

Understanding the Long Tail Latency of TCP in Large-Scale Cloud Networks

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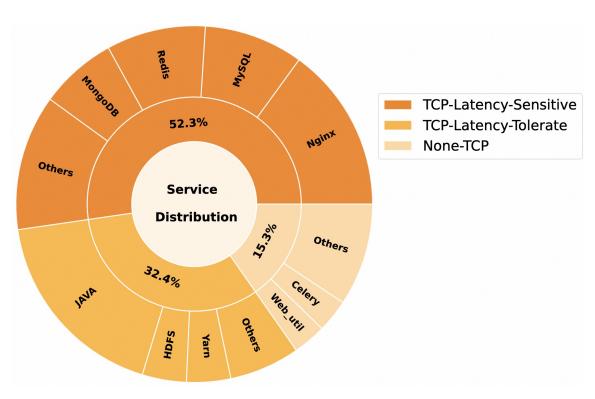






Why TCP's Tail Latency Matters in the Cloud

- TCP dominates ECS:
 - 85%+ of workloads rely on it
- 60% of TCP services are tail-sensitive
 - (e.g., Redis)
- Long tail latency
 - → SLA violation
 - → Revenue loss

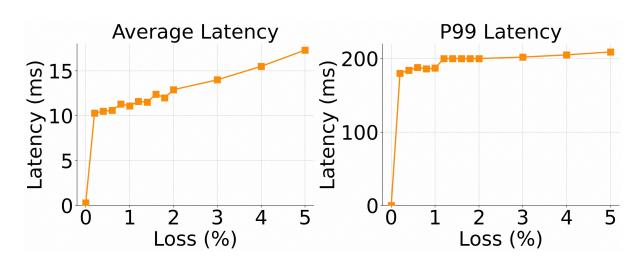


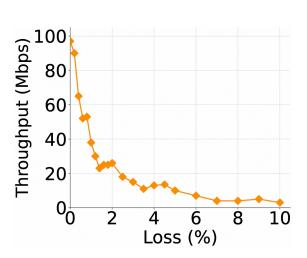




Why TCP's Tail Latency Gets Worse in Cloud

- Single-path transmission can overload paths
- Loss detection slow (RTO ~200ms, triple dup ACK)
- Congestion control reacts late (loss-driven)





Empirical: 1% loss → P99 latency increases 150×







Why is TCP's Tail Latency So Hard to Eliminate

Cloud-scale:

O(1M) of links,

O(100k) servers per

region

Network Instabilities are common and inevitable

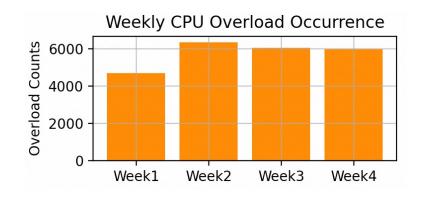


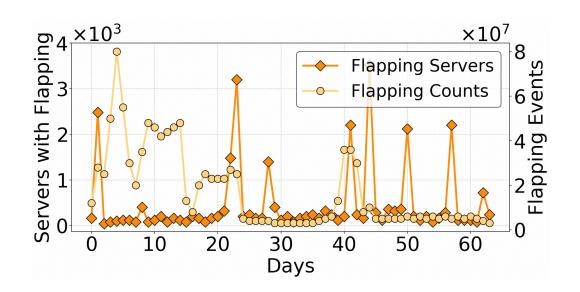


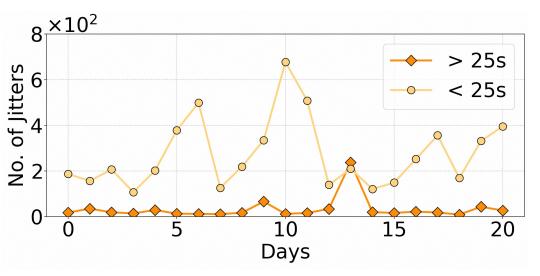


Root Causes: Network Instabilities

- Elephant flows: thousands of times/week
- NIC flapping: millions of times/day
- Network jitter: hundreds of times/day







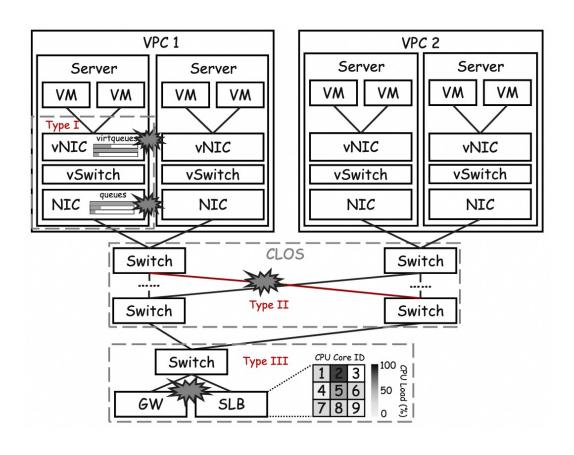






Result of Instability: Packet Loss

- Packet loss occurs at multiple points
 - Type I: Packet loss in a single server
 - Physical NIC
 - Front-end (vNIC) and back-end (vSwitch)
 - Type II: Packet loss in physical networks
 - Physical Link
 - Switch
 - Type III: Packet loss in Middlebox
 - Gateway, Load balancer, NAT, ...







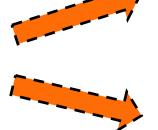


Result of instability: Packet Loss

Physical NIC

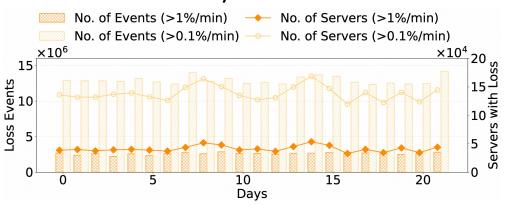
Type I: Packet loss

in a single server



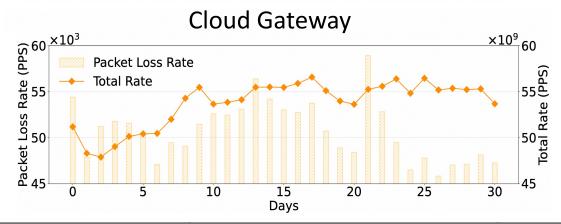
Type III: Packet loss in Middlebox





vNIC front-end and vSwitch back-end

| Solution | Average | P90 | P99 | P999 |
|-----------------|---------|-------|-------|------|
| Kernel-based VM | 0.015% | 0% | 0.17% | 2.0% |
| SmartNIC I | 0.006% | 0% | 0.02% | 1.0% |
| SmartNIC II | 0.072% | 0.03% | 0.29% | 2.0% |







Why Existing Solutions Fall Short

Our Goal

Mitigating long tail latency in unstable, large-scale cloud networks while maintaining complete transparency to end users

Limitations

Limited Performance Improvement

Coarse-grained multipath^{[1] [2]} Lack of receiver-side reordering^{[3][4]} Random path selection^{[1][5]}

Intrusiveness to Users

Dependency on ECN^{[6][7]} Kernel modifications at end hosts^{[1][8]}

Poor Compatibility and Scalability

Custom switch functionalities^{[9][10]} Centralized control plane^{[11][12]}

^[1] Qureshi et al. PLB: Congestion signals are simple and effective for network load balancing. SIGCOMM 2022, pp. 207–218.

^[2] Google Cloud. Introducing Falcon: A reliable, low-latency hardware transport. Google Cloud Blog, 2023.

^[3] Shalev et al. The Tail at AWS Scale. IEEE Micro, 2024.

^[4] Le et al. STrack: A Reliable Multipath Transport for AI/ML Clusters. arXiv:2407.15266, 2024.

^[5] Vanini et al. Let it flow: Resilient asymmetric load balancing with flowlet switching. NSDI 2017. pp. 407–420.

^[6] Katta et al. Clove: Congestion-aware load balancing at the virtual edge. CoNEXT 2017, pp. 323–335.

^[7] Kabbani et al. Flowbender: Flow-level adaptive routing for improved latency and throughput in datacenter networks. CoNEXT 2014, pp. 149–160.

^[8] Ford et al. TCP extensions for multipath operation with multiple addresses. Technical Report, 2013.

^[9] Alizadeh et al. CONGA: Distributed congestion-aware load balancing for datacenters. SIGCOMM 2014, pp. 503-514.

^[10] Song et al. Network Load Balancing with In-network Reordering Support for RDMA. SIGCOMM 2023, pp. 816–831.

^[11] Al-Fares et al. Hedera: Dynamic flow scheduling for data center networks. NSDI 2010, pp. 89–92.

^[12] Curtis et al. Mahout: Low-overhead datacenter traffic management using end-host-based elephant detection. IEEE INFOCOM 2011, pp. 1629–1637.





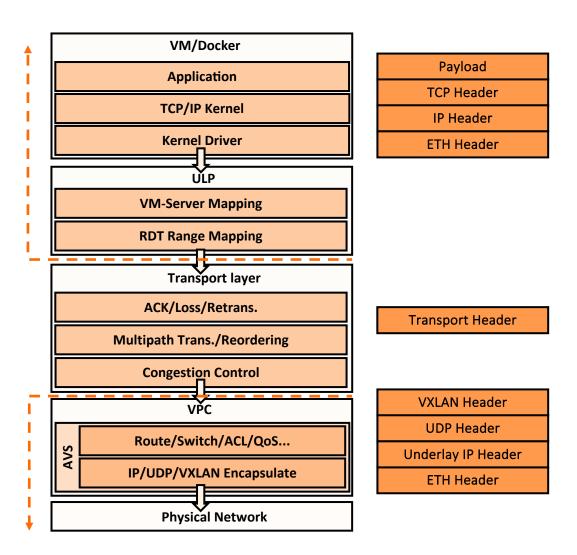


Bifrost

- RTT-Aware multipath transmission
- Hybrid hardware-software reordering!
- ACK aggregation via delayed bitmap



- 1. High performance guarantee
- 2. Non-intrusive to users
- 3. W/o requiring support from network devices

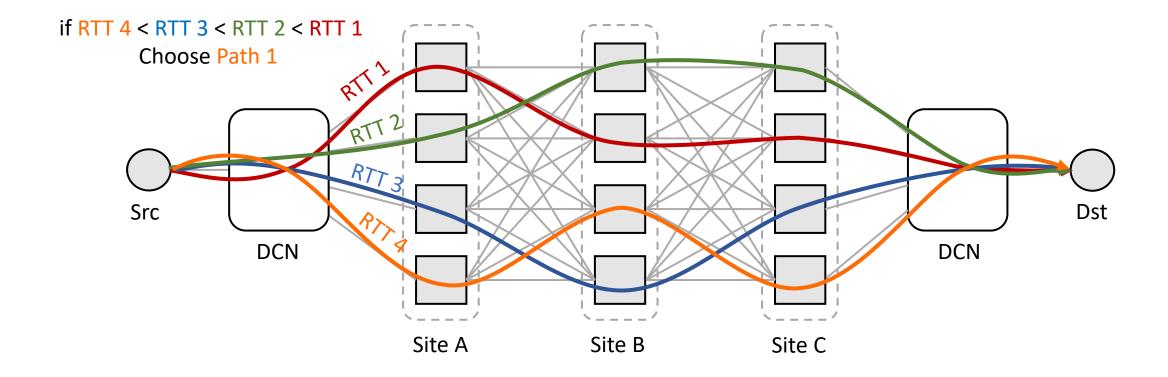








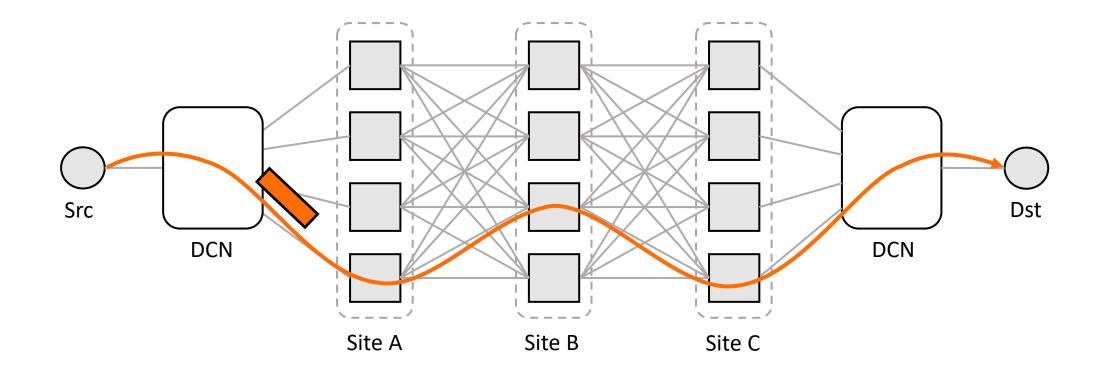
- Partition flow into equal-sized packet groups for scheduling
- Dynamically select the lowest RTT paths across groups





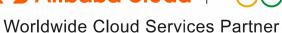


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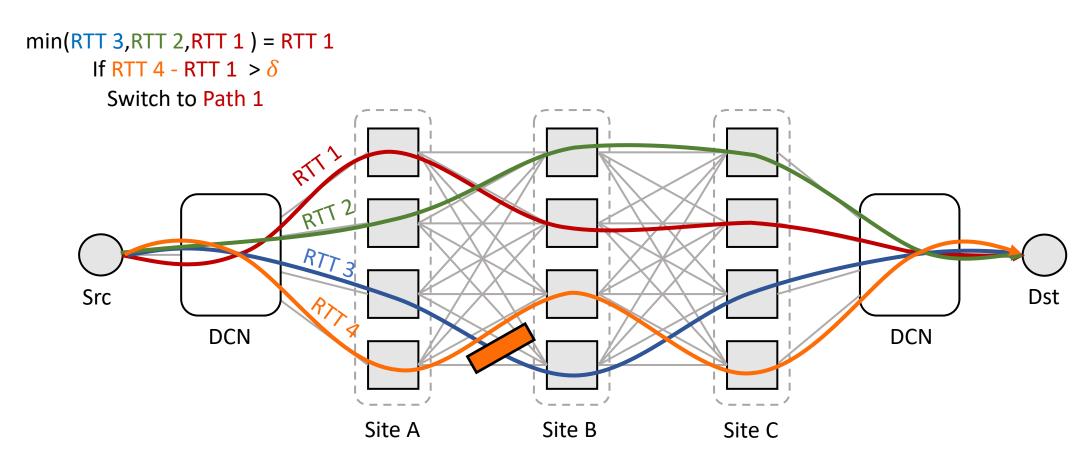






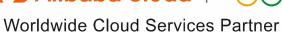


Path stickiness optimization

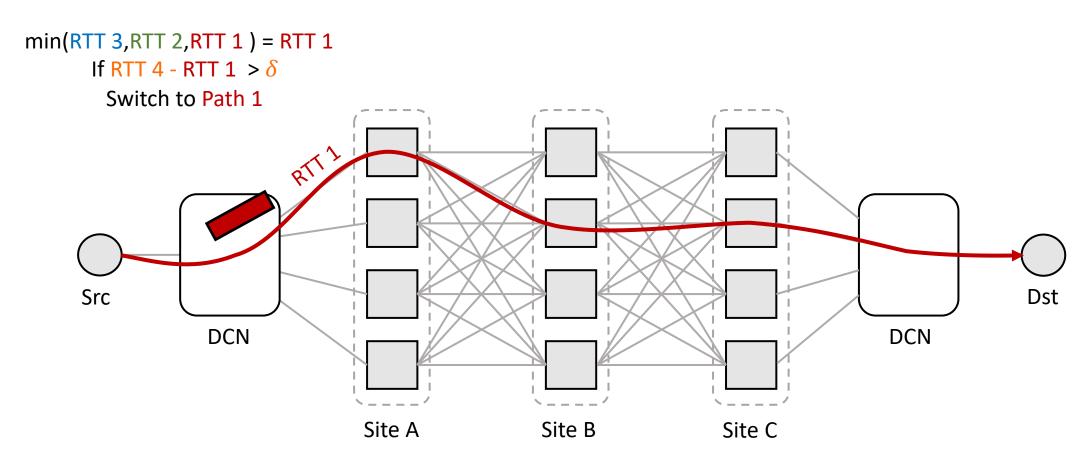








Path stickiness optimization

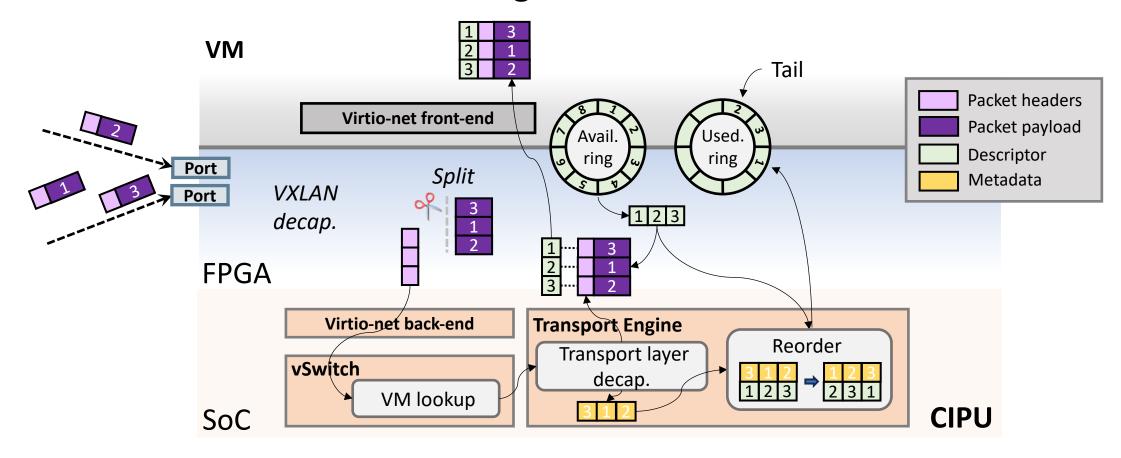




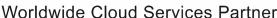


Hybrid Hardware-Software Reordering

- Hardware receives packets
- Software enforces reordering via metadata.







ACK Aggregation Via Delayed Bitmap

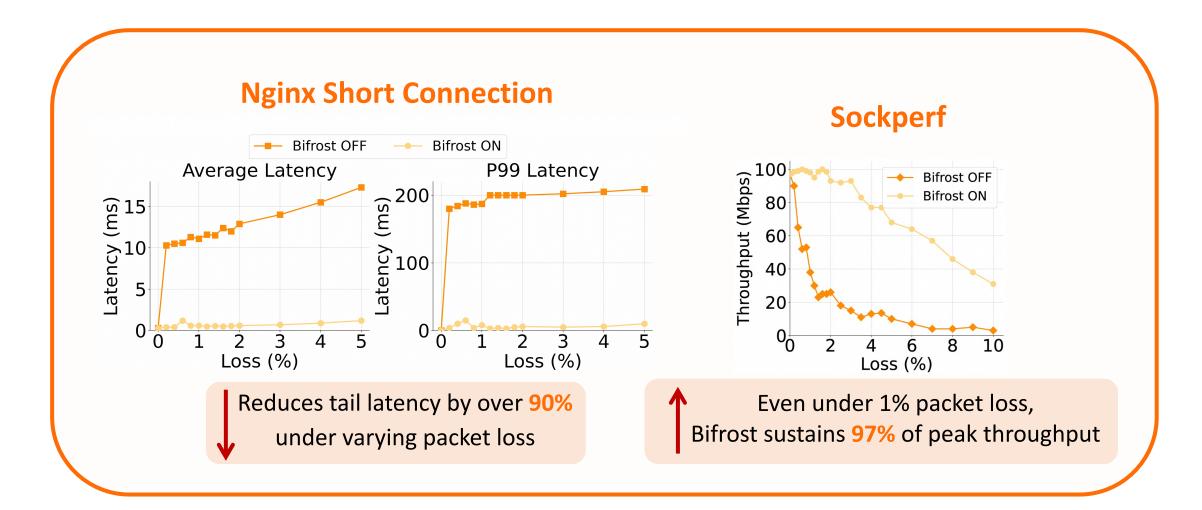
- Delayed ACK generation to capture vNIC losses
 - ACK is sent only after packet reaches the VM, not at physical NIC reception
- ACK format follows bitmap ACK standard
 - Sender triggers fast retrans. upon detecting gaps in the bitmap
 - ACK(3|0100) -> retrans. packet 3
- Precise Retransmission Timeout (RTO)
 - RTO is set to 4ms, roughly 2× RTT in data center networks
 - Significantly faster than traditional TCP RTO (200ms)







Preliminary Evaluation







Conclusion

- Cloud-scale network instability causes frequent tail latency spikes in TCP services, impacting SLA-critical applications like Redis.
- We present Bifrost, a scalable and non-intrusive transport layer that combines RTT-aware multipath, hybrid reordering and delay ACK.
- Evaluation shows Bifrost reduces P99 latency by >90% and sustains 97% throughput under loss, significantly improving both transport and application performance.

Background & Motivation

Challeng

Bifrost Desig

Preliminary Evaluation

Conclusion



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