



Achieving Bounded Fairness for Multicast and TCP Traffic in the Internet

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Outline

- Problem Statement
- Design Objectives
 - ◆ Essential Fairness
- Random Listening Algorithm (RLA)
 - ◆ Essentially Fair to TCP
 - ◆ Multicast Fairness
 - ◆ Performance

Problem Statement

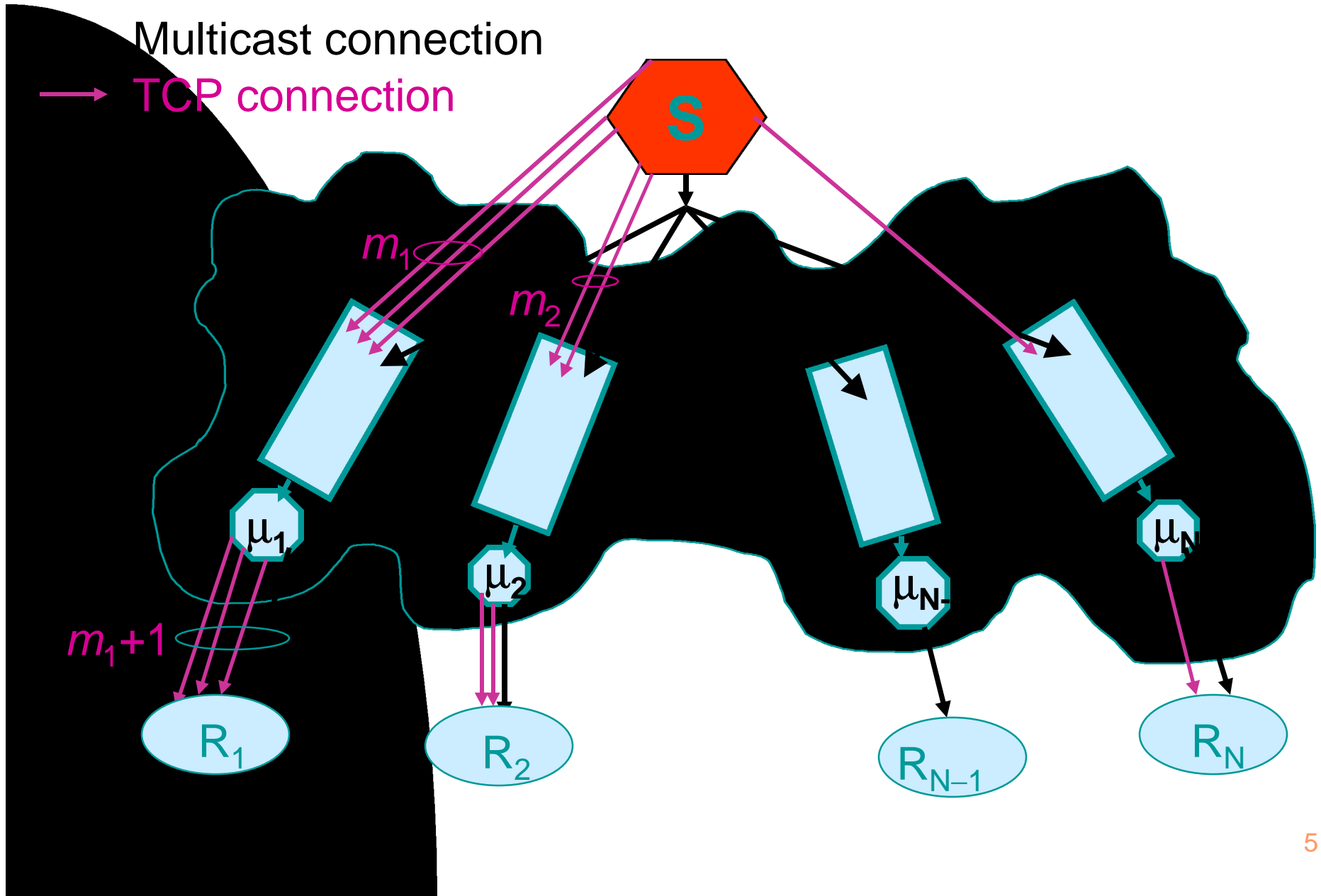
- Based on current Internet infrastructure.
 - ◆ Drop-tail gateways prevail.
 - ◆ RED gateways preferred.
- Based on end-to-end packet loss information.
- TCP background traffic.

Different receivers lose packets at different times.

Design Objective – TCP-like

- Responsive to congestion.
- Probe available bandwidth.
- Multicast fairness.
- Fair to TCP.

Restricted Topology



Concepts

- **Soft bottleneck:**

- ◆ $k = \arg \min_i \{m_i / (m_i + 1)\}, i = 1, \dots, N.$

- **One share of bandwidth:**

- ◆ $I_{min} = m_k / (m_k + 1).$

- **Absolute fairness** (equal share): $I_{RLA} = I_{min}.$

- **Essential fairness** (bounded share):

- ◆ $0 < I_{RLA} < m_k, a I_{min} \leq I_{RLA} \leq b I_{min}, a \leq b \leq N.$

Random Listening Algorithm

Based on TCP SACK.

- **Loss detection:** selective ACK
- **Congestion detection:**
 - ◆ losses within 2 RTT are grouped.
- **Window adjustment policy:** random listening.
 - ◆ Upon a congestion signal : $W \leftarrow W/2$ w.p. $1/n$.
 - ✦ $n = \#$ of rcvrs reporting congestion frequently.
 - ◆ Once a pkt is ACKed by all rcvrs, $W \leftarrow W+1/W$.

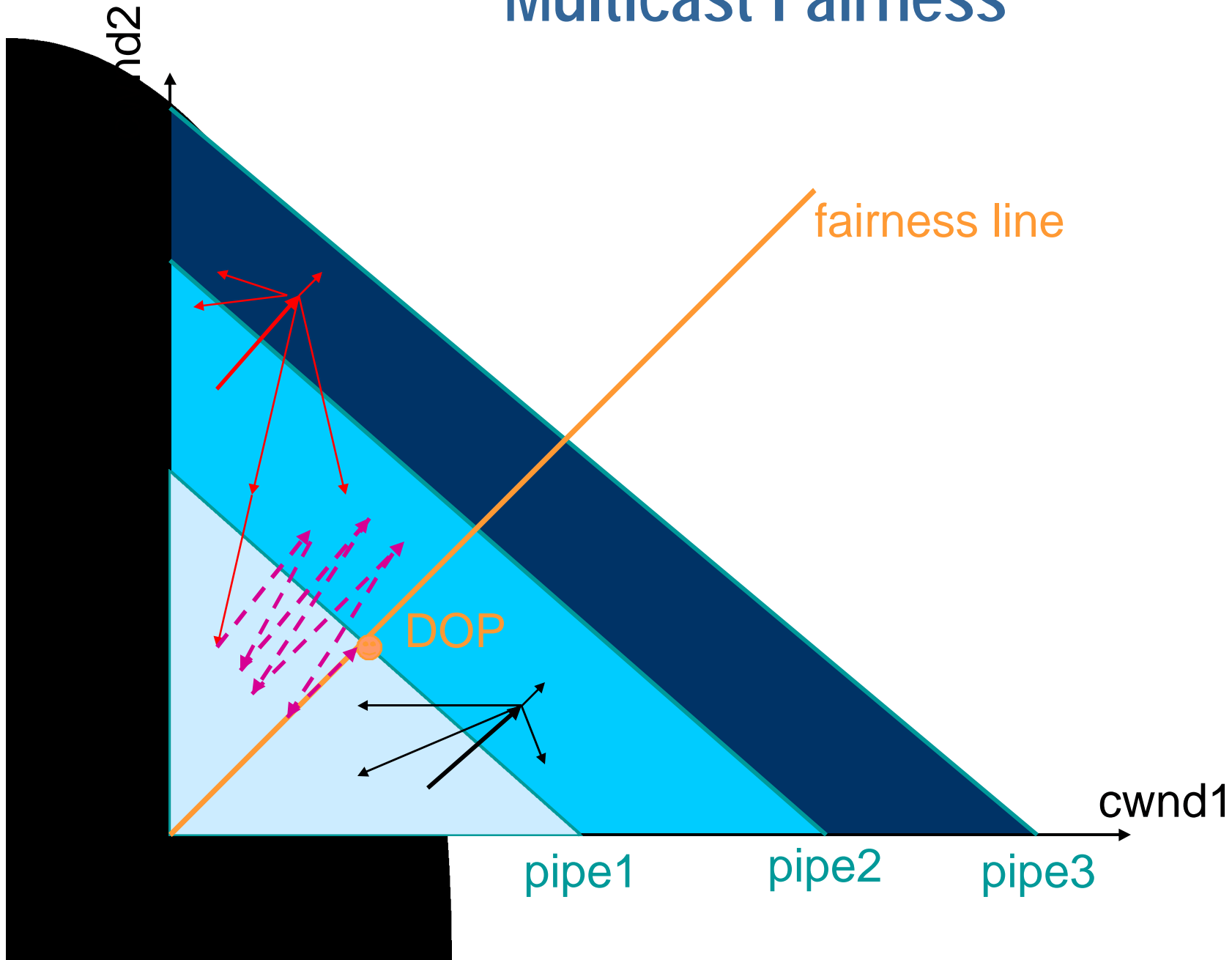
Properties of RLA

- Share with TCP through drop-tail:
 - ◆ congestion frequency.
- Share with TCP through RED:
 - ◆ same loss probability.
- Essentially fair to TCP:
 - ◆ RED: $1/3 I_{TCP} \leq I_{RLA} \leq \sqrt{3n} I_{TCP}$.
 - ◆ drop-tail: $1/4 I_{TCP} \leq I_{RLA} \leq 2n I_{TCP}$.

Properties of RLA (cont)

- Multicast fairness.
- Simple, similar to TCP.
- No need to locate soft-bottleneck.

Multicast Fairness

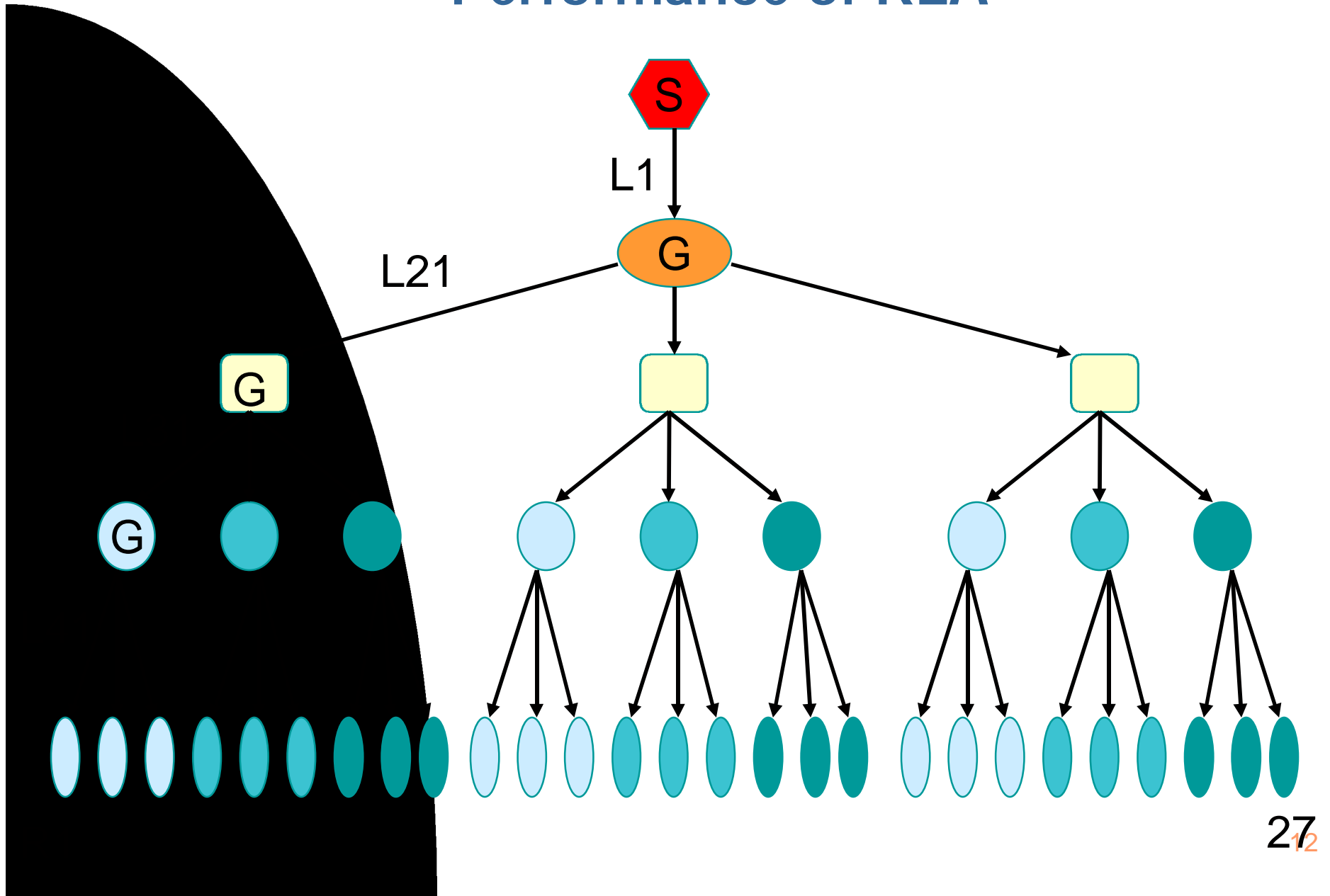


Multicast Fairness

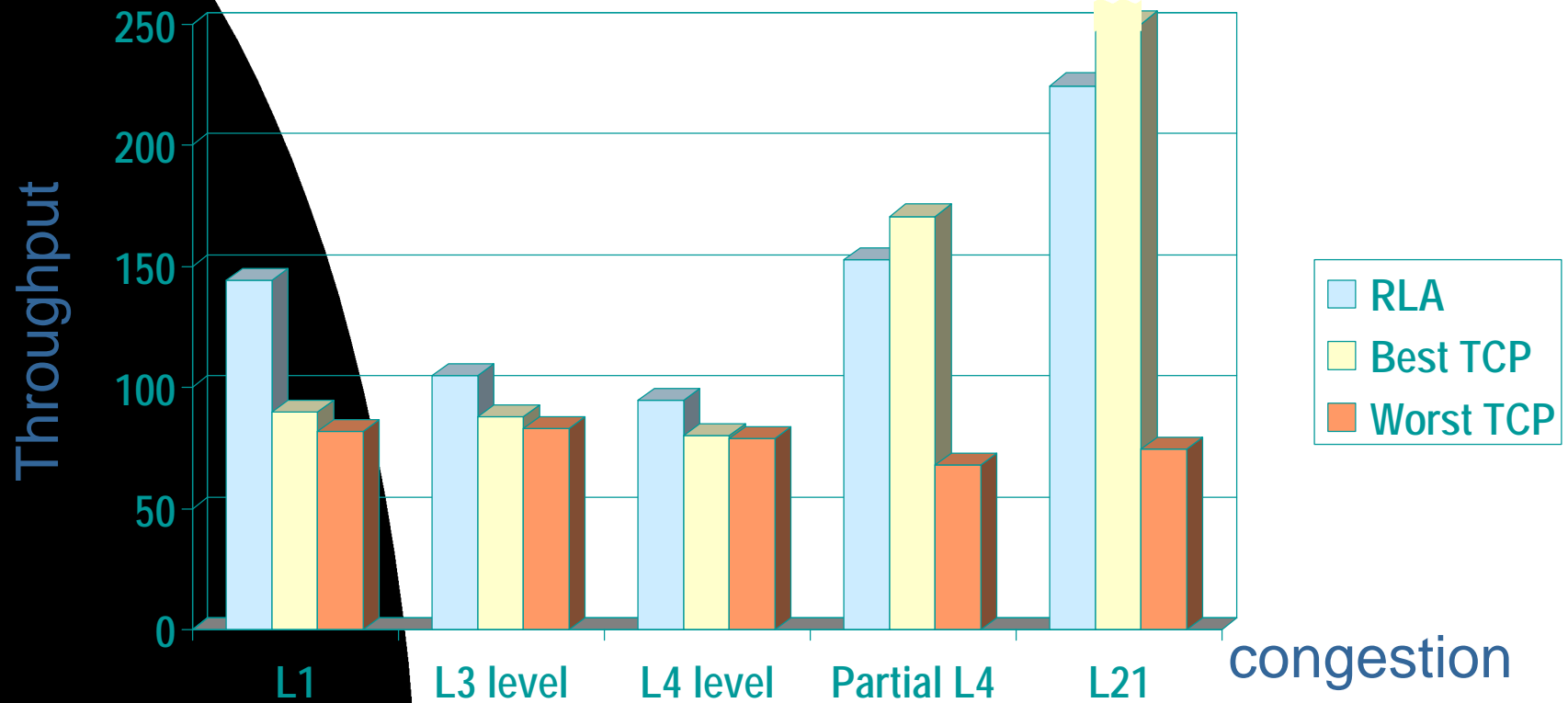
Two dimensional Markovian process:

- desired operating point (DOP) is recurrent.
- *avg cwnd*'s of the two senders are the same.
- average drifts point towards DOP.
- probability mass focuses on area around DOP.

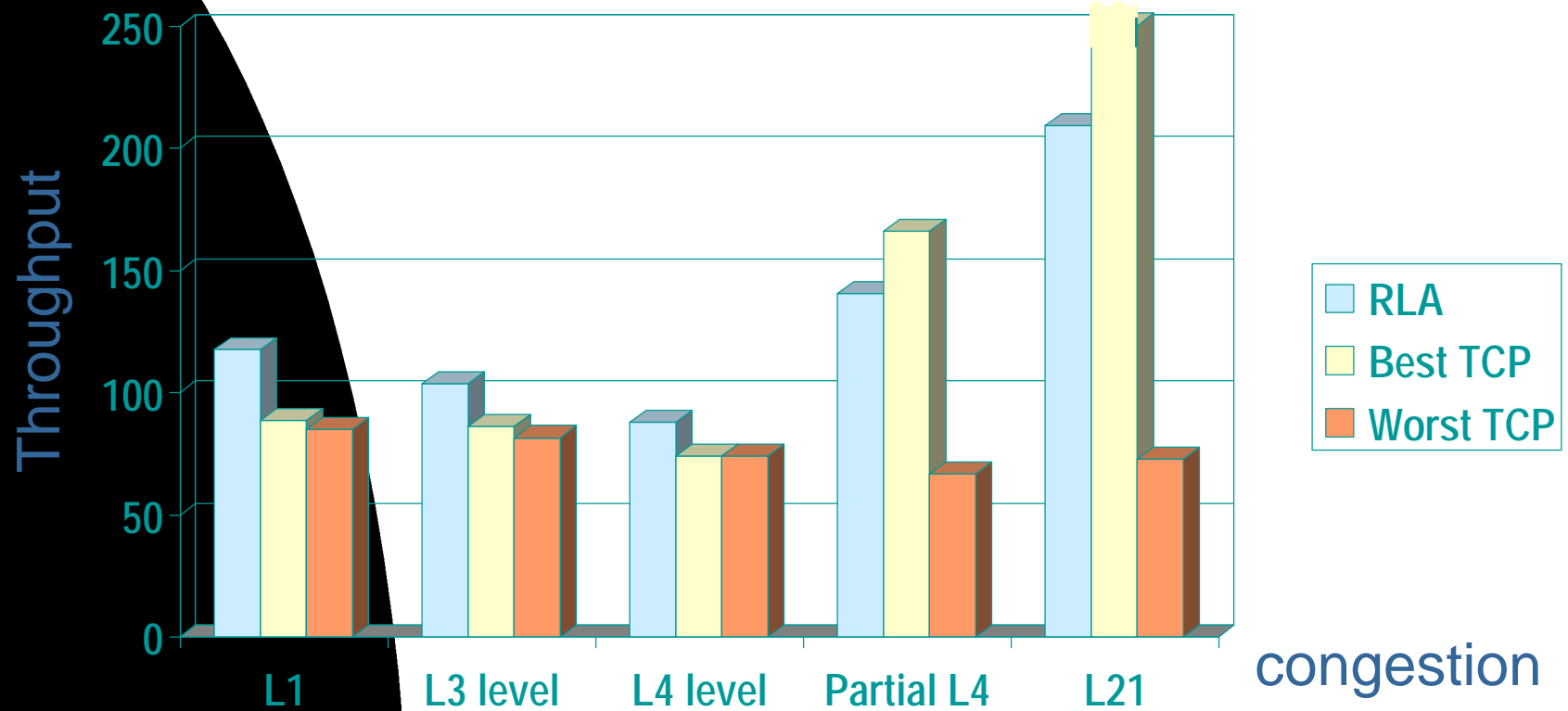
Performance of RLA



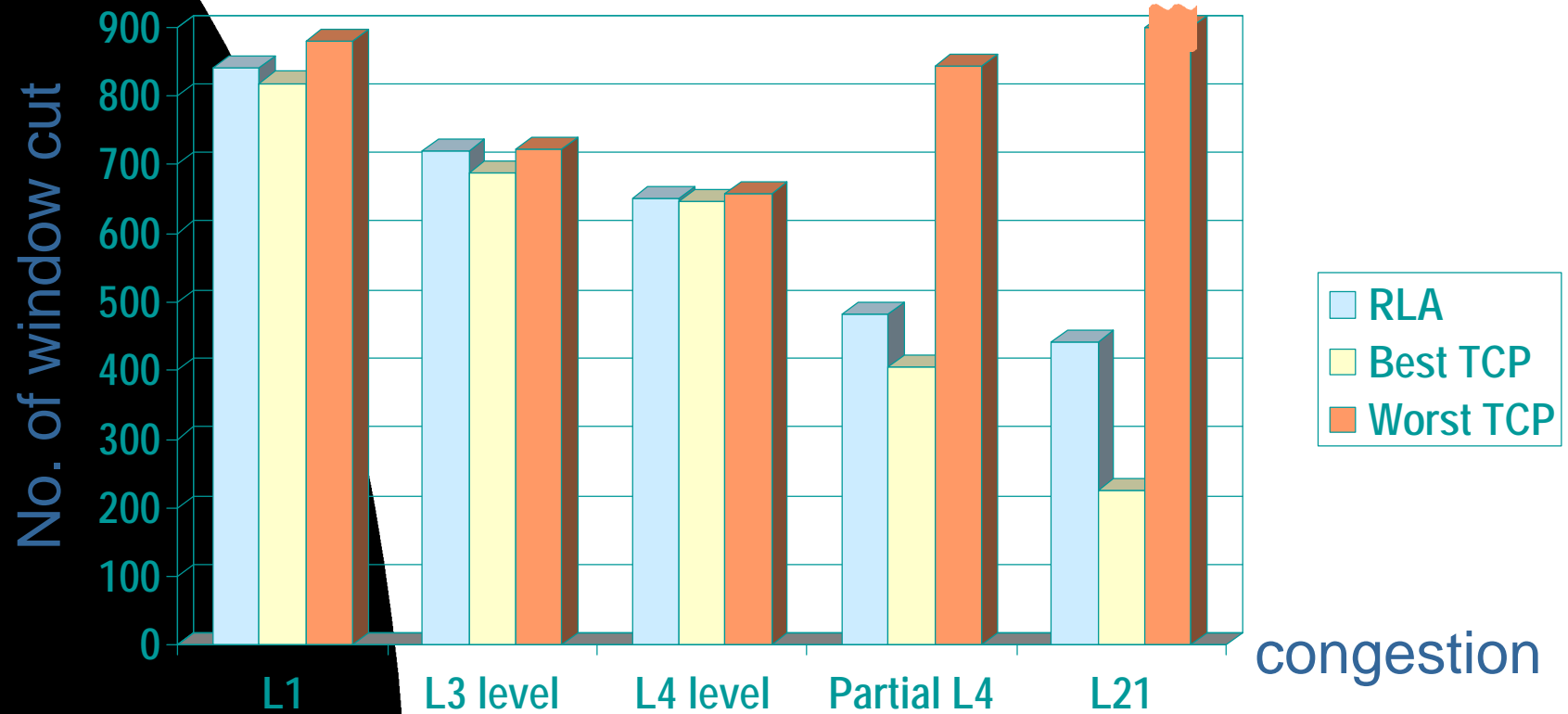
Result: Drop-tail Gateways



Results: RED Gateways

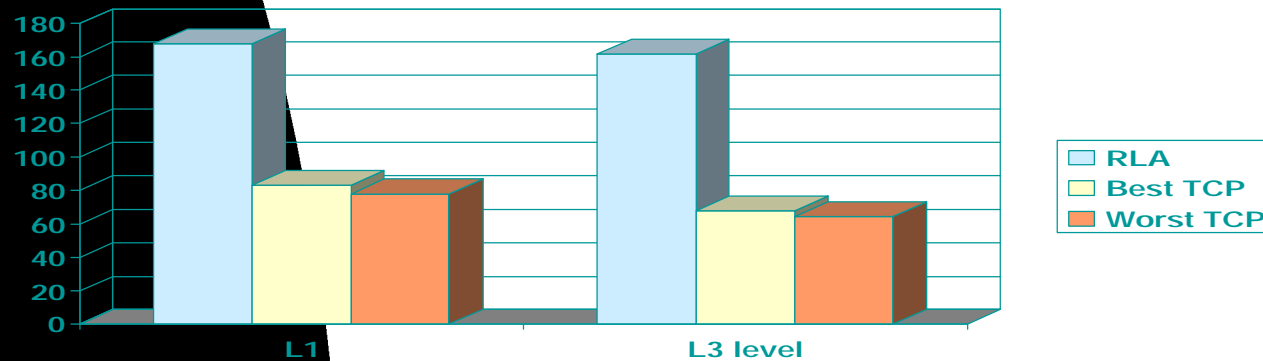


Results: Drop-tail Gateways



Results: Different RTTs

- Generalized RLA:
 - ◆ random listening threshold set to $(rtt_i / rtt_{max})^2 / n$.
- Preliminary Performance:



Summary

- Desirable Performance:
 - ◆ homogenous receivers: $I_{RLA} < 4 I_{TCP}$ for any n .
 - ◆ unbalanced congestion: $I_{RLA} \leq O(n) I_{TCP}; (DT)$
 - ◆ reasonably fair to TCP as shown by simulation.
- “**Random listening**” idea applies to other forms of control as well.