

Inter-Datacenter Bulk Transfers with NetStitcher

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joint work with:

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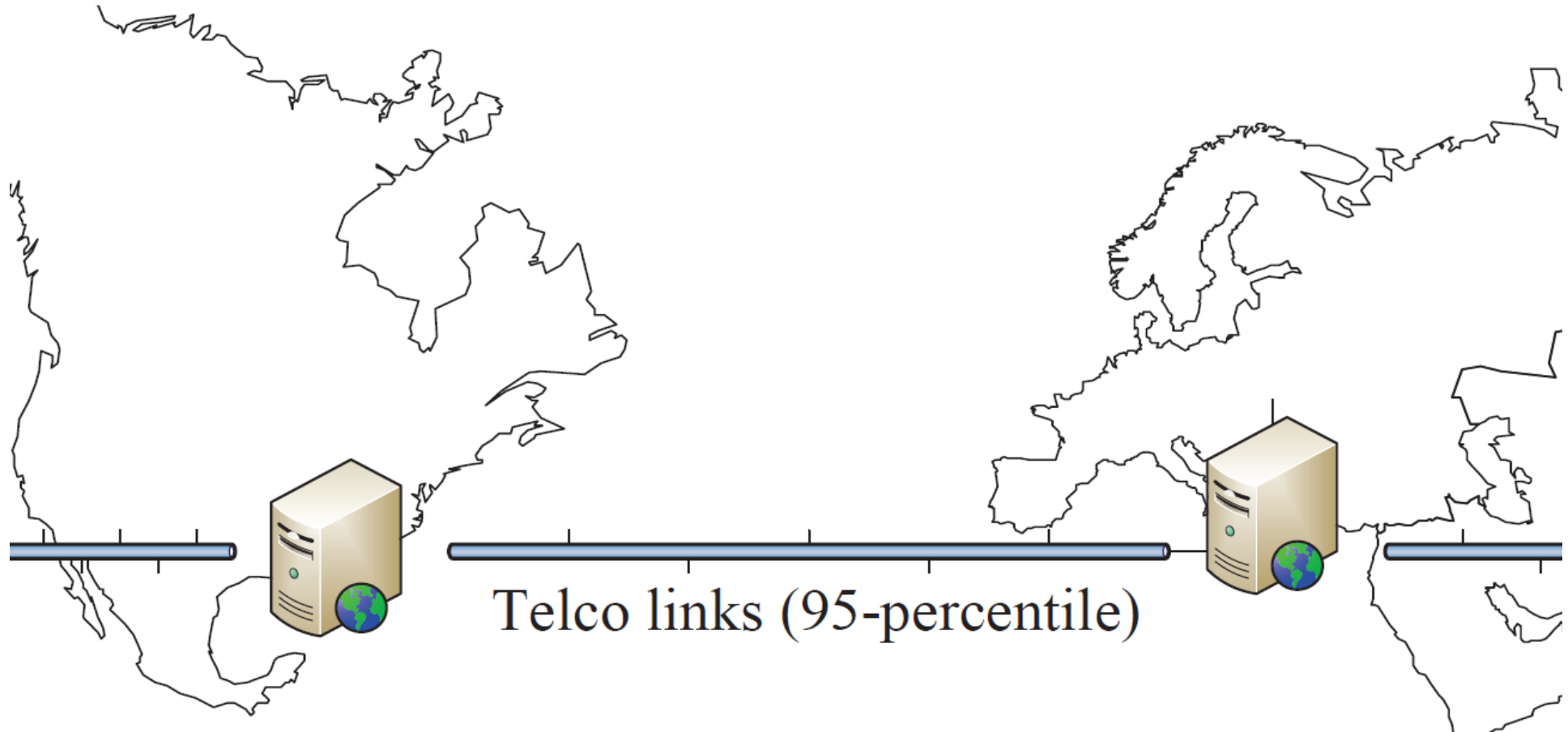


Trend 1 – Fault tolerance to catastrophic failures



With incredible redundancy, comes incredible cost. Even with incredible costs, failure modes remain that can eliminate the facility entirely. The only cost effective solution is to run redundantly across multiple data centers. Redundancy without physical separation is not sufficient and making a single facility bullet proof has expenses asymptotically heading towards infinity with only tiny increases in availability as the expense goes up. The only way to get the next nine is have redundancy between two data centers. This approach is both more available and considerably more cost effective.

Trend 2 – PoP replication for improved user QoS



The Issue: Traffic Growth Is Outpacing Business Growth

A Forrester Consulting Thought Leadership Paper Commissioned By Infineta Systems

The Future Of Data Center Wide-Area Networking

Why Skyrocketing WAN Bandwidth Requires A New Approach To Controlling Traffic

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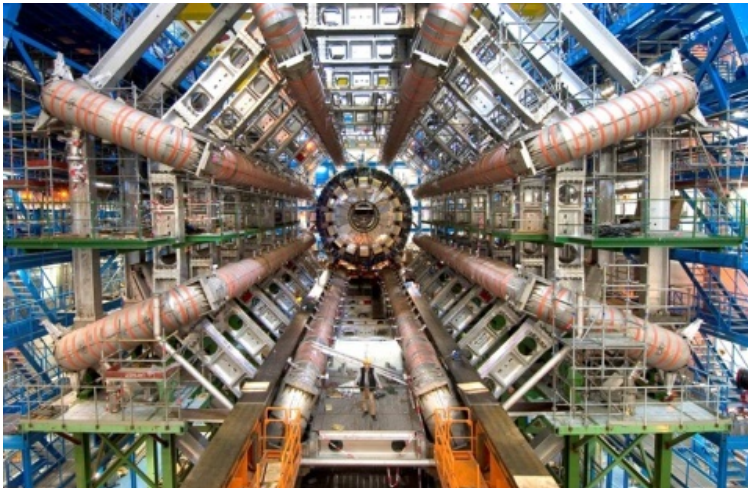
Making Leaders Successful Every Day

Storage, and the replication of it, places enormous demands on the WAN. More than half of our respondents have over a petabyte of storage in their primary data center. Where does the traffic come from? The replication of all that stored data between sites. The vast majority (77%) of the firms we surveyed have three or more data centers running in an active/active or active/passive configuration, and most had replication processes for BC/DR, backups, archives, and data warehouses.

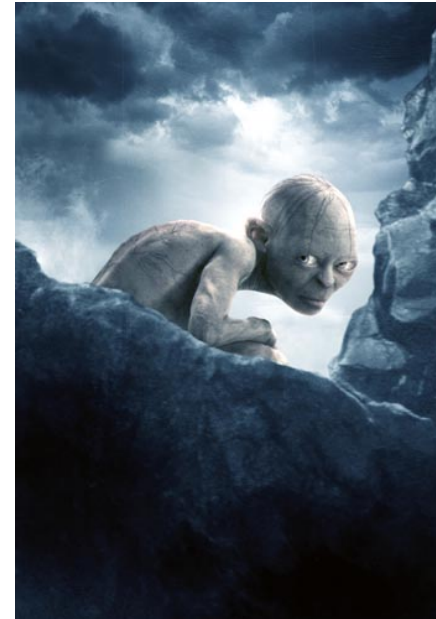
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Additional applications

- Scientific computing



- Distributed production/delivery of movies

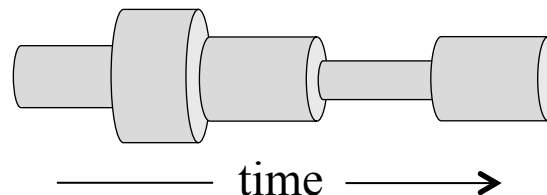
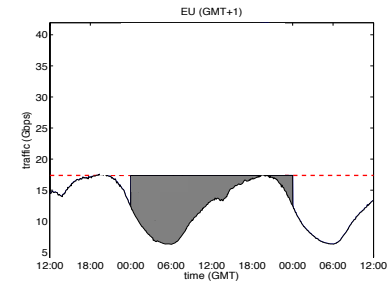


NetStitcher in a nutshell

- A system for carrying bulk data for large customers
 - Volume ~ Tbytes / Pbytes
 - Delivery time ~ hours / days
- Main idea
 - Peak load dimensioning & backup paths → lots of leftover bandwidth
- Create a volume service for interconnecting datacenters



X TBs from A to B within the next Y hours



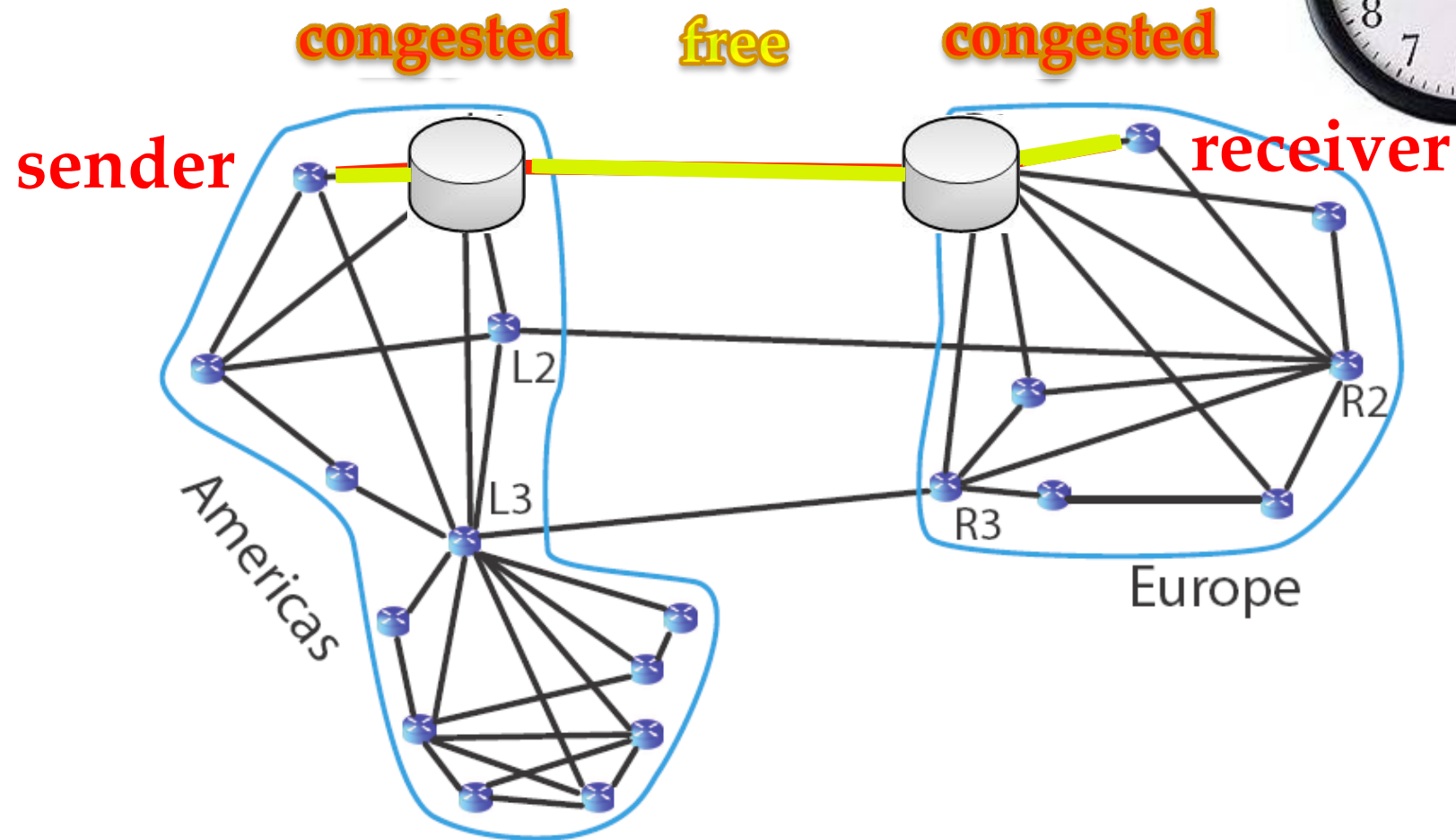
Let's play devil's advocate for a while

Why do we need
NetStitcher ?

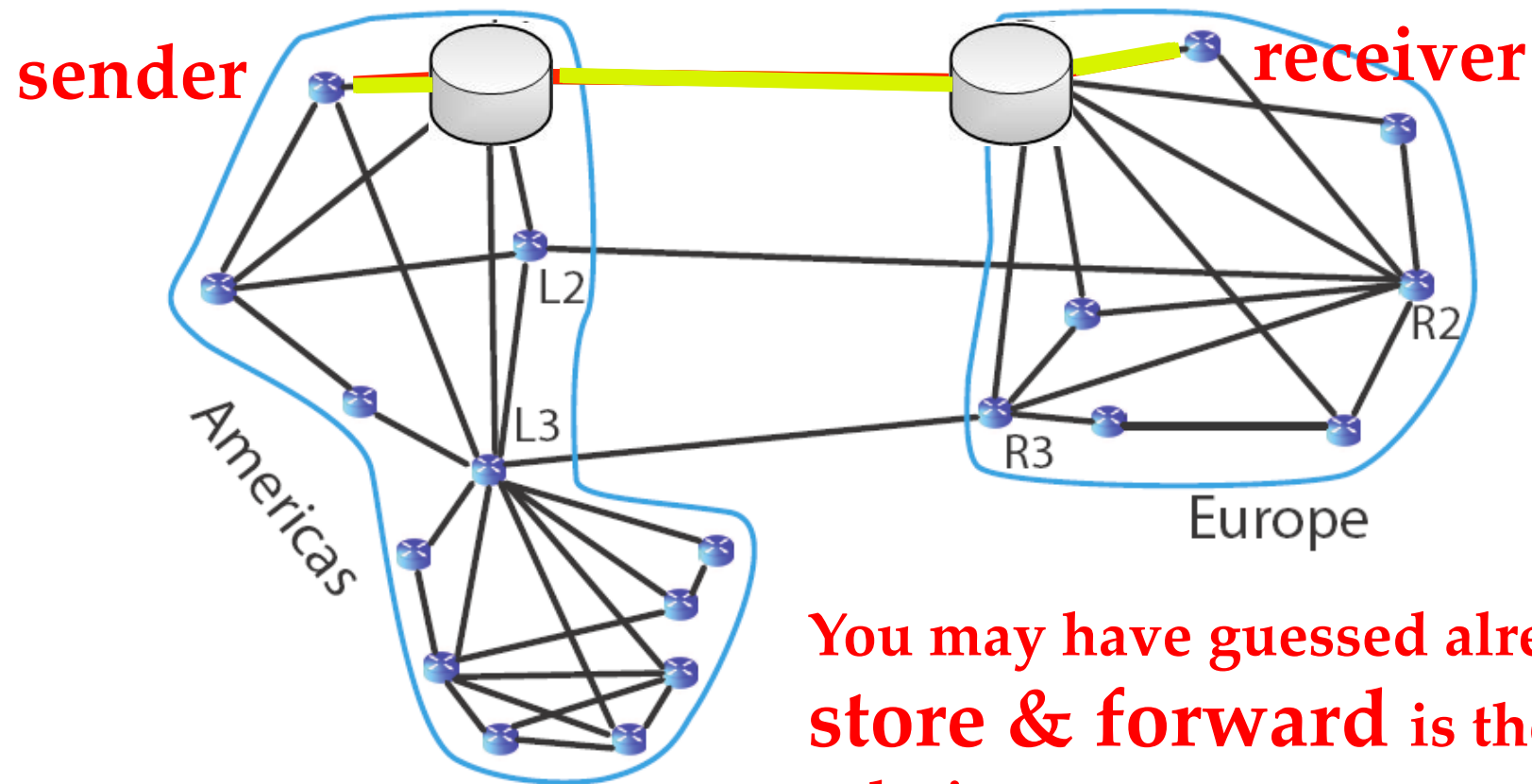


- Wide area bandwidth is not cheap enough to cover traffic growth
 - This applies independently of pricing scheme
 - 95-percentile based, flat, or owned lines
- Difficult to collect leftover bandwidth using End-to-End solutions

Leftover b/w appears whenever and wherever

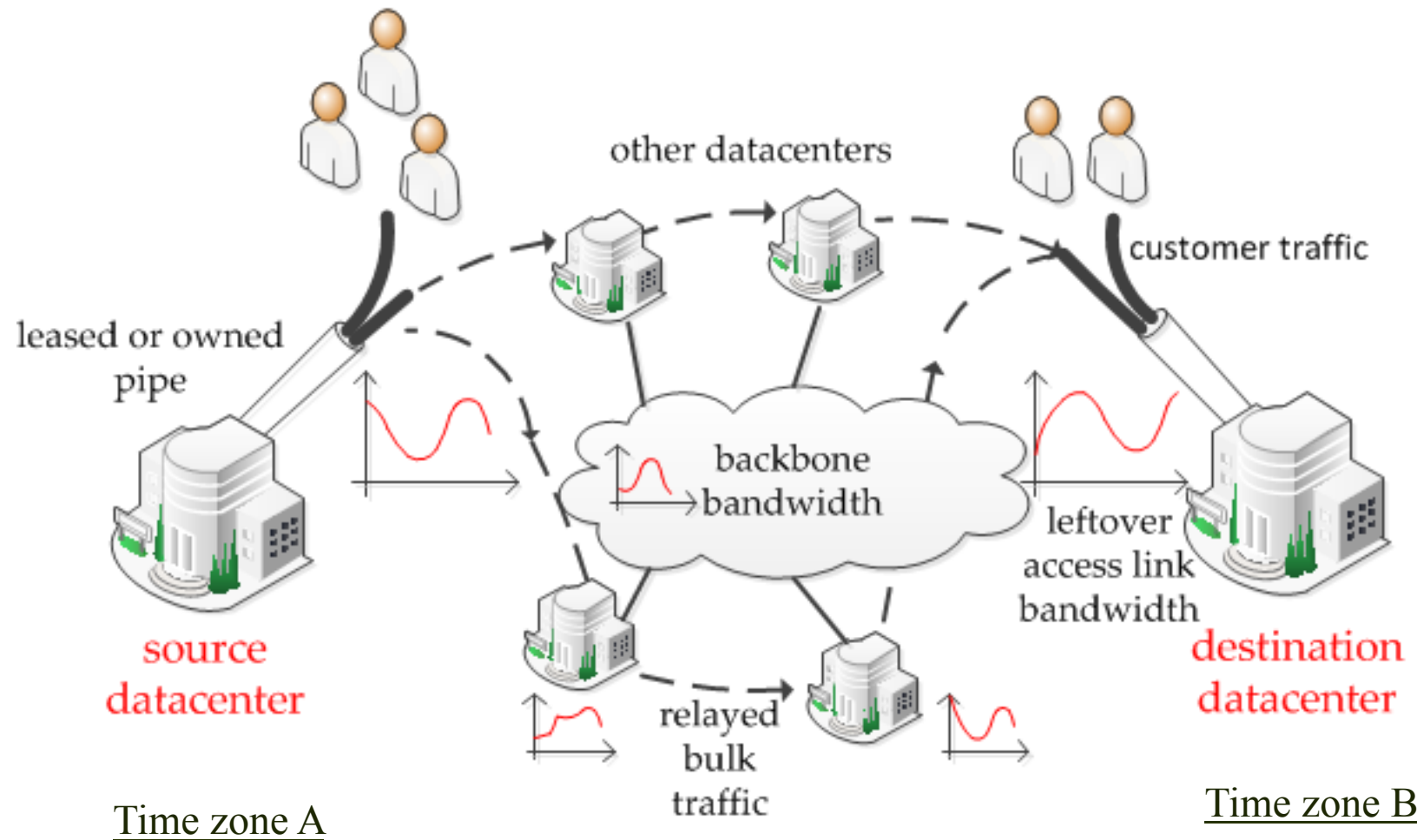


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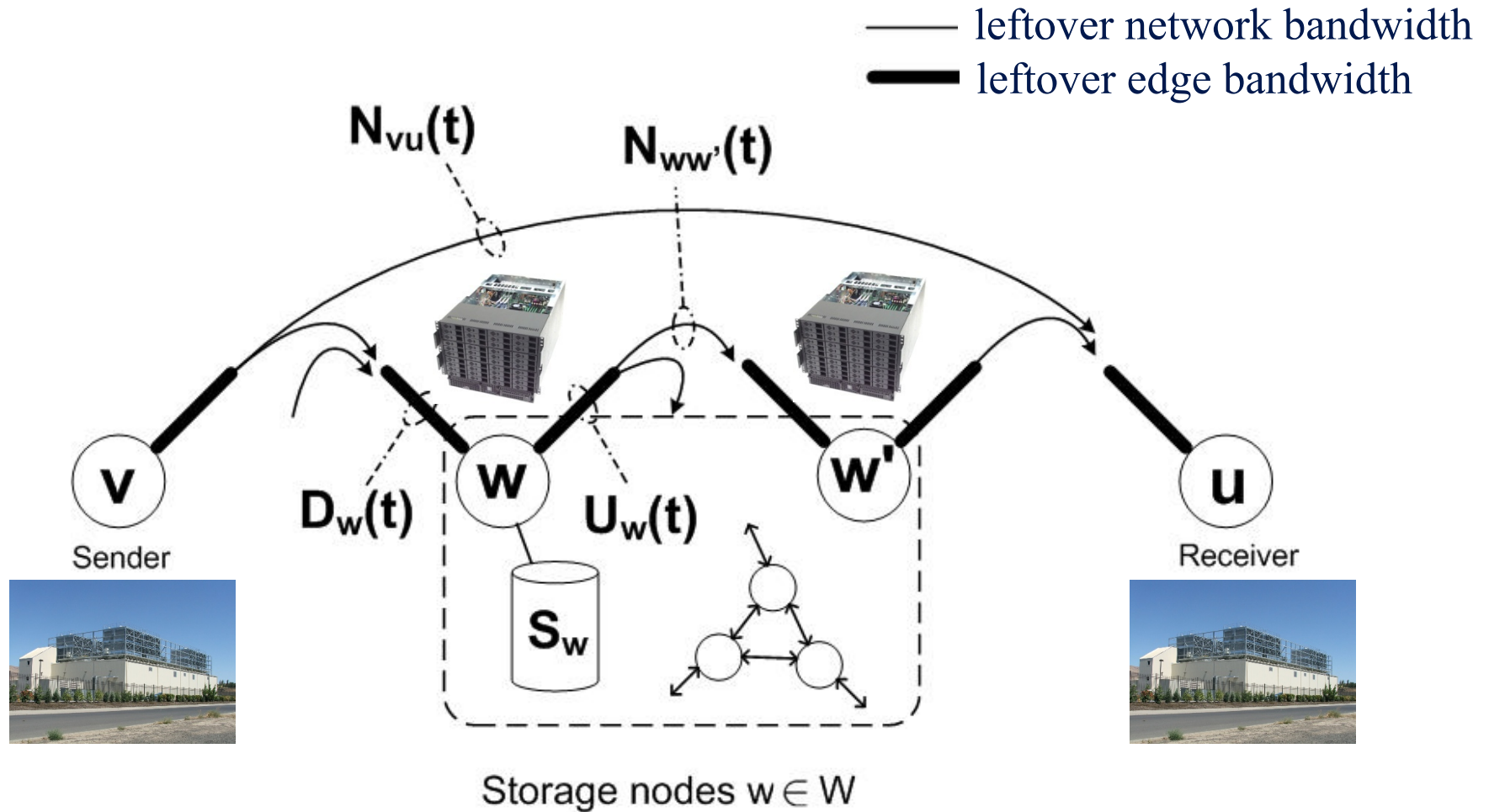


You may have guessed already:
store & forward is the
solution

Stitching together leftover bandwidth is tricky

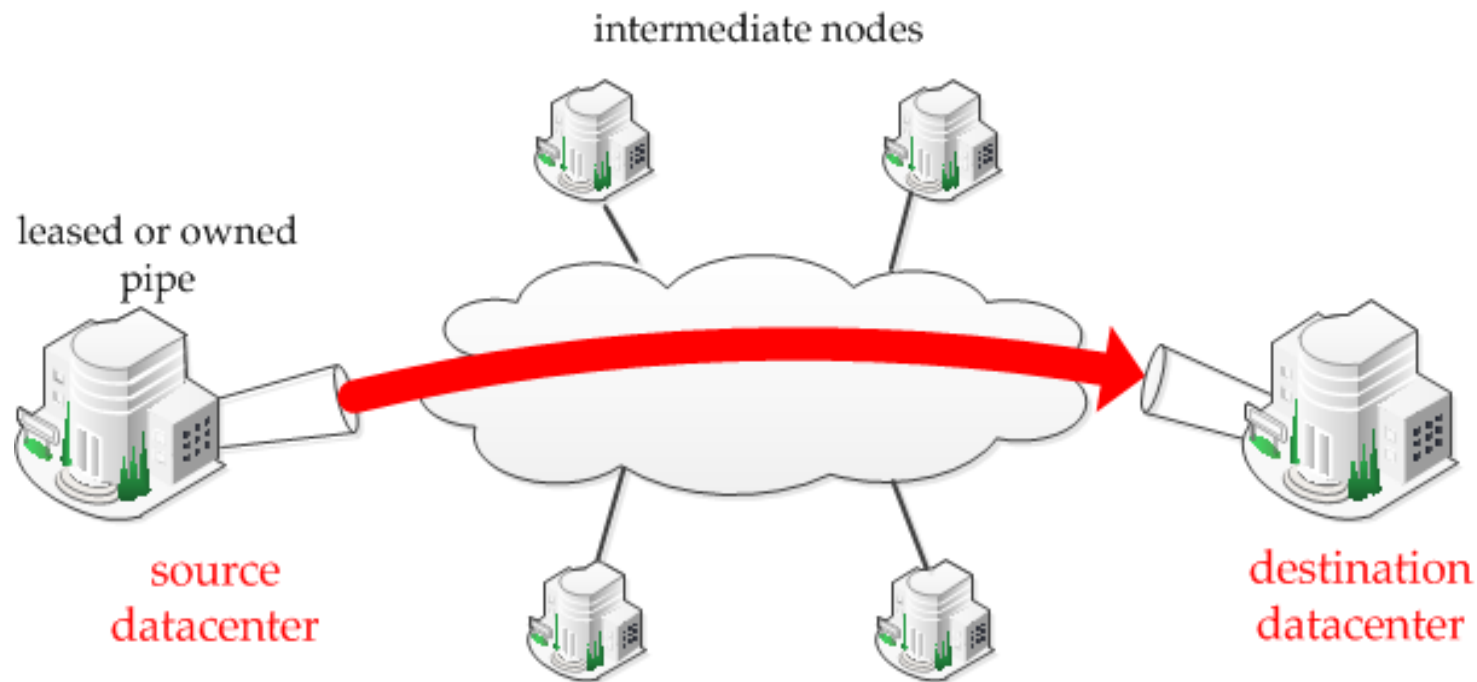


A storage overlay that is aware of leftover bandwidth



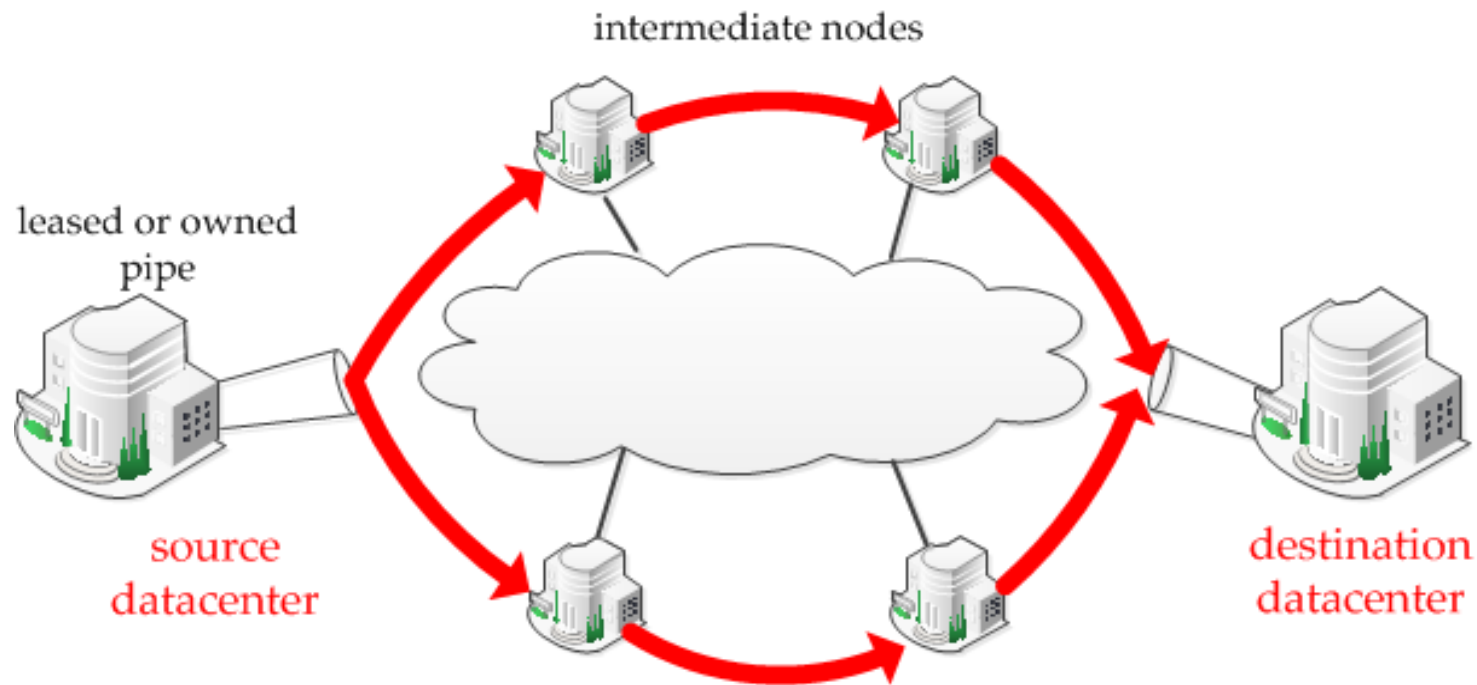
NetStitcher's bag of tricks

- No in-network constraints and time-aligned sender and receiver bandwidth availability
 - NetStitcher can perform end-to-end transfer



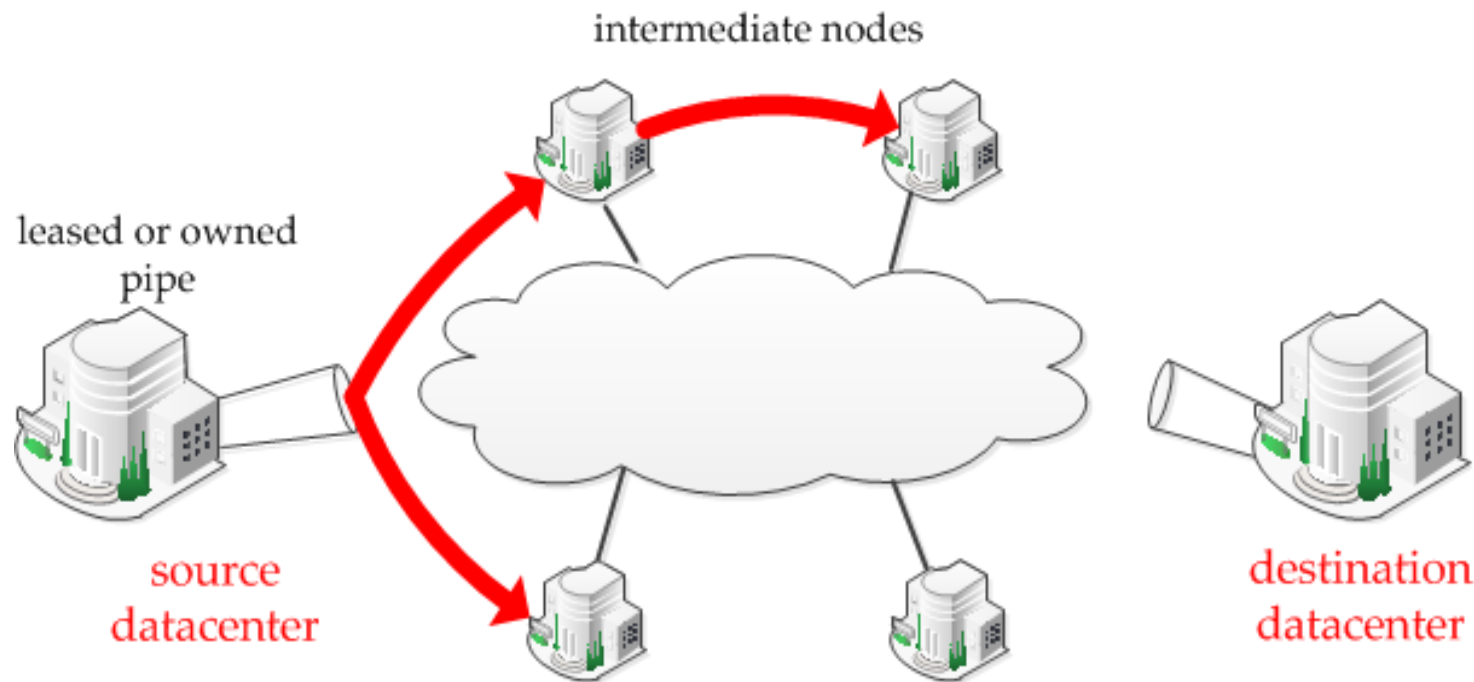
NetStitcher's bag of tricks

- In-network constraints and time-aligned sender and receiver bandwidth availability
 - NetStitcher can perform multi-path overlay routing



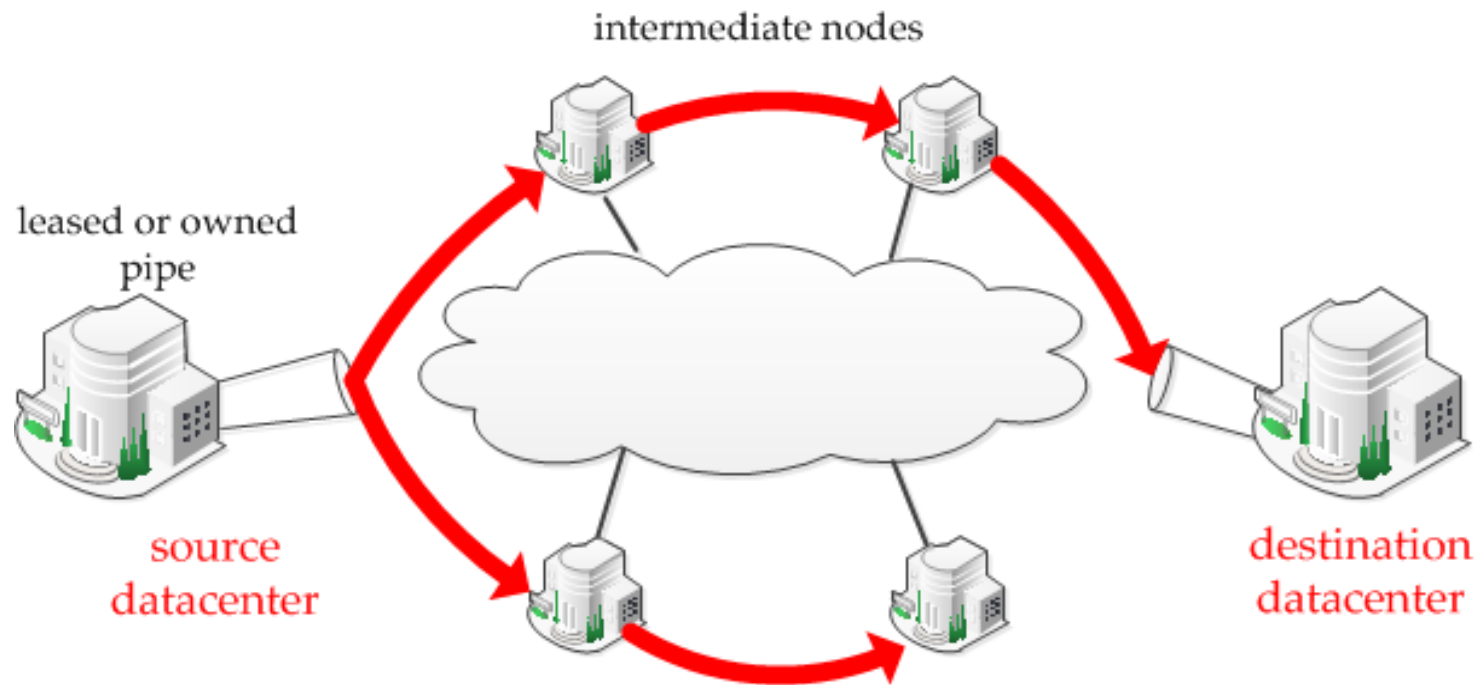
NetStitcher's bag of tricks

- In-network constraints and misaligned sender , receiver and intermediate node bandwidth availability
 - NetStitcher can perform multi-path and multi-hop store and forward



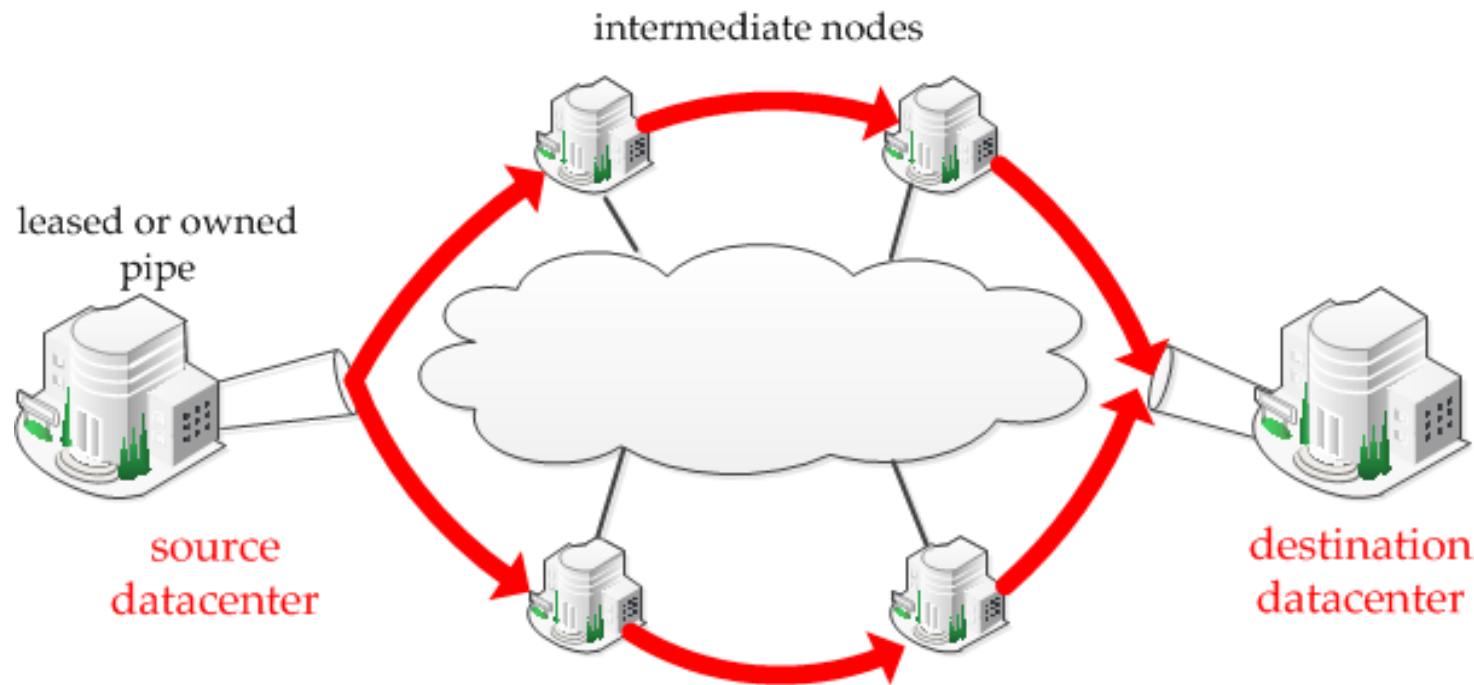
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NetStitcher's bag of tricks

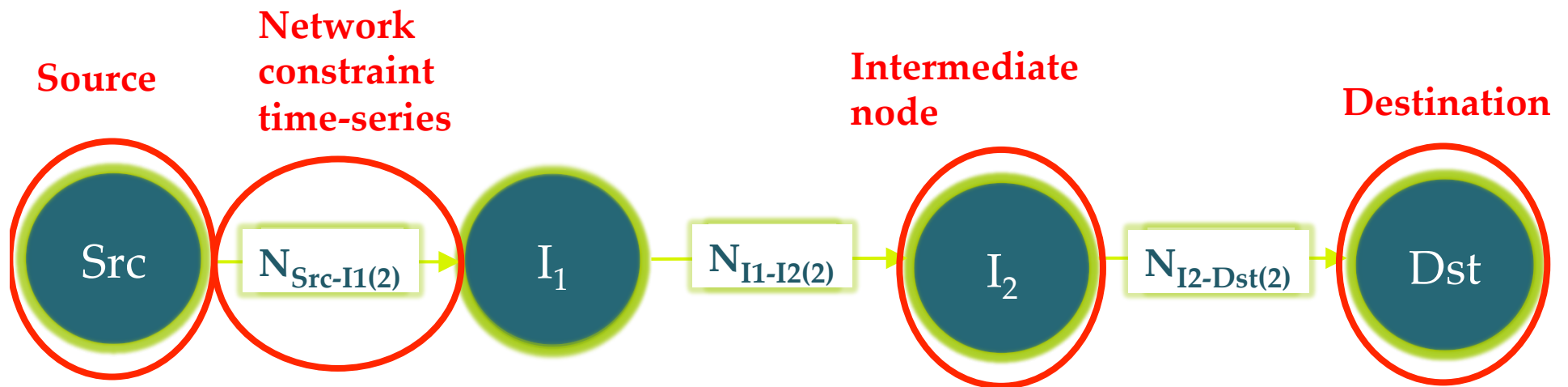
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How do we
schedule around all
these constraints?

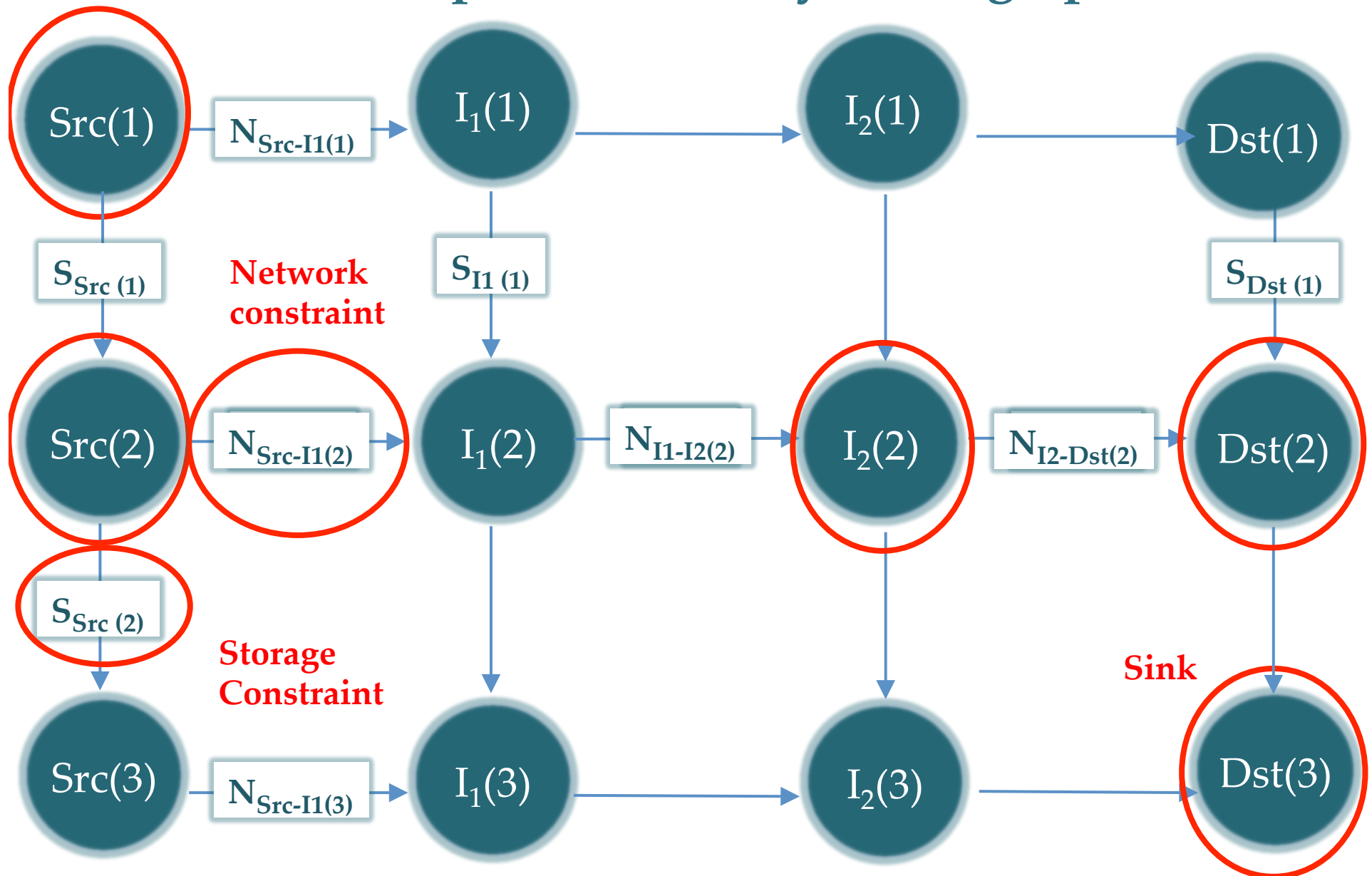


Time expansion of a dynamic graph

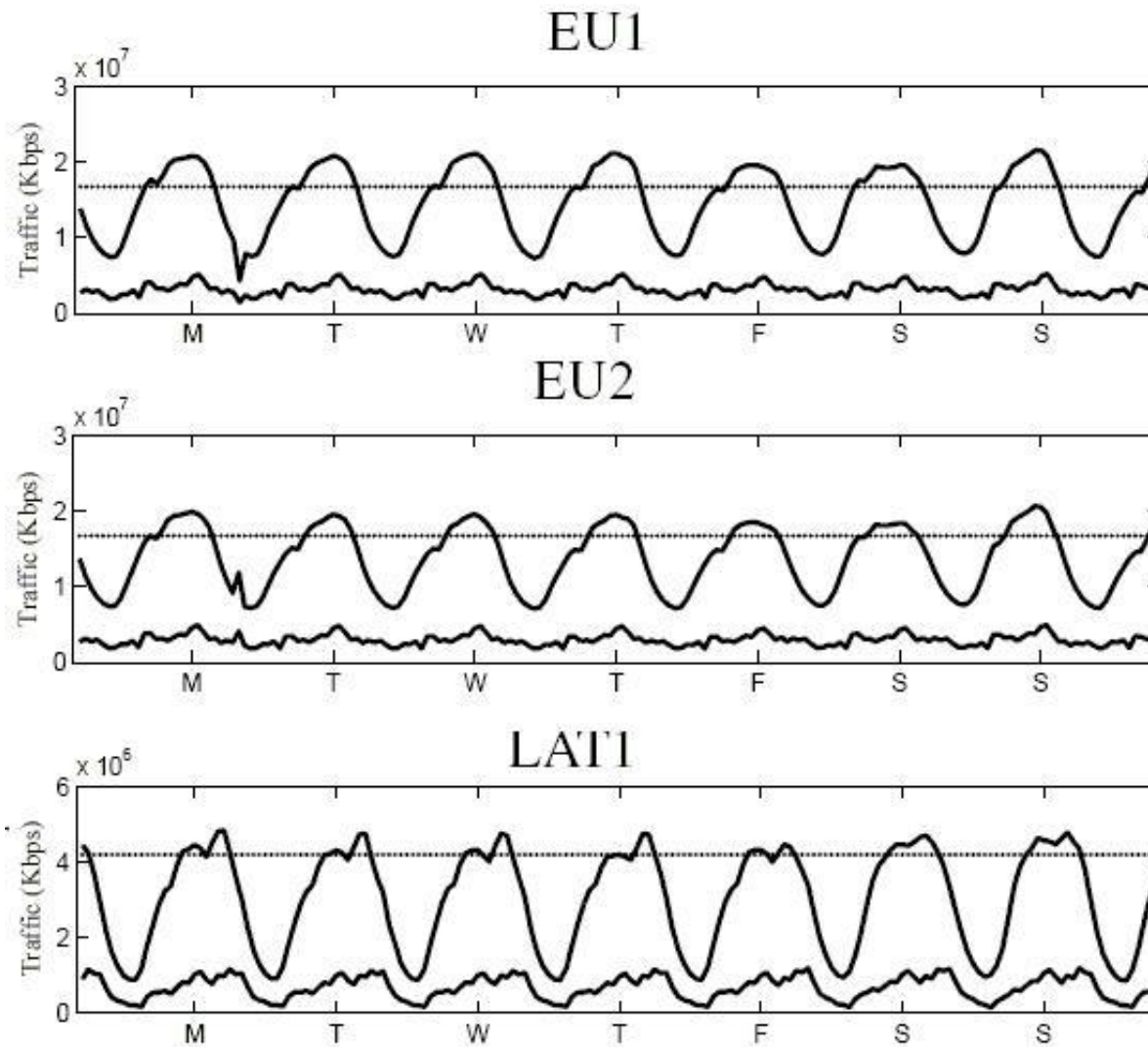


Source

Time expansion of a dynamic graph



But we need to predict the future (of bandwidth)



International backbone traffic



Prediction is easy when data are bulk

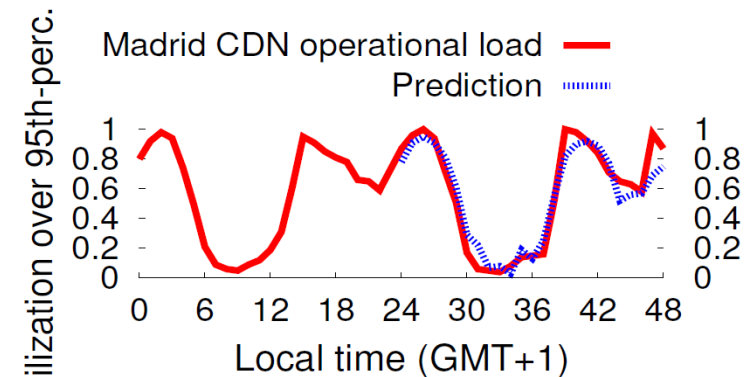
1. Periodic patterns

2. We care about **VOLUMES** not **RATES**

$$\text{VOLUME} = \int_{t_0}^{t_0+T} \text{RATE}(t) dt$$

In our NetStitcher implementation we use:

- A simple Sparse Periodic Auto-regression Predictor (Chen et al., NSDI'08)
- Recomputation of transmission schedule
- End-game “pull mode” to handle occasional churn/prediction failure



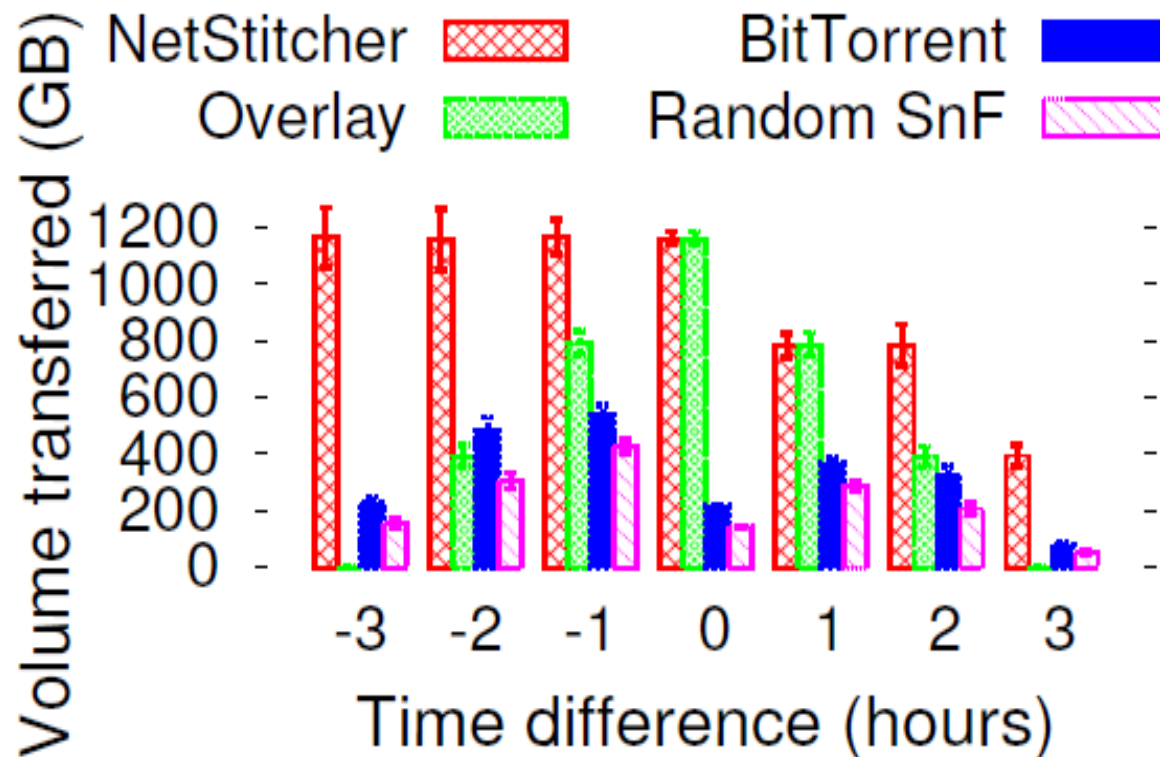
Case study 1: Equinix datacenters



- 62 datacenter at 22 locations all over North America

How much data can we backup?

- 3 hours used for backup (3-6 am local time at the datacenter)
- 1 Gbps network access capacity



NetStitcher can move ×5 more bytes

Case study 2: Telefonica CDN

- 49 servers in Europe, Latin America and USA. GMT-1 to GMT-8
- Need to send a 4.2 TB file over 24h.
- Beyond leftover → 95th-percentile pricing with \$7/Mbps/month
- Storage cost: \$0.055 GB/month



NetStitcher 80-90% cheaper between Europe & US

Conclusion: A practical application of DTNs

- The utilization of a network can be improved but for this we need:
 1. Delay elastic traffic to go into off-peak hours
 2. In-network storage
 3. High-level knowledge of traffic behaviour around the day

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More info at: <http://people.tid.es/Nikolaos.Laoutaris>