

Greedy Forwarding for Mobile Social Networks Embedded in Hyperbolic Spaces

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ABSTRACT

In this work, we design and evaluate a novel greedy forwarding algorithm using metrics in hyperbolic spaces. Hyperbolic geometry has a natural topological reflection of scale-free networks, and greedy algorithm failed in Euclidean space becomes possible in hyperbolic one. We show that mobile social networks can be successfully embedded in such spaces, and obtains competitive performance in terms of message delivery ratio and cost. Under this result, we thus intuitively reveal the fundamental reason that why the famous BUBBLE Rap achieves the optimal performance.

Categories and Subject Descriptors

C.2.0 [Computer-Communication Networks]: General;
C.4 [Performance of Systems]: [Modeling techniques, Performance attributes]

Keywords

Greedy forwarding, mobile social networks, hyperbolic spaces

1. INTRODUCTION

We evidence a world with a great amount of portable devices carried by people forming the mobile social networks[1, 2]. In this work, we propose a message forwarding method in such environment by mapping the users into a metric space and then deploying greedy forwarding algorithm. In order to design forwarding mechanisms that enhance the message delivery possibility in a limited time interval without wasting network resources, one should learn from users' contact history. Several studies developed forwarding methods based on social and temporal factors gained by contact history, such as [3, 4].

We focus on an alternative approach using greedy forwarding algorithm with nodes mapped into 2D hyperbolic space as geographic metrics. Distance between nodes represents contact probabilities in the future. Then greedy algorithm is implemented by choosing next hop which maximize the distance to the destination [5]. This simple but robust method only requires devices to store location information,

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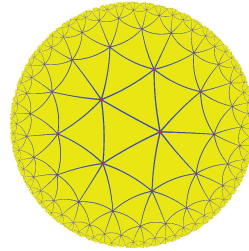


Figure 1: Hyperbolic Space

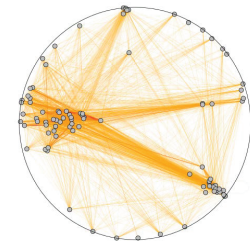


Figure 2: Mapping results of Reality Trace

meanwhile still achieves high message delivery ratio. The fundamental reason makes this method work is the relationship between hyperbolic space and scale-free networks. We know that even low-dimensional hyperbolic spaces are akin to infinite-dimensional Euclidean ones. Fig. 1 is an example of the hyperbolic plane in Poincaré disk model, the edges in the figure are actually the same length in hyperbolic space.

If ignoring the transient period and treat user contacts as encounter processes, we can form a probability graph that the coefficient p_{ij} of edge joins node i, j represents the contact probability between i, j , and $\sum_j p_{ij}$ represents the expected degree of node i . Such a graph is similar with regular scale-free graph. Therefore, it is possible to implement this mapping method as a metric of greedy forwarding algorithm in mobile social networks.

In BUBBLE Rap[3], they designed social-based forwarding algorithms that utilizes community and centrality as metrics and obtained high delivery performance. Meanwhile, there is no theoretical analysis of why this greedy method works. We will show in Section 3 that our forwarding mechanism behaves almost the same as BUBBLE Rap. That is to say, since our greedy method has been proved successfully working in hyperbolic spaces theoretically [6], BUBBLE Rap is also proved to be close to optimal solutions.

2. METHOD

2.1 Hyperbolic model

Boguñá et al.[6] proposed a method to synthesize network with any desired number of nodes N , average degree k , average clustering \bar{C} and power-law distribution exponent $\gamma \geq 2$ by distributing N nodes uniformly (under the hyperbolic metric) within a hyperbolic disc and link them with probability which decreases with the hyperbolic distance between them. Following this approach, we first distribute angular coordinates uniformly in $[0, 2\pi)$ and distribute radial coordinates with density $\rho(r) = \alpha e^{\alpha(r-R)}$, where $\alpha = (\gamma - 1)/2$

and hyperbolic dist radius $R = 2 \ln \left(\frac{N}{c} \right)$, where

$$c = \bar{k} \frac{\sin(\pi T)}{2T} \left(\frac{\gamma - 2}{\gamma - 1} \right)^2,$$

and T controls clustering. Then, connect each pair of nodes i and j with probability

$$p(x_{ij}) = \frac{1}{1 + e^{(x_{ij}-R)/2T}},$$

where x_{ij} is the hyperbolic distance satisfying

$$\cosh x_{ij} = \cosh r_i \cosh r_j - \sinh r_i \sinh r_j \cos \Delta\theta_{ij}.$$

It can be proved that the generated graph follows power-law distribution.

2.2 Maximum likelihood estimation

We fit the probability network to the model above by finding set of coordinates $\{r_i, \theta_i\}$ maximizing the likelihood that coordinates generated by the hyperbolic model match the probability network $\{p_{ij}\}$. This likelihood is

$$\mathcal{L}(\{r_i, \theta_i\} | \{p_{ij}\}) = \frac{P(\{r_i, \theta_i\}) \mathcal{L}(\{p_{ij}\} | \{r_i, \theta_i\})}{P(\{p_{ij}\})}, \quad (1)$$

where

$$P(\{r_i, \theta_i\}) = \frac{1}{(2\pi)^N} \prod_{i=1}^N \rho(r_i)$$

and

$$\mathcal{L}(\{p_{ij}\} | \{r_i, \theta_i\}) = \prod_{i < j} [p(x_{ij}) p_{ij} + (1 - p(x_{ij}))(1 - p_{ij})].$$

Maximization of the likelihood in 1 is solved by the following Metropolis-Hastings algorithm.

Algorithm 1 Maximum Likelihood Estimation

```

calculate parameters  $N, \{p_{ij}\}, \bar{k}, \gamma, T$ 
set  $\{r_i, \theta_i\}$  randomly uniformly on the disk
while likelihood not converge do
  for  $i$  in all nodes do
    set  $r_i, \theta_i$  that maximize  $\mathcal{L}$ 
  end for
end while

```

2.3 Greedy forwarding

In the greedy forwarding, nodes simply choose the next hop from neighbors which maximize the contact probability

$$p = 1 + \frac{\lambda_2}{\lambda_1 - \lambda_2} e^{-\lambda_1} + \frac{\lambda_1}{\lambda_2 - \lambda_1} e^{-\lambda_2},$$

where λ_1 denotes intensity of encounter process from the current node to the next node, and λ_2 from the next node to the destination node which can be derived from $p(x_{ij})$ [7].

3. RESULTS

We use two real-world human mobile traces of **Reality** for the performance evaluation. We show the mapping result of **Reality** in Fig. 2, where different colors represent the contact frequency. The result show that there are two big mobile communities and several small ones in **Reality**. Node pairs that interact often are mapped close to each other on the hyperbolic plane, with hot nodes located near the center. If using the greedy algorithm to forward message from

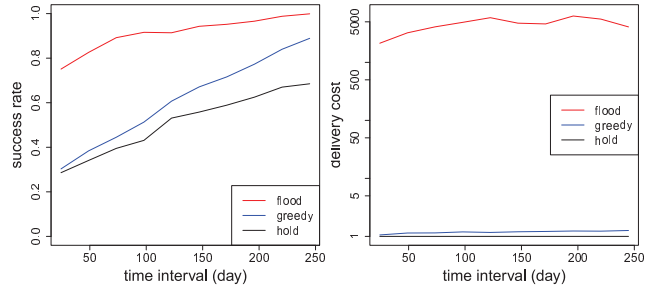


Figure 3: Forwarding results of **Reality**

one node to another, the forward path naturally follows *popular node at local community – popular node at the destination’s community – destination* sequence, which is similar with BUBBEL Rap. Thanks to the huge space of the hyperbolic plane, we are able to locate nodes almost uniformly (in hyperbolic space). Thus, greedy forwarding is possible.

Fig. 3 shows the forwarding results, where 1,000 messages with random source and destination are generated randomly for each time interval. The success ratio stands for the proportion of the messages that are sent to the destination within specified time interval. Meanwhile, the delivery cost stands for the average number of copies generated during the forwarding process. *Hold* algorithm stores the message until source node meet destination node, which represents the lower bound of delivery cost. *Flood* algorithm flood message to every node encountered, which represents the upper bound of success ratio. The results show that our algorithm performs well with few forwarding cost, especially when time interval increases.

4. REFERENCES

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