

# SureMail: Notification Overlay for Email Reliability

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HotNets 2005

# Motivation

- **Silent** email loss
  - email “vanishes” without sender/recipient knowledge
  - can be costly even if relatively rare
    - missed opportunities, misunderstanding, or worse
- **Nontrivial** problem
  - anecdotal evidence
  - measurement studies
    - 0.69% loss rate [Lang & Moors 2004]
    - 0.1-5% loss rate [Afergan & Beverly 2005]
  - commercial offerings to address the problem
    - e.g., Pivotal Veracity, Zenprise

HotNets e-ticket

*"We have sent it through again. If you do not receive it with in an hour or two, please let us know."*

Funding proposal

*"No I never got and I never acked it... My last mail from you was on XYZ."*

IMC 2005 decision notification

*"I recd reviews for one paper (#X) but not that of #Y."*

IMAP server upgrade problems

*"Some unanticipated migration problems occurred that may have caused some lost or delayed email."*

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# Silent Email Loss

- Why email loss?
  - spam filtering: big problem  $\Rightarrow$  aggressive filtering
    - MS: 90% of emails discarded before hitting user mailboxes
    - AOL: 100 emails per month to maintain IP white-listing
  - server failures and upgrades
    - SMTP is not end-to-end reliable
- (Non-)Delivery status notifications
  - compounds spam problem
  - raises privacy concerns
- So email loss is often silent

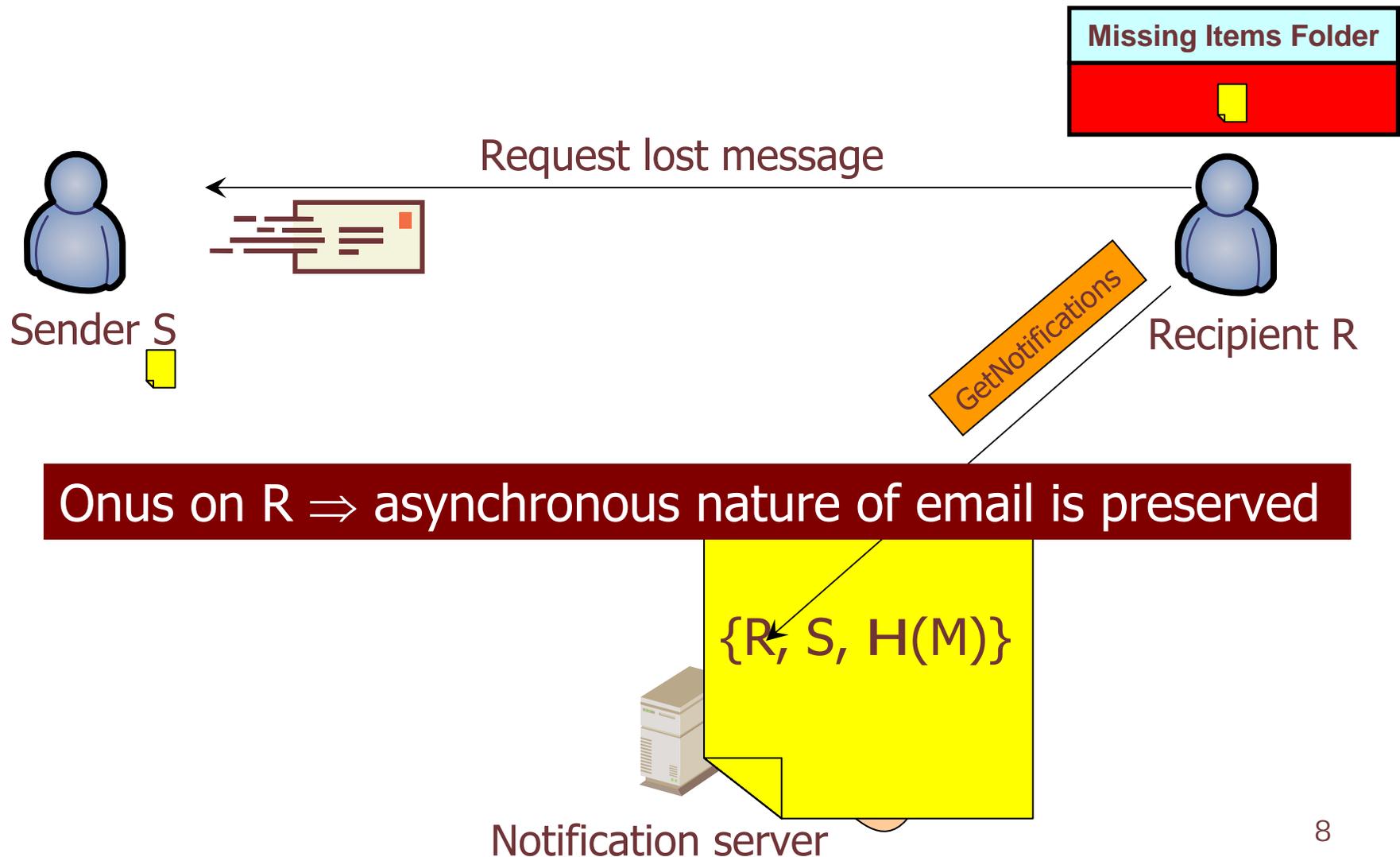
# Fixing the Problem

- Improve the email delivery infrastructure
  - more reliable servers
    - e.g., cluster-based (Porcupine [Saito '00])
  - server-less systems
    - e.g., DHT-based (POST [Mislove '03])
  - total switchover might be risky
- “Smarter” spam filtering
  - moving target  $\Rightarrow$  mistakes inevitable
  - non-content-based filtering still needed to cope with spam load

# SureMail

- Address the problem from the outside
  - add separate **notification** overlay
  - emails & email delivery infrastructure left untouched
  - eases deployment, bounds the worst case
- Design requirements:
  - minimize demands on infrastructure and users
  - preserve asynchronous operation and privacy
  - maintain defenses against spam and viruses
  - minimize overhead

# Basic Operation



# Notification Overlay

- Decentralized
  - limited collusion among the constituent nodes
- Efficient notification server lookup
  - e.g.,  $R \rightarrow H(R)$  in a DHT setup
- Agnostic to actual implementation
  - end-host-based (e.g., always-on user desktops)
  - infrastructure-based (e.g., “NX servers”)

# Challenges

- ❑ Privacy
  - information on users' email habits or even just whether an email address is active could be leaked
- ❑ Notification spam
  - spammers could spoof notifications and burden users
  - "annoyance attacks" discredit notifications in general
- ❑ Even the notification infrastructure isn't trusted
- ❑ No universal PKI for email users

# SureMail Goals

- Protect the recipient's identity
  - attacker shouldn't be able to learn R's identity or monitor the volume of notifications intended for R
- Protect the sender's identity
  - attacker shouldn't be able to learn S's identity or monitor the volume of notifications posted by S
- Block notification spam
  - attacker shouldn't be able to spoof notifications

$\{R, S, H(M)\}$

# Assumptions

- ❑ No email eavesdroppers
  - bigger problems otherwise
- ❑ Limited collusion among notification nodes
  - needed only to avoid leaking information on whether or how many notifications R is receiving

# Key Mechanisms

#1: Email-based handshake

#2: Decoupled registration and notification

#3: Email-based shared secret

#4: Reply-based shared secret

$\{H(R), S, H(M)\}$

## #3: Email-based shared secret

**Goal:** prevent snooping on sender identity



Email  $M_{old}$  from S to R is known only to S and R

- $H(M_{old})$  could serve as implicit identifier of S to R
- But it doesn't quite serve as authenticator for S:
  - $D_{not}$  knows  $H(M_{old})$ , so it could spoof notifications from S
  - even other attackers could do so by first sending  $M_{spoofed}$  purporting to be from S

$\{H(R), H(M_{old}), H(M)\}$

# #4: Reply-based shared secret

**Goal:** block spoofing of notifications



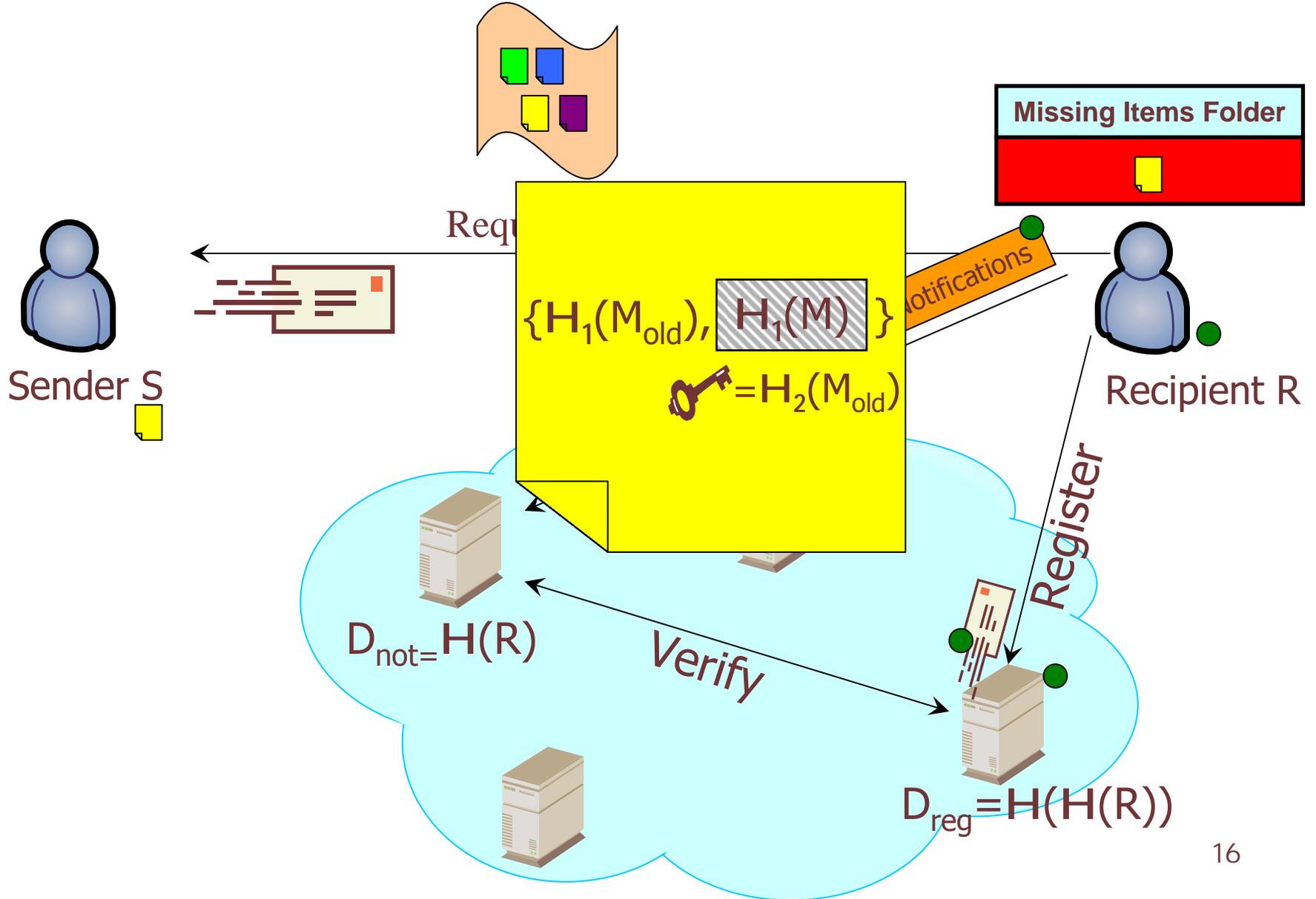
Users rarely have conversations with spammers

- R remembers (hashes of) recent emails from S that it has replied to
- If S receives a reply to  $M_{old}$  it had sent R,  $M_{old}$  can serve as a shared secret between S and R
  - $H_1(M_{old})$  as implicit identifier,  $H_2(M_{old})$  as authenticator
- Hard for a spammer (even  $D_{not}$ ) to spoof

$$\{H(R), H_1(M_{old}), \boxed{H_1(M)}\}$$

 =  $H_2(M_{old})$

# Putting it all together



# Other issues

- ❑ Reply-detection:
  - “in-reply-to” insufficient, indirect checks needed
- ❑ Reducing overhead:
  - look for implicit ACK (reply) or NACK (bounce-back)
  - post notifications selectively (for “important” emails)
- ❑ First-time “legitimate” senders:
  - indistinguishable from spammers
- ❑ Mobility:
  - reply-based scheme naturally lends itself to migration

# Status

- ❑ Ongoing measurement experiment
- ❑ Design being refined
- ❑ Implementation in the works

# Discussion

#1: Should the notification system be folded into the email infrastructure?

- Separation is advantageous:
  - provides failure independence
  - keeps the notification layer simple
    - small, fixed format notifications don't require the same kind of processing as virus-laden email
    - no direct benefit for spammers
  - provides engineering convenience
    - fewer dependencies, easier deployment

# Discussion

#2: Is there a social benefit to silent email loss because of the plausible deniability it provides?

- Any such benefit is far outweighed by the costs
  - Should cars be slightly unreliable because of the excuse it would give people when they miss an engagement?
- It is the asynchronous nature of email that is key and SureMail preserves that