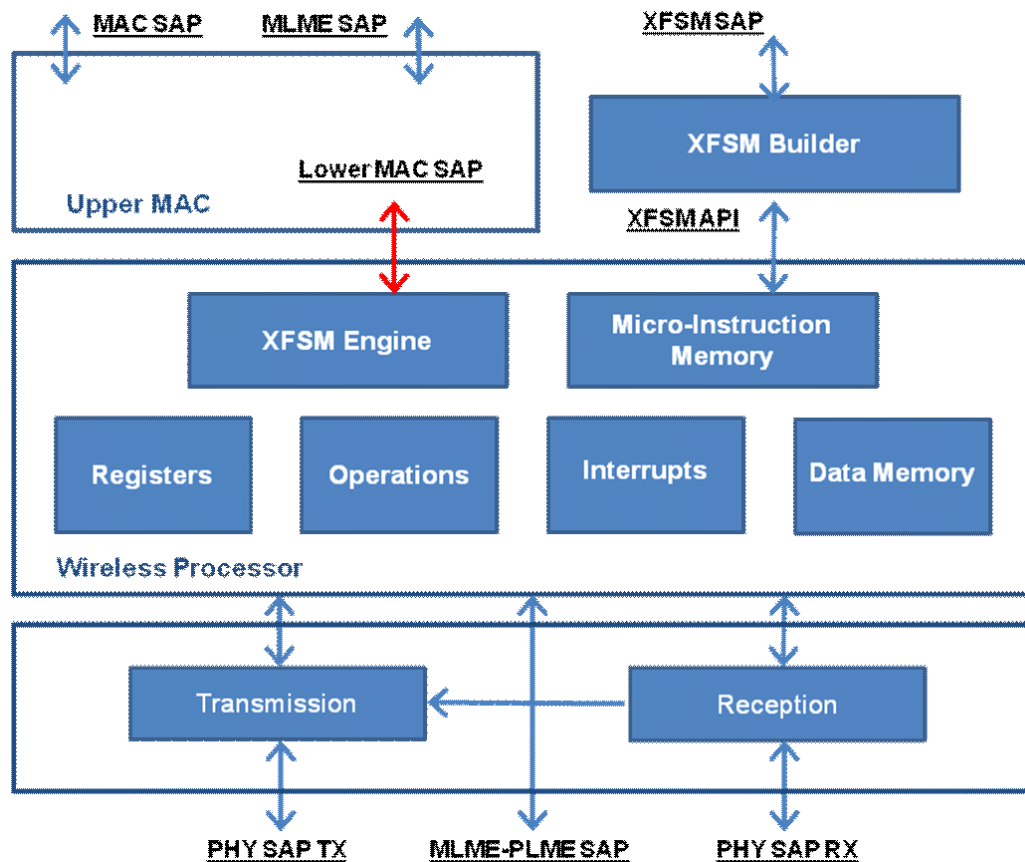
⇒ Source code for everything else

⇒ Manual & documentation, sample programs



WMP Overall architecture

from protocol-specific hard-coded device to protocol executor

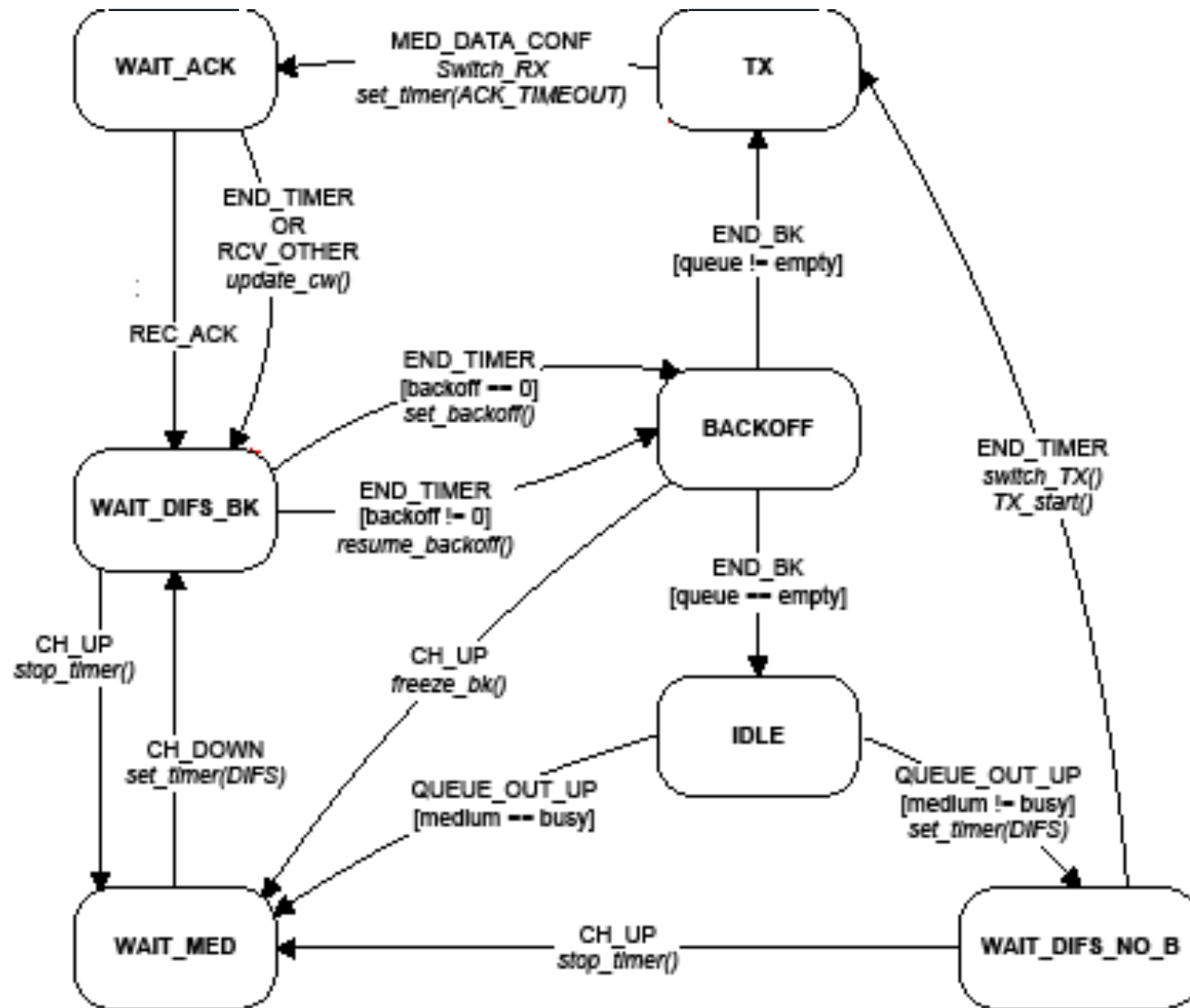


- **MAC Engine:** XFSM executor
- **Memory blocks:** data, prog
- **Registers:** save system state (*conditions*);
- **Interrupts block** passing HW signals to Engine (*events*);
- **Operations** invoked by the engine for driving the hardware (*actions*)

The MAC engine works as a Virtual MAC Machine

XFSM example: legacy DCF

simplified for graphical convenience



Actions:

set_timer, stop_timer,
set_backoff,
resume_backoff,
update_cw,
switch_TX, TX_start

Events:

END_TIMER,
QUEUE_OUT_UP,
CH_DOWN, CH_UP,
END_BK,
MED_DATA_CONF

Conditions:

medium, backoff,
queue