

# ACM SIGCOMM 2018 Workshop on Mobile Edge Communications (MECOMM'2018), Budapest, Hungary

# Anveshak: Placing Edge Servers In The Wild

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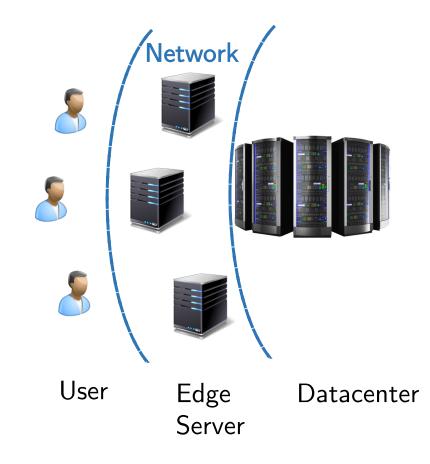


# **Edge Cloud Overview**

**Edge Cloud**: Small-scale server(s) deployed at network edge to compute user data

#### Motivation:

- ✓ Decreased latency and network traffic
- ✓ Computing data of local relevance

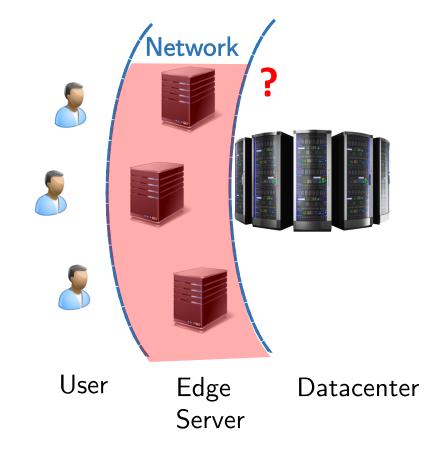


## **Edge Cloud Overview**

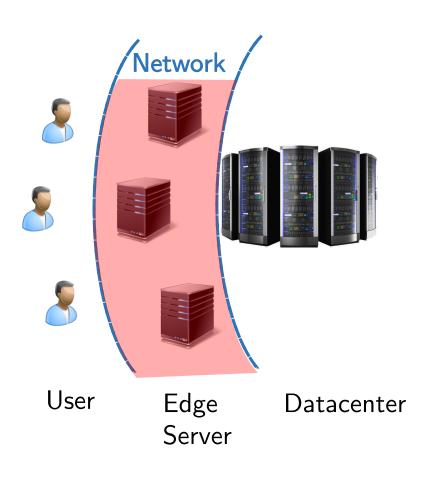
**Edge Cloud**: Small-scale server(s) deployed at network edge to compute user data

#### Motivation:

- ✓ Decreased latency and network traffic
- ✓ Computing data of local relevance



# **Edge Cloud Deployment**



- 1. What are these servers configuration?
- 2. Who owns and operates them?
- 3. Where are these servers going to be deployed in the physical world?

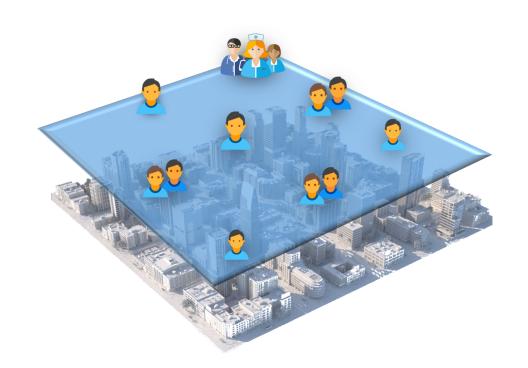
# Physical Edge Cloud Network



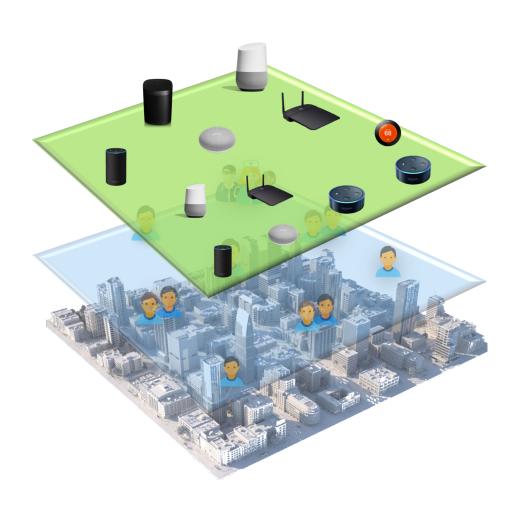
- 1. Multiple edge cloud deployments can co-exist in same physical space
- 2. Edge server availability will directly affect end utilization
  - Usage of cellular-based internet in areas where WiFi is available

#### **Users**

- 1. Subscribers and requesters of edge resources
- 2. User density directly impacts local edge server utilization
- 3. Request distribution is highly dependent on **time** and **behavior** 
  - ➤ More user requests during daytime than in night
  - ➤ More user requests in city than suburban areas

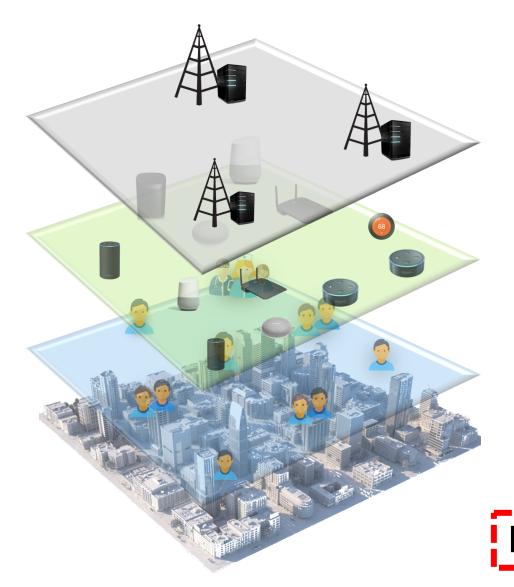


#### User-Managed Edge



- 1. Composed of: self-managing, locally-relevant edge resources
  - ➤ e.g. smart speakers, home automation, intelligent WiFi hubs
- 2. Limited computational-power and network capability
- 3. High server density
  - > Dependent on user population
  - > One server caters to small set of users

 ${\sf Management/Control} o {\sf End-users}$ 



# Service Provider Managed Edge

- 1. Composed of: high computation and network capable edge servers
- 2. Set up specifically by cloud provider in partnership with local ISP
- 3. Co-located/accessible with cellular base stations for *ease-of operation* and *maintenance*
- 4. Low server density
  - ➤ One server caters to large set of users

 ${\sf Management/Control} o {\sf Cloud provider/ISP}$ 

**New Server Deployment\*** 

\*from Service Provider perspective

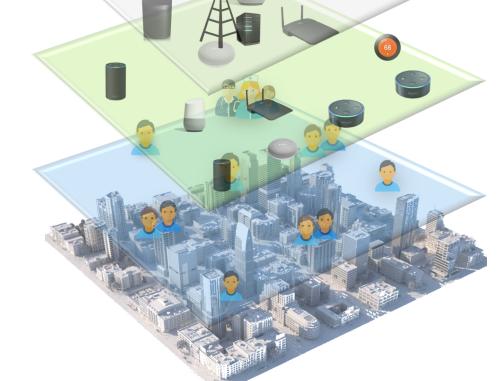
Which location to deploy new edge server?

#### Over-provisioning!

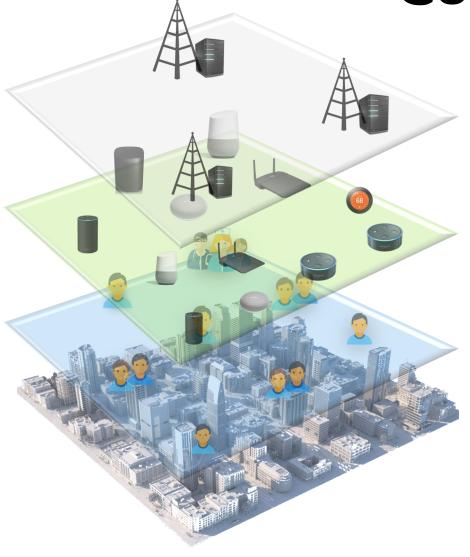
- > Connectivity can be overlapping
- Future utilization can be minimal OR

#### **Under-provisioning!**

- ➤ Maintain a Quality-of-Service
- > Must support peak request traffic



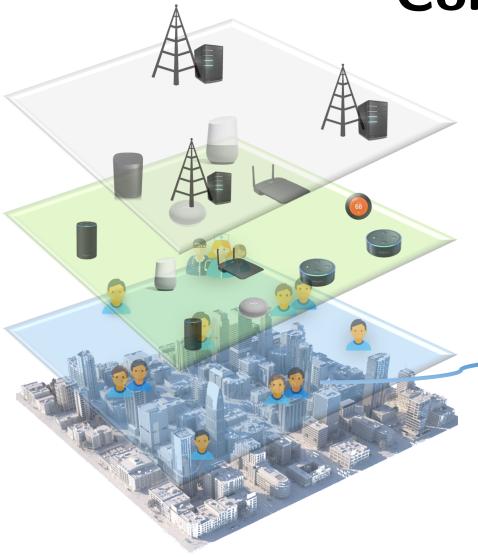
#### Considerations



# Deploying managed edge servers is expensive!

An efficient server deployment algorithm **must**:

#### Considerations

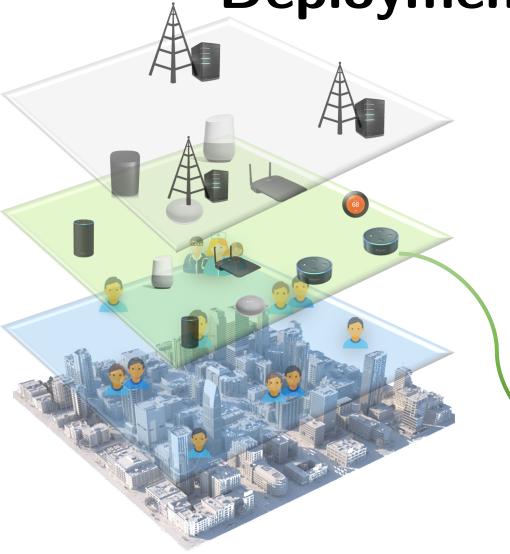


# Deploying managed edge servers is expensive!

An efficient server deployment algorithm **must**:

•1. Prioritize areas with high user requests





# Deploying managed edge servers is expensive!

An efficient server deployment algorithm **must**:

- 1. Prioritize areas with high user requests
- 2. Avoid areas with high usermanaged edge resources

# Anveshak

# System Workflow

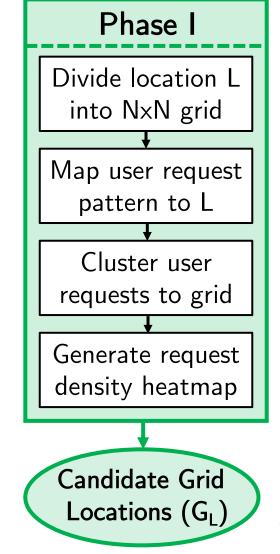


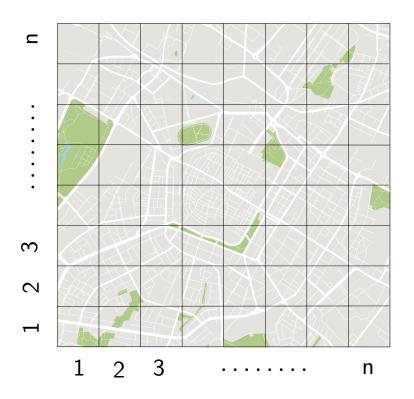
- > Three-phase waterfall-based workflow
- > Intermediate phase checkpoints for recoverability
- Swappable phase modules for incorporating improved algorithms and parameters
- Anveshak is *first-of-its-kind* framework which considers usermanaged edge in location selection



AIM: Identify and prioritize areas of high user communication requests

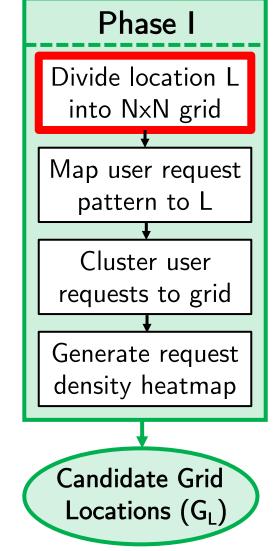
- ➤ High server utilization
- > Low user-server connection latency

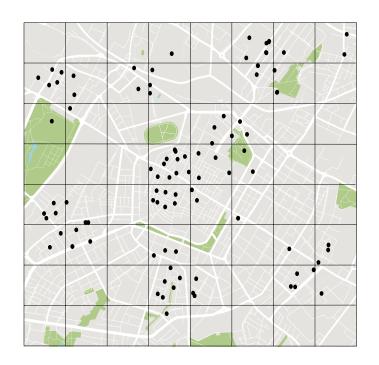




Divide map into equi-spaced NxN grid

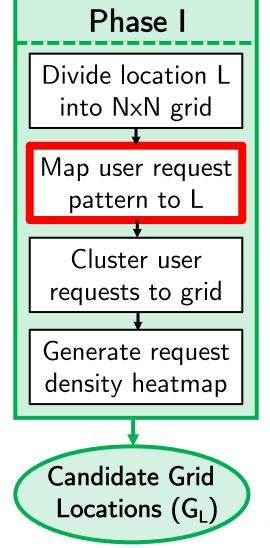
- > Grid division allows for consistent clustering
- > Grid size directly affects problem search space

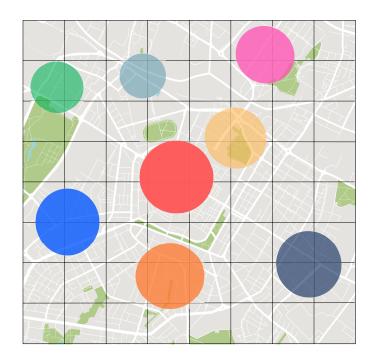




Past user communication requests such as *CDRs*, internet initiation are mapped on location

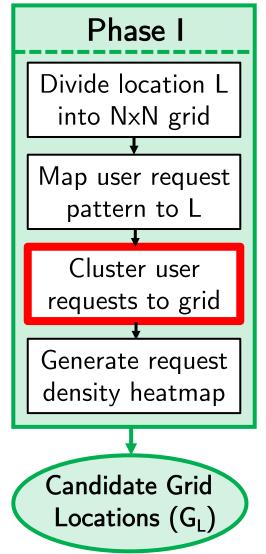
Requests are averaged over time to remove temporal outliers





Cluster user requests based on inter-request distances and densities

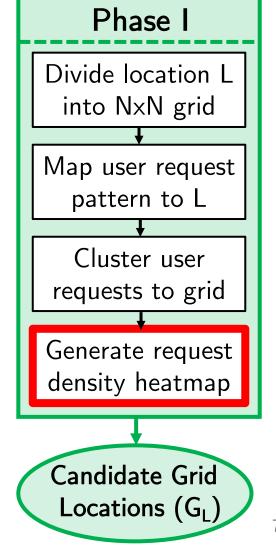
Choice of clustering algorithms and their parameters can be easily tweaked [**DBScan** used as example]

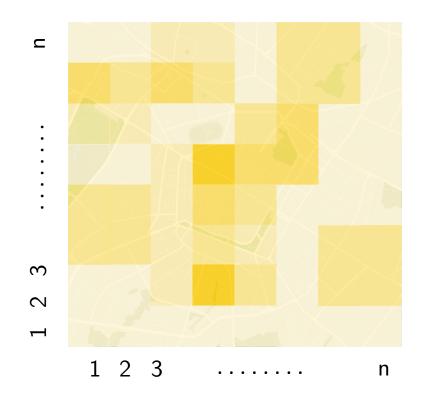




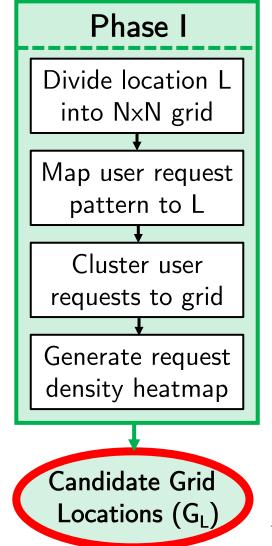
Generate heatmap for arbitrary cluster shapes

- > Handles over-lapping shapes, small/dense clusters
- > Handles any inefficiency of clustering algorithm





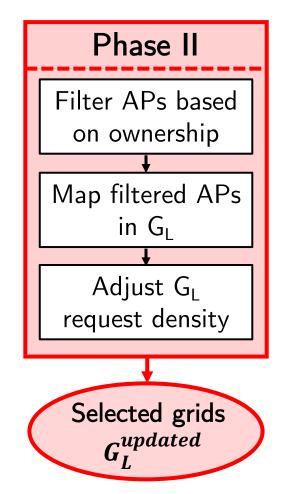
 $G_1 = \{Grid ID, Request Density\}$ 



- Anveshak estimates future deployment of usermanaged edge resources
- Availability of user-edge servers will limit utilization of deployed edge in same location
- Such devices are highly dependent of user population and interaction in an area

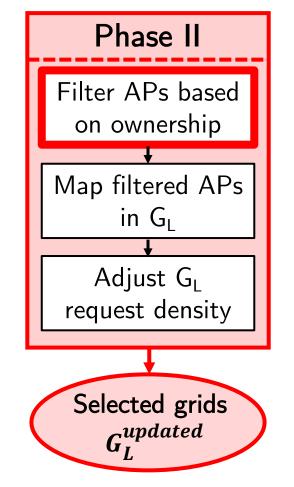
Estimated via current deployment of WiFi access points

**INPUT**:  $G_1 = \{Grid ID, Request Density\}$ 

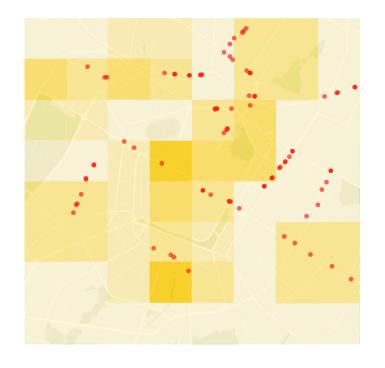


Map currently deployed WiFi access points in the same area

- Utilize open datasets for WiFi access points such as wigle.net
- Filter out mobile and temporary access points

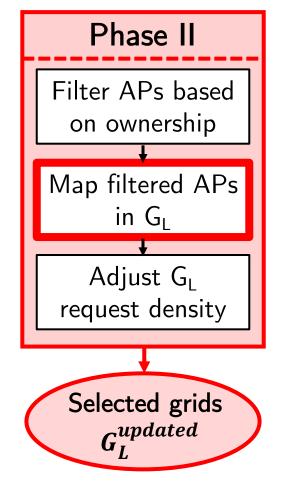


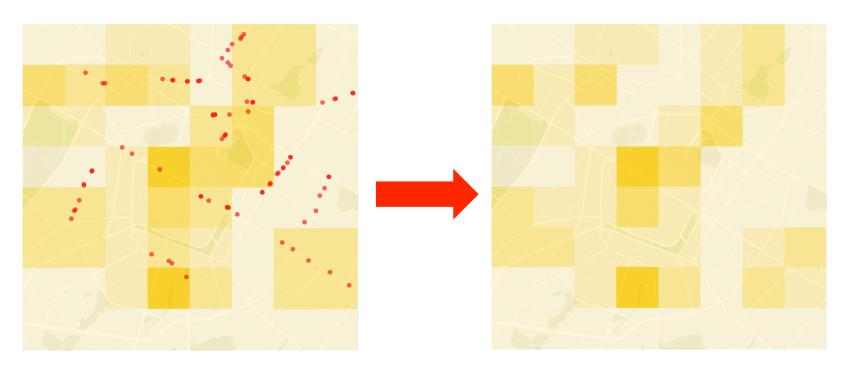




Map all filtered access points on Phase I heatmap

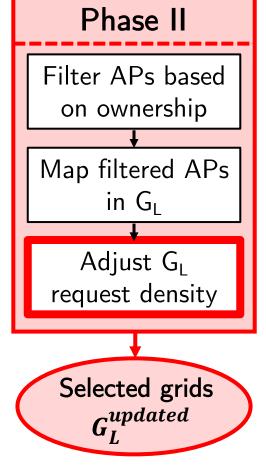
> Cluster nearby access points based on densities

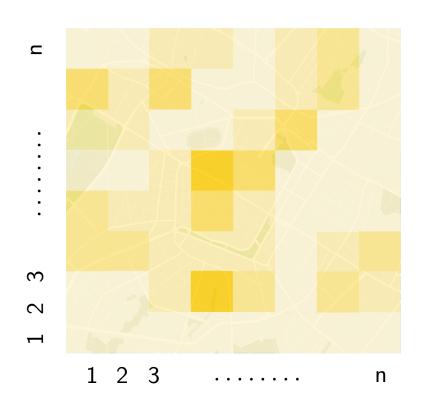




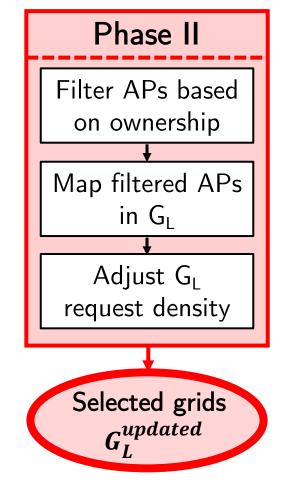
Reduce grid densities based access point availability density

Resulting heatmap denotes grids with overflowing user requests





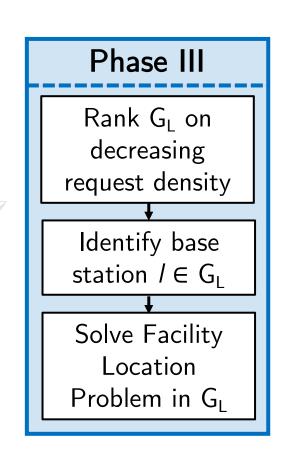
```
G_L^{updated} = \{ g_{11}, d_{11}; g_{12}, d_{12}; g_{13}, d_{13}; \dots g_{nn}, d_{nn} \}
```

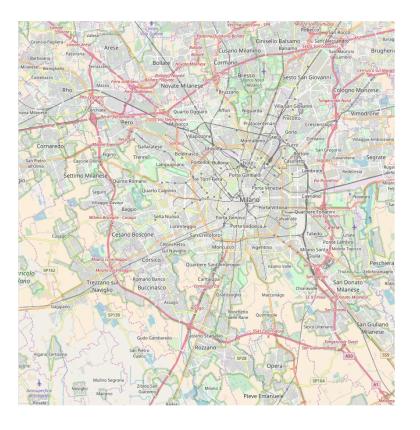


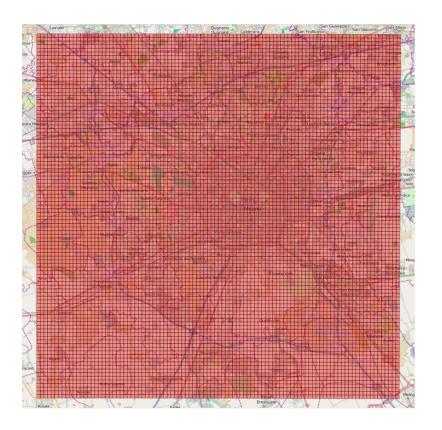
## Phase III: Edge Location Selection

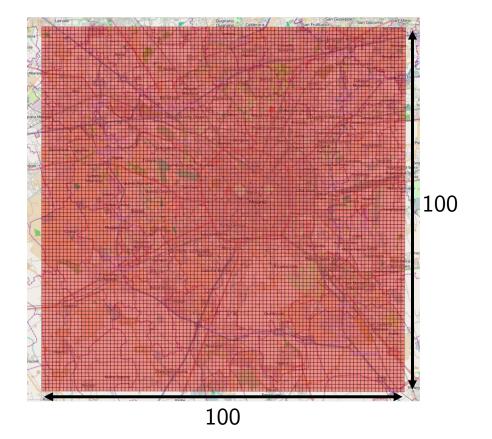
- 1. Select the **best set** of deployment locations considering connectivity to end users
- 2. Base stations are taken as possible deployment locations

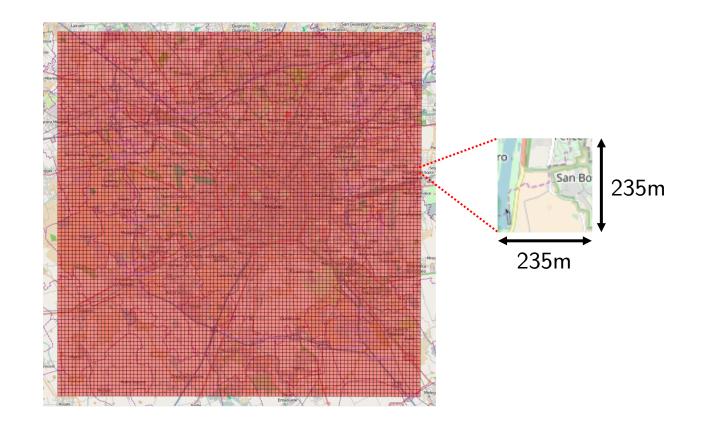
Location selection best resembles Facility Location Problem (FLP)











Evaluate Anveshak's placement of k' edge servers on n' possible base-stations in Milan, Italy

#### A. With real datasets

- 1. User requests: Published by Telecom Italia
- 2. Service-provider edge locations: Published by OpenCellid
- 3. User-managed edge locations: Published by WiGLE

Evaluate Anveshak's placement of k' edge servers on n' possible base-stations

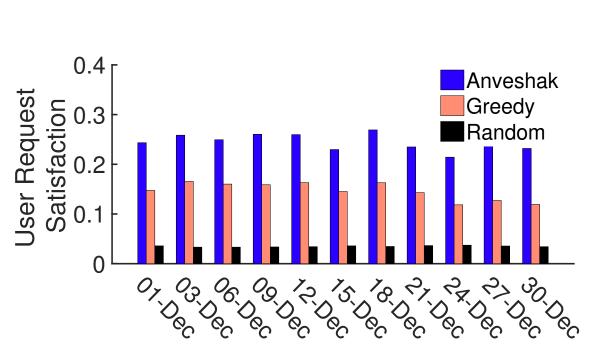
#### B. With two approaches

- 1. **Greedy:** Allocates user request densities to grids and selects top-k maximum serving base-stations
- 2. Random: Randomly chooses k valid base-stations

#### Setup

- 1. User is mapped to nearest base station using coordinate based latency approximation
- 2. Deployment framework selects **50** edge location out of 850+

### **User Request Satisfaction**

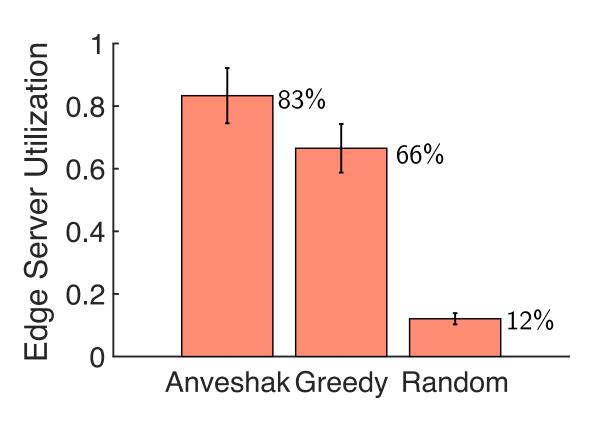


Q. How many user requests are handled by *Anveshak's* placed edge servers?

- 1. 67% more requests than *Greedy*
- 2. 25% of total requests handled by 8% of selected base stations

Can achieve 90% user satisfaction by placing 124 servers over 218 and 300+ by *Greedy* and *Random* 

#### **Server Utilization**



Q. How busy are the deployed servers?

Assumption: All user requests in a grid are first handled by user-managed edge

More user request density areas have more user-managed edge resources leading to less utilization of deployed server

#### **Conclusion**

- 1. Anveshak is a deployment framework designed to assist service providers
- 2. It efficiently identifies optimal locations for edge server placement while considering:
  - Density of user requests
  - > Density of future deployment of user-managed edge resources
- 3. We evaluate Anveshak and other deployment algorithms on real datasets
- 4. Anveshak achieves 67% increase in user satisfaction with 83% server utilization





# Thank You! Questions?

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## Backup Slides

## How is Anveshak different from CDN server placement?

#### Similarities:

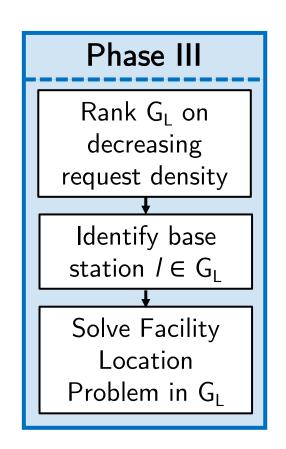
- 1. Both problems must ensure consistent and least connectivity to end-clients
- 2. Both problems must optimize for cost of deployment

#### Differences:

- 1. Unlike content, edge server handles requests which are short-lived and locally-relevant
- 2. Focuses more on server availability and network latency than network bandwidth

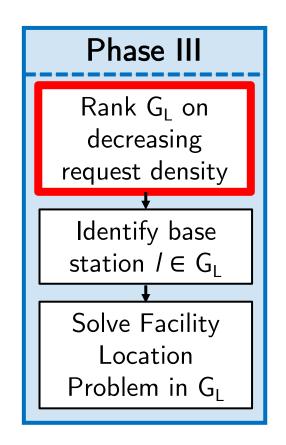
Anveshak assumes that the deployed edge server will be co-located with a base station

- From Phase II heatmap, select exact base station location which satisfies maximum user requests
- ➤ Selected base station must be able to service maximum users in 1-2 hops
- List of exact base station locations can be made available with partnering ISP

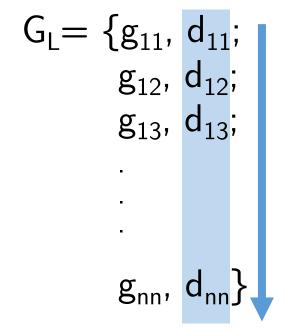


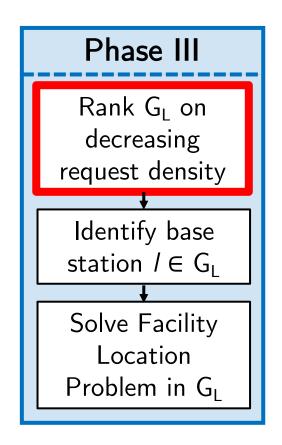
Prioritize the grid locations based on *updated* user request density

```
\begin{aligned} G_L &= \{g_{11}, \ d_{11}; \\ g_{12}, \ d_{12}; \\ g_{13}, \ d_{13}; \\ & \cdot \\ & \cdot \\ g_{nn}, \ d_{nn} \} \end{aligned}
```



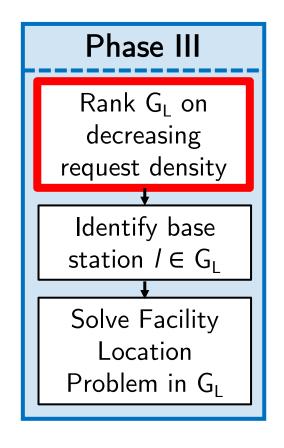
Prioritize the grid locations based on *updated* user request density





Prioritize the grid locations based on *updated* user request density

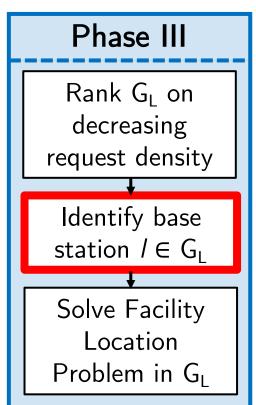
$$\begin{split} G_L &= \{ g_{13}, \ d_{13}; \\ g_{32}, \ d_{32}; \\ g_{61}, \ d_{61;} \\ &\cdot \\ &\cdot \\ g_{n7}, \ d_{n7} \} \end{split}$$

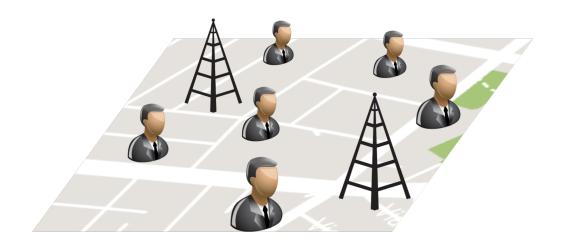


where,

$$G_L\{d_1\} > G_L\{d_2\} > ... > G_L\{d_{last}\}$$





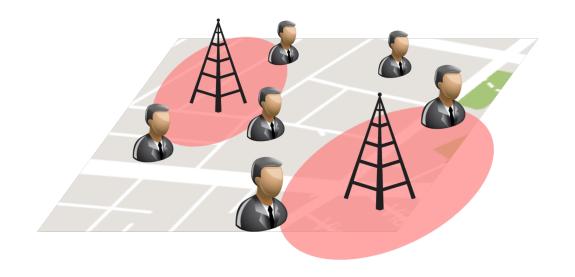


Select optimal base station  $(S_I)$  in  $G_L[i]$  satisfying maximum user requests [Phase I]

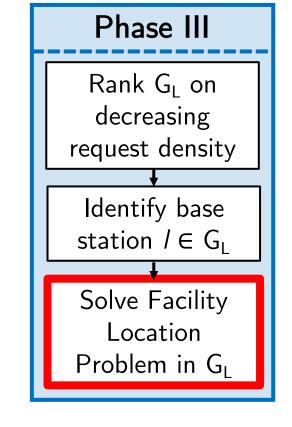
Rank G<sub>I</sub> on decreasing request density Identify base station  $l \in G_1$ Solve Facility Location Problem in G<sub>1</sub>

Phase III

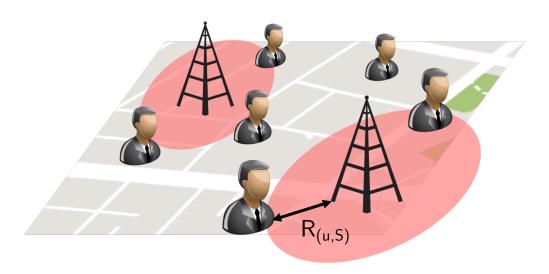
> We focus on **location** and not density



Select optimal base station  $(S_I)$  in  $G_L[i]$  satisfying maximum user requests [Phase I]

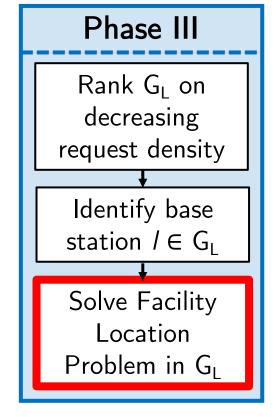


> We focus on **location** and not density



Network cost of connectivity  $(n_{(u,S)})$ 

$$\mathsf{n}_{(\mathsf{u},\mathsf{S})} = \alpha * \mathsf{R}_{(\mathsf{u},\mathsf{S})}$$

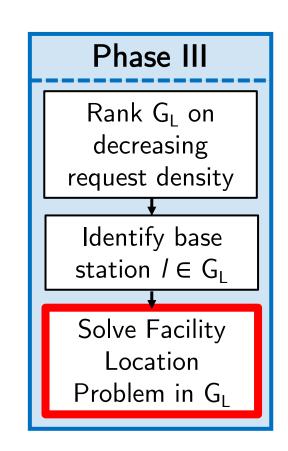


where, 
$$R_{(u_l,S_l)}^{max} = \max[\mathbf{u} - \mathbf{S}_l] \quad \forall \ \mathbf{u}_l \in \mathbf{U} \quad \& \ \mathbf{l} \in \mathbf{L}$$
  $\alpha = \max \max \text{ access cost}$ 

Objective: Minimize one-hop latency between users and selected server location

$$S_u = \min \sum_{l \in L} \{S_l \mid S_l \in S, \ n_{(u,S_l)} < n_{max}\} x_l$$

where,  $x_l \in \{0,1\}$ 

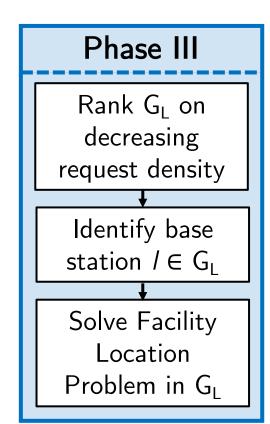


$$S_u = \min \sum_{l \in L} \{S_l \mid S_l \in S, n_{(u,S_l)} < n_{max}\} x_l$$

- ➤ NP-hard problem
- > Approximation-based solvers available [ODL, Sitation]

#### Anveshak's Advantage:

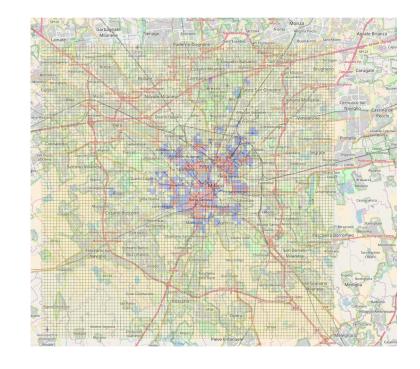
- > Grid size as additional constraint, reduce problem size
- Due to problem's future outlook, exact optimization is **not** required!



#### **Datasets**

#### (1/3) User Request Dataset

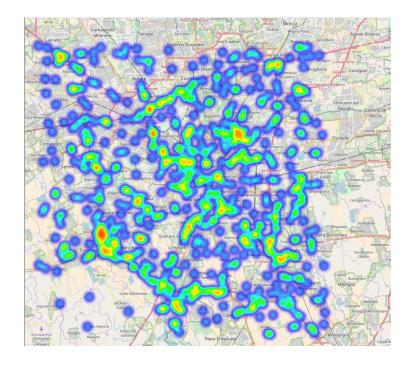
- Published by Telecom Italia\* for Milan, Italy
- Anonymized details of Call Detail Records (CDRs), internet connectivity of users in the region
- Data of November 1<sup>st</sup> to December 31<sup>st</sup> 2013



#### **Datasets**

#### (2/3) Service-provider Edge Dataset

- Published by OpenCellid\*
- Details of cellular base stations such as connectivity type (3G/4G), area etc. along with their GPS coordinates
- Post filteration: 800+ LTE base stations
- Each base station is associated to the map grid



#### **Datasets**

#### (3/3) User-managed Edge Dataset

- Published by WiGLE\*
- Details of WIFi access points such as SSID, signal strength, channel number etc. along with GPS coordinates
- 800+ base stations Milan, Italy serving LTE connection

