

Application-aware Data Plane Processing in SDN

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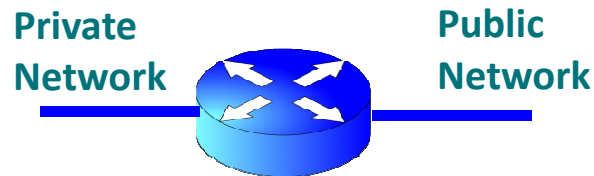
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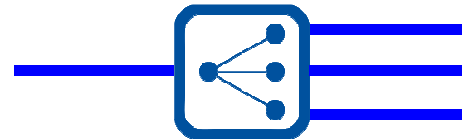
Motivating Examples

NAT



- IP & port rewrite
- Keep track of available public IP & port

Content based server selection



- Assign server IP based on URL
- Maintain flow state to support TCP splicing

Can we extend the SDN data path to directly support such network services ?

Application-aware data plane processing in SDN

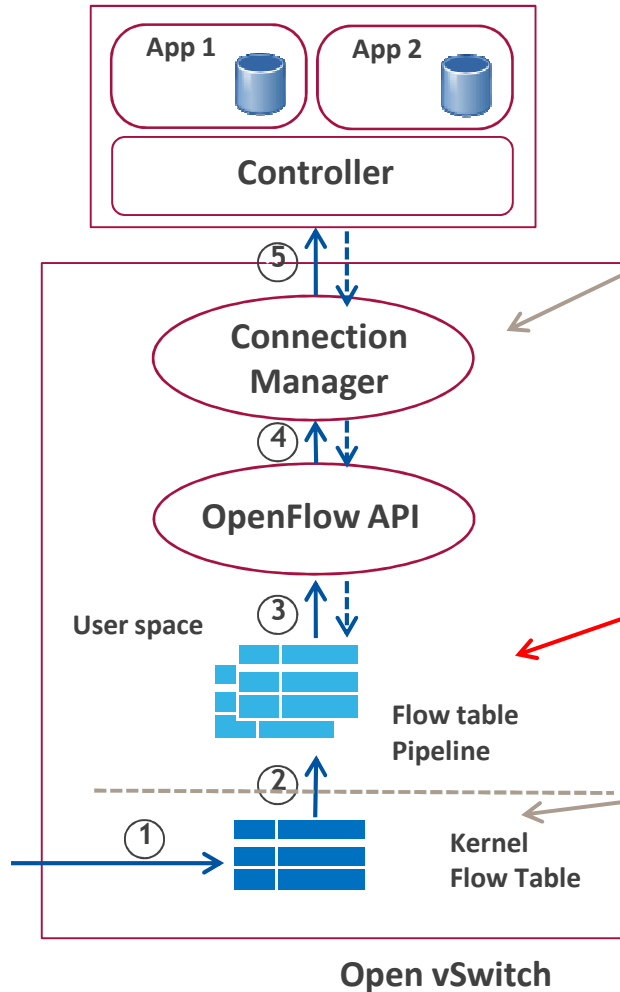
Alternatives:

- Using middleboxes
 - More equipments to maintain
 - Extra hops (detour) in data path
 - Complexity of service chaining
- Using SDN controller to implement application processing logic
 - Switch-Controller delay cause slowdown in data path
 - Control plane not designed to handle every packet -> throughput bottleneck

Application-aware data plane processing:
One system for both routing and network services with **uniform central control** and **scale-out data plane**



Open vSwitch



Design Choice: Where to intercept the packet and implement application processing logic ?

Option 1: connection manager

- Pros: modular design
- Cons: redundant coding, slow
 - Unnecessary encap/decap
 - Redundant flow table

Option 3: user space flow table

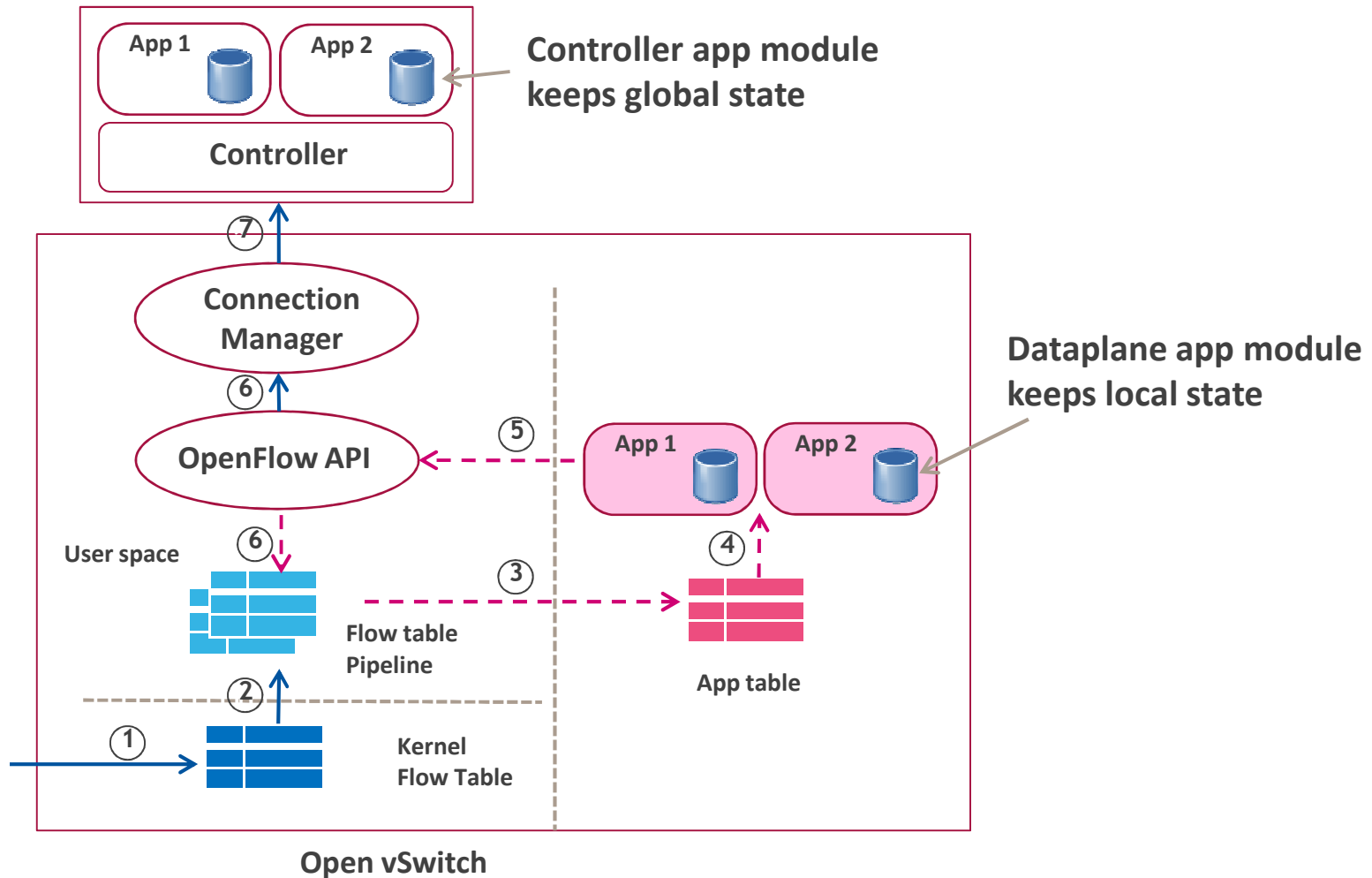
- **Good tradeoff between 1 & 2: easy implementation & reasonable performance**

Option 2: kernel flow table

- Pros: best performance
- Cons: hard to implement

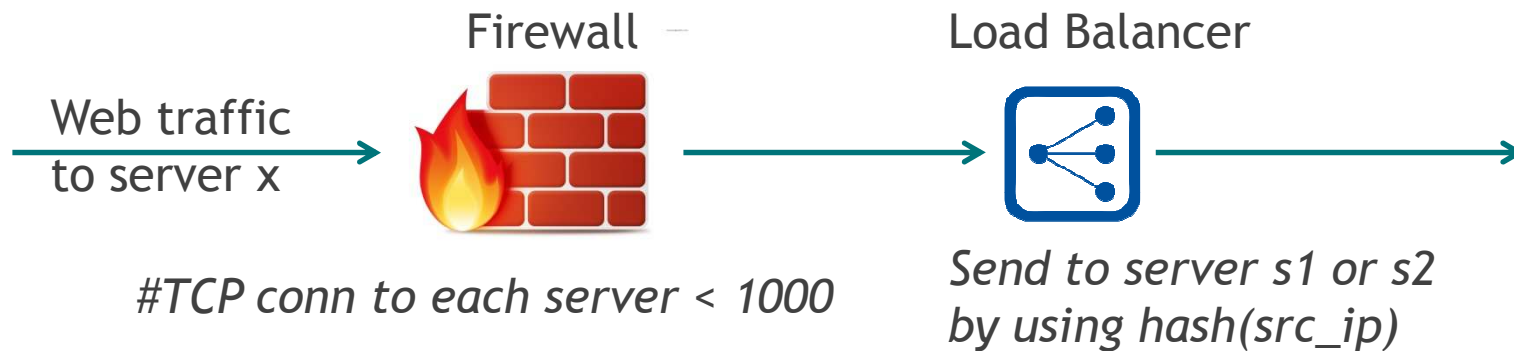


Application-aware Data Plane for Open vSwitch



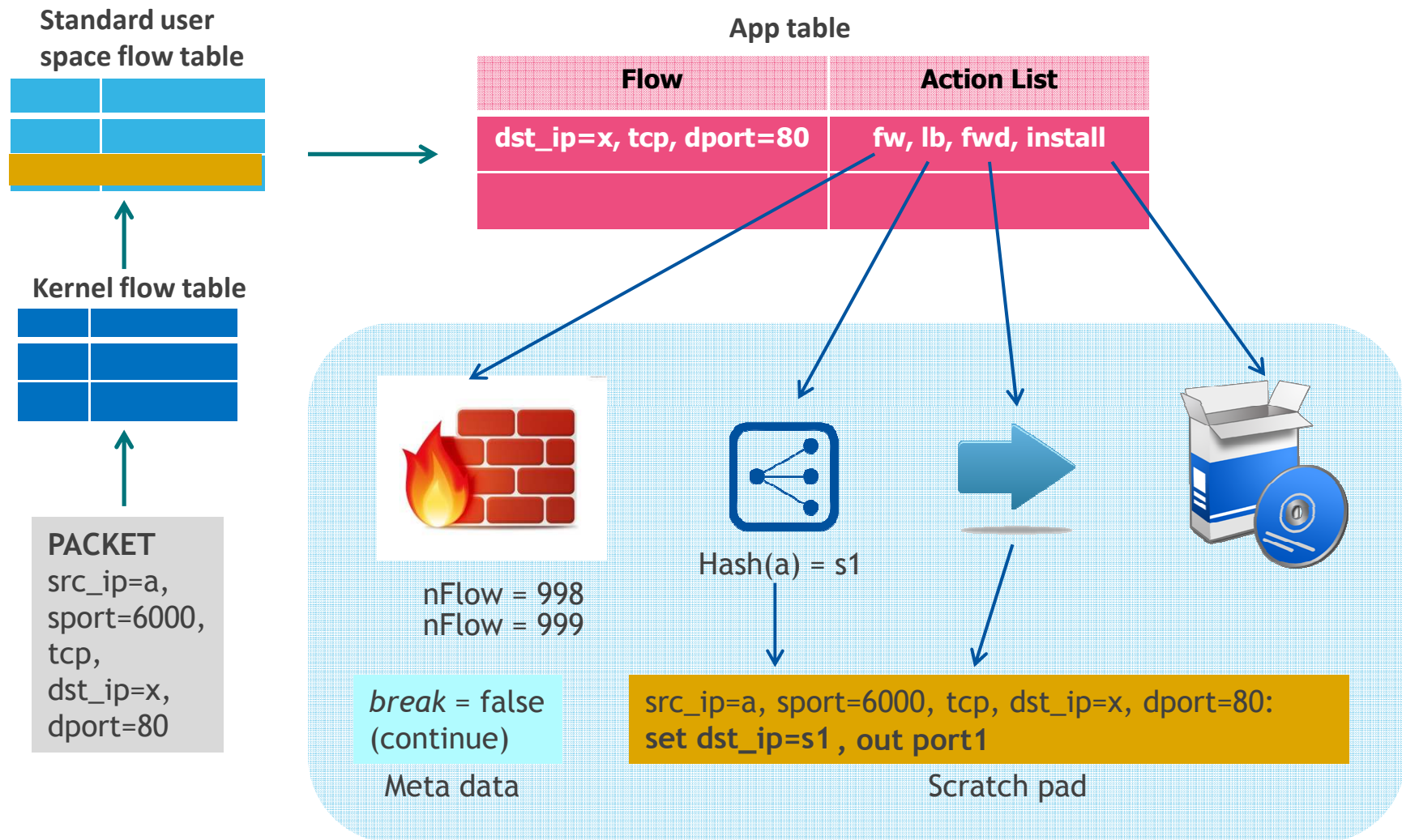
Example: Firewall & Load Balancer

Required Policy:

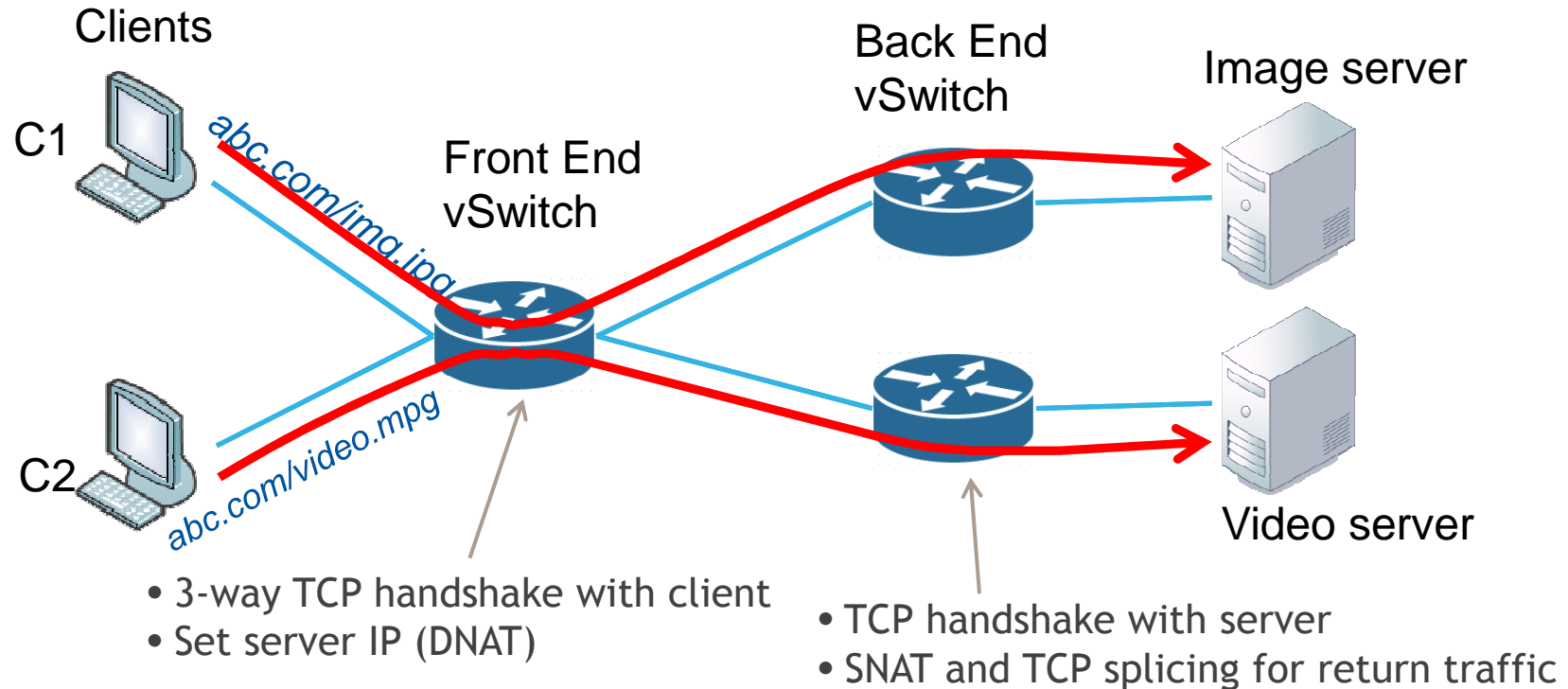


App table rule: `dst_ip=x, tcp, dport=80: fw, lb, fwd, install`

Firewall & Load Balancer: Implementation

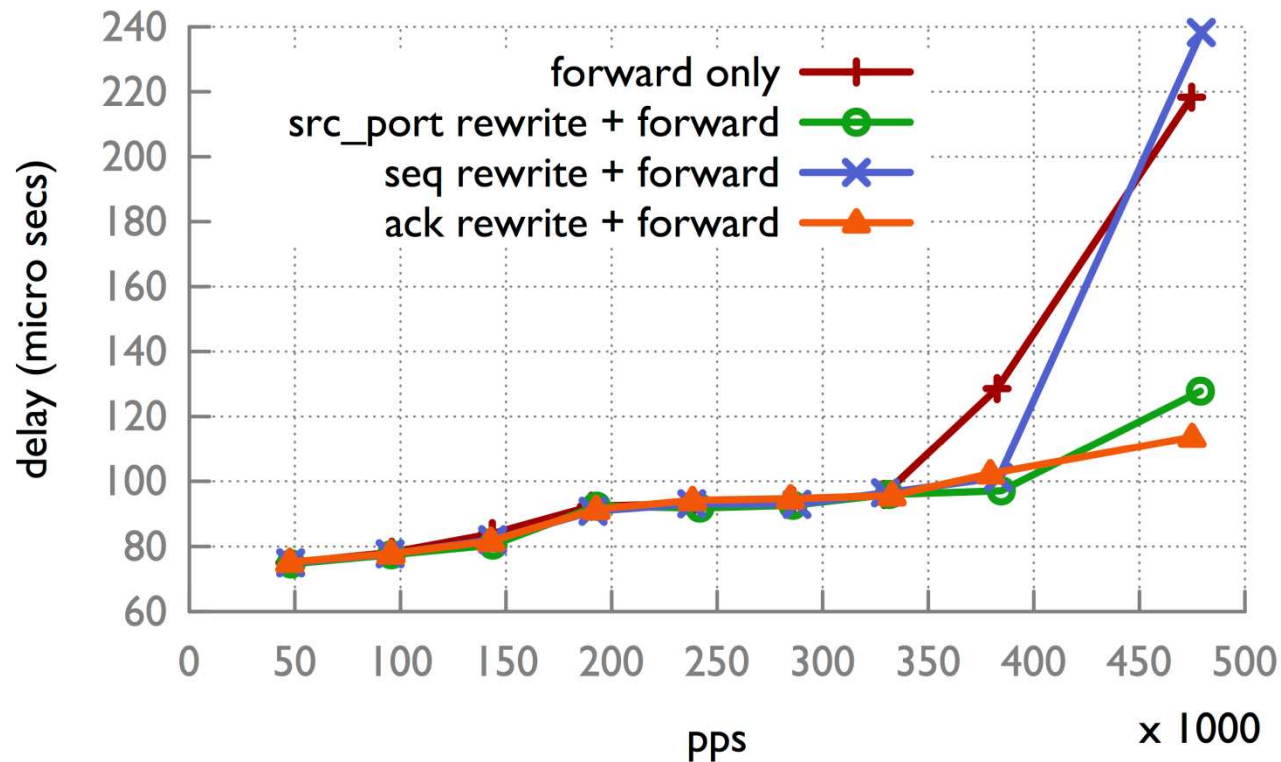


Example 2: Content Aware Server Selection



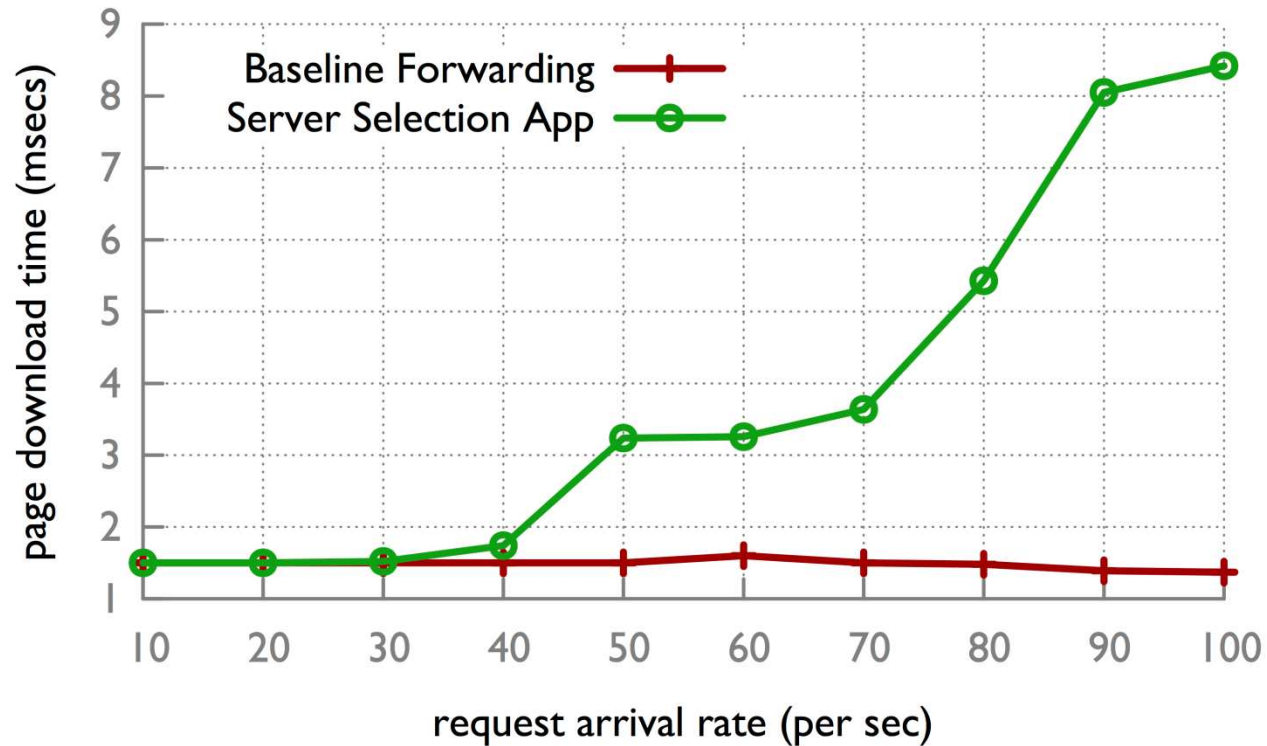
- Return traffic does not have to go through front-end vSwitch
- Both front-end & back-end vSwitches can be scaled out independently

Experimentation: TCP splicing



- TCP sequence number rewriting implemented in kernel space
- Very good performance: same as native OVS kernel actions

Experimentation: Server Selection



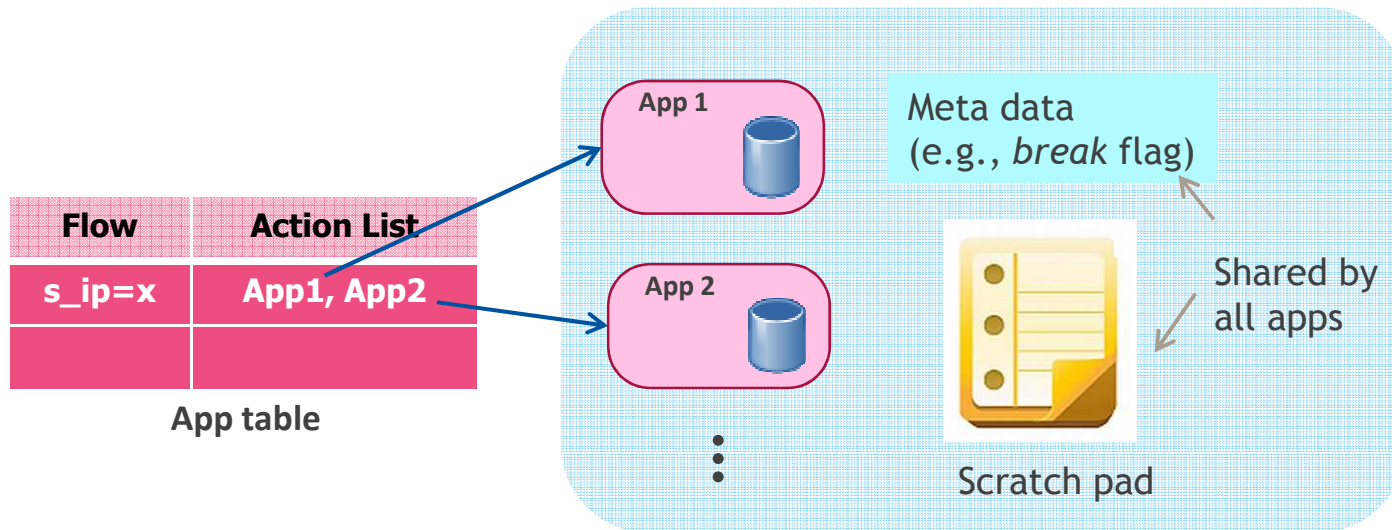
- Server selection action implemented in user space
- Performance can be significantly improved by using the new multi-threaded OVS daemon implementation (on-going work)

Conclusions and Future Work

- A first step towards enabling application-aware SDN data plane
- Overall, extending software switches like Open vSwitch to support network service applications is a promising direction
- Kernel to user space copying is still a performance barrier
 - Careful design choices to trade-off between ease of implementation and performance
 - Current design doesn't suite applications that require processing *every packet*
- Future work
 - Take advantage of muti-threaded OVS daemon in recent versions
 - Make the apps “pluggable”
 - DPDK

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Design Details: App Table, App Actions, App Chaining



- App actions implemented as *vendor actions*, only run in user space
 - Determine actions for current packet
 - Modify local app state
 - Generate/modify rule set to be installed in standard flow table
 - Remove flows from standard flow table
 - Generate packet out
 - Send Packet_In, Flow_Removed, or app update vendor msg to controller... by calling Openflow API module