Understanding Block-level Address Usage in the Visible Internet

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The Discovery of Halley's Comet



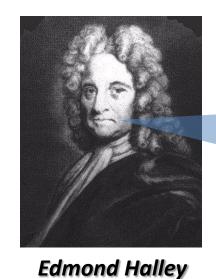






The Discovery of Halley's Comet

2 historical records (year 1531, 1607)1 observation (year 1682)



"It's the same object which returns to earth every 76 years."

3 simple observations



an astronomer

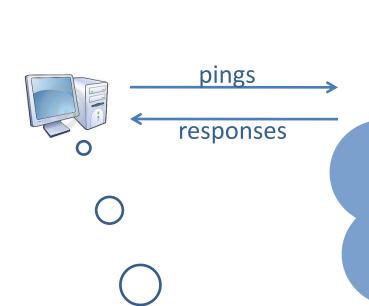


1 simple characteristic of the comet

SIMPLE observations inferred **SIMPLE** conclusion can have **TREMENDOUS** value.







Internet

Address Utilization?

Dynamic Addressing?

Our Q: what can simple observations about the Internet say?



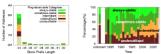


Methodology



- Active probing, pattern analysis, clustering, classification

Application



- Network management, resource allocation, Internet trend study

Validation



- USC's network, the general Internet, consistency across time





Methodology

Spatial Correlation?

Group addresses into blocks by usage

Address Utilization?

Find blocks with less than 10% time responsive

Dynamic Addressing?

Blocks switching state (up/down) frequently

Low-bitrate Identification?

Utilize standard deviation of RTTs

Application ____

More frequent probing?
Block sizes?
Block-level usage?

Resource reallocation? Efficient management?

Botnet detection?
Spam filtering?
Click fraud?

Auto content serving? Network management?

Validation

USC's network, General Internet, Consistency

USC's network, General Internet, Consistency

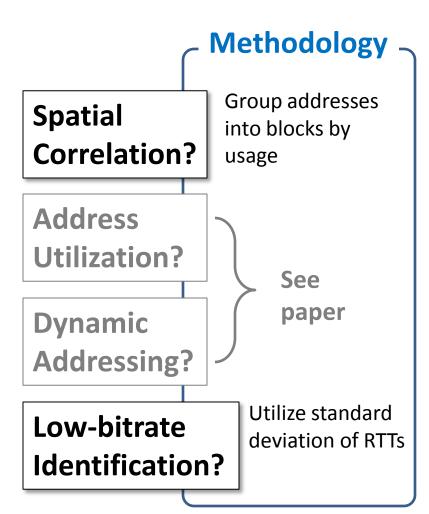
USC's network, General Internet, Consistency

USC's network, General Internet









Application -

More frequent probing?
Block sizes?
Block-level usage?

See paper

Auto content serving? Network management?

Validation

USC's network, General Internet, Consistency

> See paper

USC's network, General Internet







Related Work

• J. Heidemann, Y. Pradkin, R. Govindan, C. Papadopoulos, G. Bartlett, and J. Bannister. <u>Census and Survey of the Visible Internet.</u> In *Proceedings of the ACM Internet Measurement Conference (IMC)*, p. 169-182. Vouliagmeni, Greece, October, 2008.

What's the same?

- Collection methodology (and datasets)
- Error bounds on ping census accuracy: undercounts by about 40%
- Preliminary metrics

What's new? deeper understanding; new interpretation

- new metrics
 - block-level analysis, not just addresses
 - RTT, not just responsivness

new algorithms

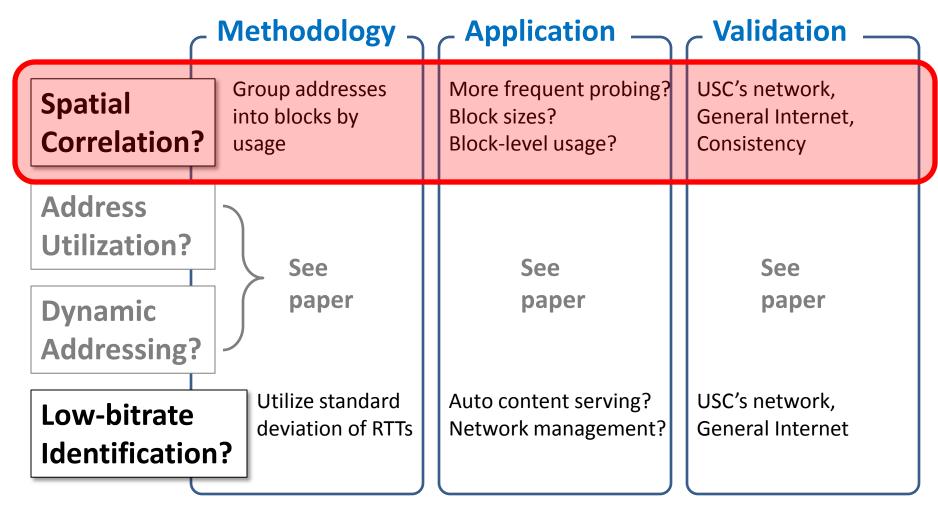
- block identification
- low-bitrate identification

new conclusions

- evaluation of block utilization
- trends of address utilization
- trends of dynamic addressing











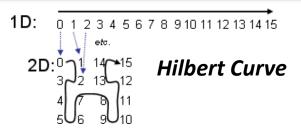


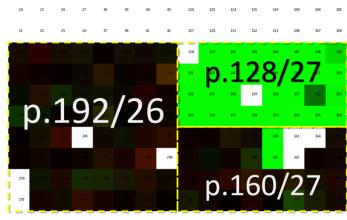
Background: What space?

- IPv4 address space
 - address block: p/n: addresses with common n-bit prefix p
 - a.b.c.d and a.b.c.(d+1) are adjacent addresses

A /24 block (*p*/24) with 256 addresses,

Layout *Hilbert Curve* keeps adjacent addresses physically near each other.







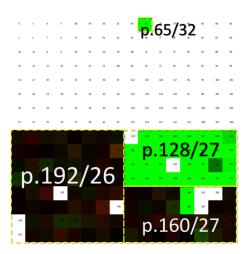


Hypothesis: Spatial Correlation

- What is Spatial Correlation?
 - adjacent addresses are likely to be used in the same way
 - ⇒ spatial correlation of address blocks
 - ⇒ usage blocks
- Usage blocks
 - are NOT allocated blocks, but correlated
 - Internet addresses are allocated in blocks (ICANN to regional registries to ISPs to you)
 - addresses in one block are usually assigned to similar users
 - are what we want to observe if exist
 - observable blocks → usage blocks







Spatial Correlation: Application

- Why care?
 - Efficiently select representative addresses to conduct more detailed study
 - Addresses in one block are used in the same way
 - So only need few representatives to probe in the future





Spatial Correlation: Methodology

Data Collection Representation **Block Identification**

Input: data for individual *addresses*

Output: address sharing similar usage grouped into *observable blocks*

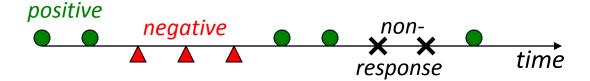




Spatial Correlation: Data Collection

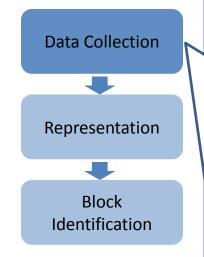
How Ping each address in random /24 blocks every 11 minutes for a week and collect the probe responses.

1% of the allocated IPv4 address space probed.



Why Systematic pings reveal more information.

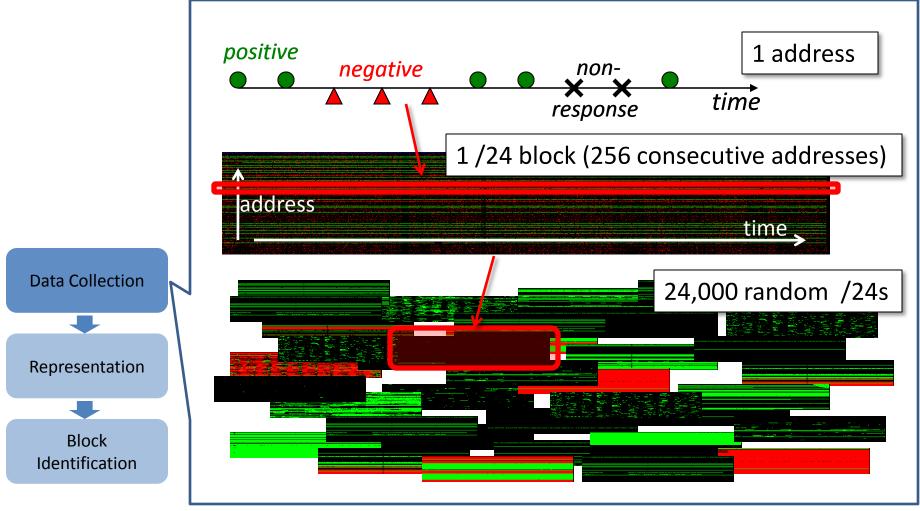
Validity of ping: IMC'08 paper established error bounds: not perfect, but often pretty good; ~40% undercount







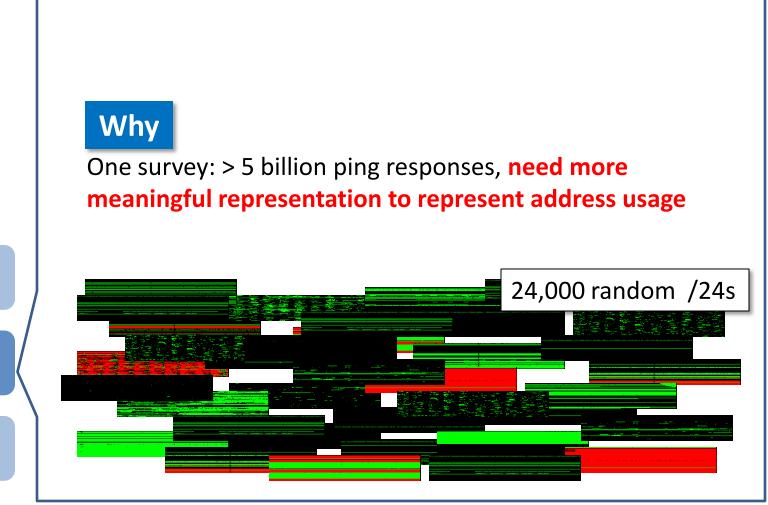
Spatial Correlation: Data Collection







Spatial Correlation: Representation





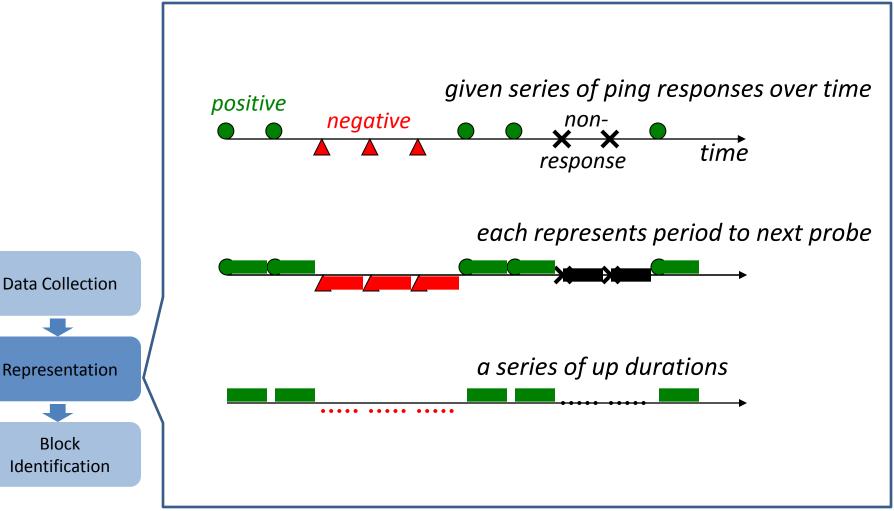
Data Collection

Representation

Block Identification



Spatial Correlation: Representation





Block



Spatial Correlation: Representation



1st duration length: 2

2nd duration length: 2

3rd duration length: 1

How 3 metrics to capture address usage



Data Collection

Block Identification

Availability (A)

:= normalized *sum* of up durations

Example:

$$= (2+2+1) / 10 = 0.5$$

Intuition:

utilization efficiency

Volatility (V)

:= normalized # of up durations

Example:

$$= 3 / (10/2) = 0.6$$

Intuition:

high V infers dynamics

Median-Up (U)

:= median up duration

Example:

= median(2,2,1) = 2

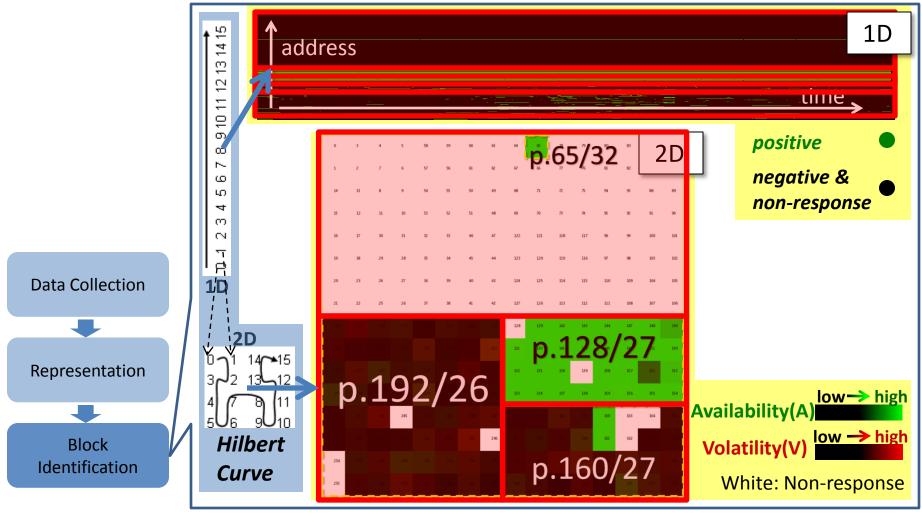
Intuition:

typical duration





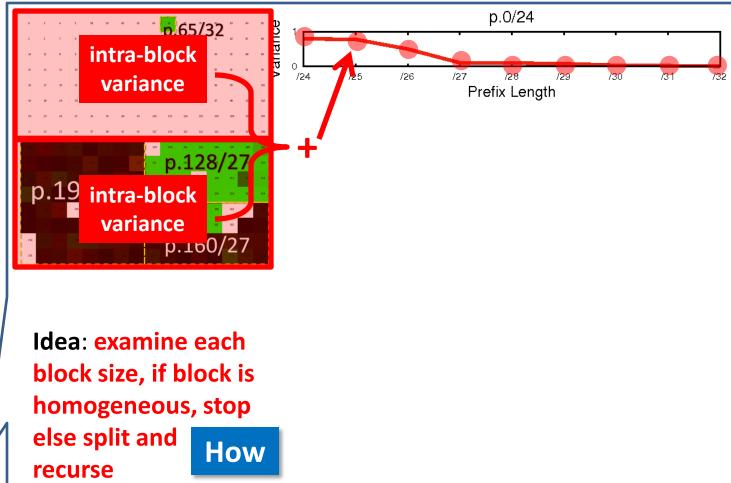
Spatial Correlation: Block Identification

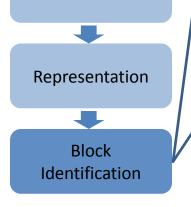






Spatial Correlation: Block Identification



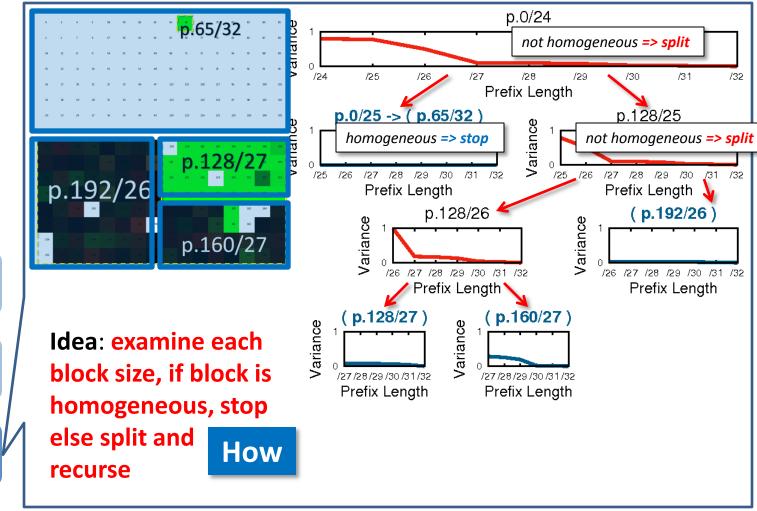


Data Collection





Spatial Correlation: Block Identification





Data Collection

Representation

Block

Identification



Spatial Correlation: Validation

- Validation is hard
 - Where to find ground truth?
 - decentralized management
 - usage block ground truth?
- Use three complementary ways:
 - Compare to USC's network (operator provided truth)
 - Compare to general Internet (hostname inferred truth)
 - Evaluate different samples and dates
 - is 1% of the Internet enough? yes!
 - trends change some over time
 - details: paper section 5.3





Spatial Correlation: USC's Network

• Why

- quite solid truth (operator provided)
- knowledge of both allocated blocks and usage
 blocks

- How
 - compare observable blocks (result to validate)
 with usage blocks (ground truth)





Spatial Correlation: USC's Network

category:		blocks		percentage		_
ground truth usage blocks			243	100%		_
	lse negative not in use	mostly	$\frac{105}{19}$	43%	approac	ch is
talse-neg	not responding	non-use	28		incompl	lete
blocks we	fow responding	(23%)	19			
missed to	single-block multi-us /25 to /27	sage	46			
identify	/25 to /27	sometimes	9			
	/28 to /32	error (20%)	37			
	ocks identified		147		100%	but what is
	correctly identified		138	57%	94%	found is correct
false-pos.:	false positive		9		6.1%	Touriu is correct
blocks we	multi-block single	-usage	9			
wrongly identified	very accurate v	when it re	ach	es a c	onclus	ion





Spatial Correlation: General Internet

- Why
 - unbiased truth (randomly selected)
- How
 - Infer usage blocks from hostnames
 - dhcp-host-xxx.example.net
 - compare observable blocks (result to validate)
 with usage blocks (ground truth)





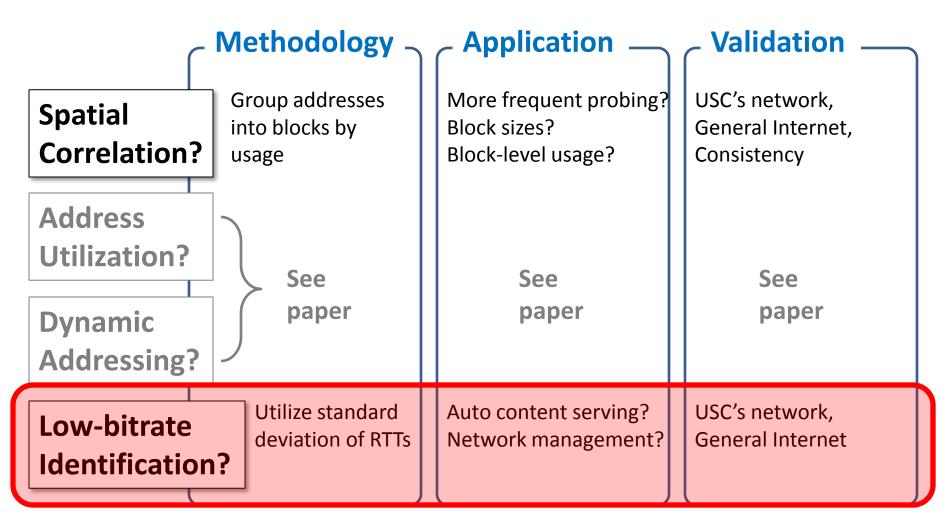
Spatial Correlation: General Internet

category:	blocks	percentage	
/24 randomly selected	100	100%	
decided (/24 inferred from hostname)	37	37% 100%	
correct	25	68%	
wrong (false negative)	12	32%	
few responding	6	mostly correct	
single-block multi-usage	6	(and more	
undecided	63	63% than USC)	
no hostname	45	ground truth is	
few hostnames	7	hard to infer	
potential $/24$ inferred	7	nard to line	
correct	7		
has sub-/24 groupings	4		

methodology more complete when evaluate with unbiased sample











Background: What is low-bitrate?

- Addresses are connected to Internet through edge access links
- Different access link type has different bitrate
 - Dial-up: 56Kb/s
 - ADSL (typical): 3,000/768 kbit/s
 - *GPRS*: 57.6 Kb/s
 - *UMTS 3G*: 384 kbit/s

 We define low-bitrate as less than 100Kb/s, such as dial-up and GPRS.





Low-bitrate: Application

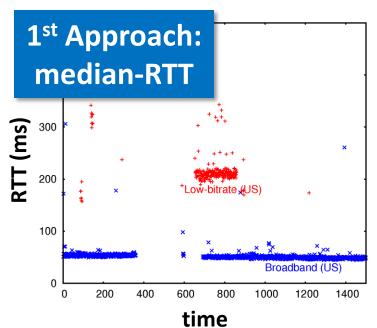
- Why care?
 - For the researchers
 - help understand trends in technology deployment
 - For the business
 - automatically match content and layout
 - For network management
 - low-bitrates links are correlated with short connecttimes and sparse usage.



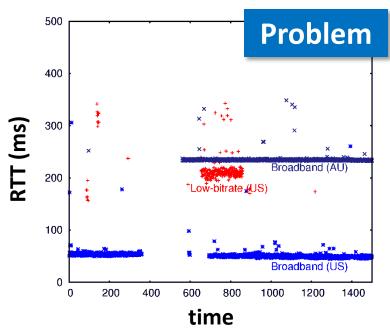


Methodology: Formalizing RTT -> Edge Bitrate

RTT = transfer + queuing + propagation



transfer distinguishes low-bitrate vs. broadband



but internationally *propagation* time dominates





Methodology: Formalizing RTT -> Edge Bitrate

RTT = transfer + queuing + propagation

edge-bitrate dependent, distance dependent, and varying but **consistent** 500 400 400 Low-bitrate (AU RTT (ms) Broadband (AU) Broadband (AU) 200 200 Low-bitrate (US) variance predicts 100 100 low-bitrate Broadband (US) Broadband (US) 80 200 400 600 1000 1200 1400 100 **Solution** time CDF of RTTs (%) (or consistency predicts broadband)





Low-bitrate: Validation

category:	blocks	percentage
hostname-inferrable edges	36	100%
low-bitrate blocks (6 dial, 2 mobile)	8	
$R_{\mu_{1/2},\sigma}^*(b) > \delta$ (true positive)	8	22%
$R_{\mu_{1/2},\sigma}^{*}(b) \leq \delta$ (false negative)	0	0%
broadband (21 dsl, 4 cable, 3 3G)	28	
$R_{\mu_{1/2},\sigma}^*(b) > \delta$ (false positive)	0	0%
$R^*_{\mu_{1/2},\sigma}(b) \leq \delta$ (true negative)	28	78%
clear hostname	25	what is found
confusing hostname	3	is all correct

can accurately find low-bitrate links





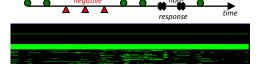




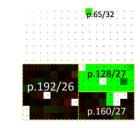


Conclusion

SIMPLE observations (*pings*)



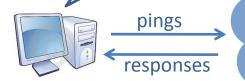
can tell ...



VALUABLE truths about the Internet.

spatial correlation, address utilization dynamic addressing, low-bitrate

Visit www.isi.edu/ant
for our dataset and more information!



Internet





